



Contents lists available at ScienceDirect

Journal of Algebra

www.elsevier.com/locate/jalgebra



Stability of depths of powers of edge ideals



Tran Nam Trung

Institute of Mathematics, VAST, 18 Hoang Quoc Viet, Hanoi, Viet Nam

ARTICLE INFO

Article history:

Received 15 July 2013

Available online 4 February 2016

Communicated by Bernd Ulrich

MSC:

13D45

05C90

05E40

05E45

Keywords:

Depth

Monomial ideal

Stanley–Reisner ideal

Edge ideal

Simplicial complex

Graph

ABSTRACT

Let G be a graph and let $I := I(G)$ be its edge ideal. In this paper, we provide an upper bound of n from which $\text{depth } R/I(G)^n$ is stationary, and compute this limit explicitly. This bound is always achieved if G has no cycles of length 4 and every its connected component is either a tree or a unicyclic graph.

© 2016 Elsevier Inc. All rights reserved.

Introduction

Let $R = K[x_1, \dots, x_r]$ be a polynomial ring over a field K and I be a homogeneous ideal in R . Brodmann [2] showed that $\text{depth } R/I^n$ is a constant for sufficiently large n . Moreover

$$\lim_{n \rightarrow \infty} \text{depth } R/I^n \leq \dim R - \ell(I),$$

E-mail address: tntrung@math.ac.vn.

<http://dx.doi.org/10.1016/j.jalgebra.2016.01.009>

0021-8693/© 2016 Elsevier Inc. All rights reserved.

where $\ell(I)$ is the analytic spread of I . It was shown in [6, Proposition 3.3] that this is an equality when the associated graded ring of I is Cohen–Macaulay. We call the smallest number n_0 such that $\text{depth } R/I^n = \text{depth } R/I^{n_0}$ for all $n \geq n_0$, the *index of depth stability* of I , and denote this number by $\text{dstab}(I)$. It is of natural interest to find a bound for $\text{dstab}(I)$. As until now we only know effective bounds of $\text{dstab}(I)$ for few special classes of ideals I , such as complete intersection ideals (see [5]), square-free Veronese ideals (see [8]), polymatroidal ideals (see [10]). In this paper we will study this problem for *edge ideals*.

From now on, every graph G is assumed to be simple (i.e., a finite, undirected, loopless and without multiple edges) without isolated vertices on the vertex set $V(G) = [r] := \{1, \dots, r\}$ and the edge set $E(G)$ unless otherwise indicated. We associate to G the quadratic squarefree monomial ideal

$$I(G) = (x_i x_j \mid \{i, j\} \in E(G)) \subseteq R = K[x_1, \dots, x_r]$$

which is called the edge ideal of G .

If I is a polymatroidal ideal in R , Herzog and Qureshi proved that $\text{dstab}(I) < \dim R$ and they asked whether $\text{dstab}(I) < \dim R$ for all Stanley–Reisner ideals I in R (see [10]). For a graph G , if every its connected component is nonbipartite, then we can see that $\text{dstab}(I(G)) < \dim R$ from [4]. In general, there is no an absolute bound of $\text{dstab}(I(G))$ even in the case G is a tree (see [20]). In this paper we will establish a bound of $\text{dstab}(I(G))$ for any graph G . In particular, $\text{dstab}(I(G)) < \dim R$.

The first main result of the paper shows that the limit of the sequence $\text{depth } R/I(G)^n$ is the number s of connected bipartite components of G and $\text{depth } R/I(G)^n$ immediately becomes constant once it reaches the value s . Moreover, $\text{dstab}(I(G))$ can be obtained via its connected components.

Theorem 4.4. *Let G be a graph with p connected components G_1, \dots, G_p . Let s be the number of connected bipartite components of G . Then*

- (1) $\min\{\text{depth } R/I(G)^n \mid n \geq 1\} = s$.
- (2) $\text{dstab}(I(G)) = \min\{n \geq 1 \mid \text{depth } R/I(G)^n = s\}$.
- (3) $\text{dstab}(I(G)) = \sum_{i=1}^p \text{dstab}(I(G_i)) - p + 1$.

The second one estimates an upper bound for $\text{dstab}(I(G))$. Before stating our result, we recall some terminologies from graph theory. In a graph G , a *leaf* is a vertex of degree one and a *leaf edge* is an edge incident with a leaf. A connected graph is called a *tree* if it contains no cycles, and it is called a *unicyclic* graph if it contains exactly one cycle. We use the symbols $v(G)$, $\varepsilon(G)$ and $\varepsilon_0(G)$ to denote the number of vertices, edges and leaf edges of G , respectively.

Theorem 4.6. *Let G be a graph. Let G_1, \dots, G_s be all connected bipartite components of G and let G_{s+1}, \dots, G_{s+t} be all connected nonbipartite components of G . Let $2k_i$ be the*

maximum length of cycles of G_i ($k_i := 1$ if G_i is a tree) for all $i = 1, \dots, s$; and let $2k_i - 1$ be the maximum length of odd cycles of G_i for every $i = s + 1, \dots, s + t$. Then

$$\text{dstab}(I(G)) \leq v(G) - \varepsilon_0(G) - \sum_{i=1}^{s+t} k_i + 1.$$

It is interesting that this bound is always achieved if G has no cycles of length 4 and every its connected component is either a tree or a unicyclic graph (see [Theorem 5.1](#)).

Our approach is based on a generalized Hochster formula for computing local cohomology modules of arbitrary monomial ideals formulated by Takayama [\[24\]](#). The efficiency of this formula was shown in recent papers (see [\[7,12,17–19\]](#)). Using this formula and an explicit description of it for symbolic powers of Stanley–Reisner ideals given in [\[17\]](#), we are able to study the stability of depths of powers of edge ideals.

The paper is organized as follows. In [Section 1](#), we give some useful formulas on $\text{dstab}(I(G))$ for the case when all components of G are either nonbipartite or bipartite. We also recall the generalized Hochster formula to compute local cohomological modules of monomial ideals formulated by Takayama. In [Section 2](#) and [Section 3](#) we set up an upper bound of the index of depth stability for connected graphs which are either nonbipartite or bipartite, respectively. The core of the paper is [Section 4](#). There we compute the limit of the sequence $\text{depth } R/I(G)^n$. Then combining with results in [Sections 2](#) and [3](#) on the index of depth stability of connected graphs we obtain a bound of $\text{dstab}(I(G))$ for all any graph G . In the last section, we compute the index of depth stability of trees and unicyclic graphs.

1. Preliminary

We recall some standard notation and terminology from graph theory here. Let G be a graph. The ends of an edge of G are said to be incident with the edge, and vice versa. Two vertices which are incident with a common edge are adjacent, and two distinct adjacent vertices are neighbors. The set of neighbors of a vertex v in G is denoted by $N_G(v)$ and the degree of a vertex v in G , denoted by $\deg_G(v)$, is the number of neighbors of v in G . If there is no ambiguity in the context, we write $\deg v$ instead of $\deg_G(v)$. The graph G is bipartite if its vertex set can be partitioned into two subsets X and Y so that every edge has one end in X and one end in Y ; such a partition (X, Y) is called a bipartition of G . It is well-known that G is bipartite if and only if G contains no odd cycle (see [\[1, Theorem 4.7\]](#)).

Let I be a homogeneous ideal in a polynomial ring $R = K[x_1, \dots, x_r]$ over the field K . As introduced in [\[9\]](#) we define the index of depth stability of I to be the number

$$\text{dstab}(I) := \min\{n_0 \geq 1 \mid \text{depth } S/I^n = \text{depth } S/I^{n_0} \text{ for all } n \geq n_0\}.$$

In this paper we will establish a bound of $\text{dstab}(I(G))$ for any graph G . First we have some information about $\text{dstab}(I(G))$ when every component of G is nonbipartite.

Lemma 1.1. *Let G be a graph with connected components G_1, \dots, G_t . If all these components are nonbipartite, then*

- (1) $\text{dstab}(I(G)) = \min\{n \geq 1 \mid \text{depth } R/I(G)^n = 0\};$
- (2) $\text{dstab}(I(G)) = \sum_{i=1}^t \text{dstab}(I(G_i)) - t + 1.$

Proof. (1) Let $\mathfrak{m}_i := (x_j \mid j \in V(G_i))$ and $R_i := K[x_j \mid j \in V(G_i)]$, i.e., \mathfrak{m}_i is the maximal homogeneous ideal of R_i , for $i = 1, \dots, t$. Let $\mathfrak{m} := (x_j \mid j \in V(G))$ be the maximal homogeneous ideal of R , so that $\mathfrak{m} = \mathfrak{m}_1 + \dots + \mathfrak{m}_t$.

By [4, Corollary 3.4] we have $\mathfrak{m}_i \in \text{Ass}(R_i/I(G_i)^{n_i})$ for some integer $n_i \geq 1$. Let $n_0 := \sum_{i=1}^t (n_i - 1) + 1$. By [4, Corollary 2.2] we have $\mathfrak{m} \in \text{Ass}(R/I(G)^n)$ for all $n \geq n_0$. On the other hand, the sequence $\{\text{Ass}(R/I(G)^n)\}_{n \geq 1}$ is increasing by [15, Theorem 2.15] and note that $\text{depth } R/I(G)^n = 0$ if and only if $\mathfrak{m} \in \text{Ass}(R/I(G)^n)$, this implies $\text{dstab}(I(G)) = \min\{n \geq 1 \mid \text{depth } R/I(G)^n = 0\}$.

(2) By Part 1 we also have $\text{dstab}(I(G_i)) = \min\{n \geq 1 \mid \mathfrak{m}_i \in \text{Ass}(R/I(G_i)^n)\}$ for each component G_i . On the other hand, by [4, Corollary 2.2] we have $\mathfrak{m} \in \text{Ass}(R/I(G)^n)$ if and only if we can write $n = \sum_{i=1}^t (n_i - 1) + 1$ where the n_i are positive integers such that $\mathfrak{m}_i \in \text{Ass}(R_i/I(G_i)^{n_i})$. Thus the statement follows. \square

Next, we consider bipartite graphs. Note that all connected components of such graphs are bipartite as well. Bipartite graphs have a nice algebraic characterization.

Lemma 1.2. (See [22].) *A graph G is bipartite if and only if $I(G)^n = I(G)^{(n)}$ for all $n \geq 1$.*

Using this characterization we obtain.

Lemma 1.3. *Let G be a bipartite graph with s connected components. Then*

- (1) $\min\{\text{depth } R/I(G)^n \mid n \geq 1\} = s$, and
- (2) $\text{dstab}(I(G)) = \min\{n \geq 1 \mid \text{depth } R/I(G)^n = s\}.$

Proof. Since G is bipartite, by Lemma 1.2 we have $I(G)$ is normally torsion-free, and so by [13] the Rees ring $\mathcal{R}[I(G)]$ of $I(G)$ is Cohen–Macaulay. Then by [14] the associated graded ring of $I(G)$ is Cohen–Macaulay as well. Hence, by [6, Proposition 3.3] we have

- (1) $\min\{\text{depth } R/I(G)^n \mid n \geq 1\} = r - \ell(I(G))$, and
- (2) $\text{dstab}(I(G)) = \min\{n \geq 1 \mid \text{depth } R/I(G)^n = r - \ell(I(G))\}.$

On the other hand, $r - \ell(I(G)) = s$ (see [25, Page 50]). Thus the lemma follows. \square

In the general case, our main tool to study $\text{dstab}(I(G))$ is a generalized version of a Hochster's formula (see [23, Theorem 4.1 in Chapter II]) to compute local cohomology modules of monomial ideals given in [24].

Let $\mathfrak{m} := (x_1, \dots, x_r)$ be the maximal homogeneous ideal of R and I a monomial ideal in R . Since R/I is an \mathbb{N}^r -graded algebra, $H_{\mathfrak{m}}^i(R/I)$ is an \mathbb{Z}^r -graded module over R/I . For every degree $\alpha \in \mathbb{Z}^r$ we denote by $H_{\mathfrak{m}}^i(R/I)_{\alpha}$ the α -component of $H_{\mathfrak{m}}^i(R/I)$.

Let $\Delta(I)$ denote the simplicial complex corresponding to the Stanley–Reisner ideal \sqrt{I} . For every $\alpha = (\alpha_1, \dots, \alpha_r) \in \mathbb{Z}^r$ we set $G_{\alpha} := \{i \mid \alpha_i < 0\}$ and we denote by $\Delta_{\alpha}(I)$ the simplicial complex of all sets of the form $F \setminus G_{\alpha}$, where F is a face of $\Delta(I)$ containing G_{α} such that for every minimal generator x^{β} of I there exists an $i \notin F$ such that $\alpha_i < \beta_i$. To represent $\Delta_{\alpha}(I)$ in a more compact way, for every subset F of $[r]$ let $R_F := R[x_i^{-1} \mid i \in F \cup G_{\alpha}]$ and $I_F := IR_F$. This means that the ideal I_F of R_F is generated by all monomials of I by setting $x_i = 1$ for all $i \in F \cup G_{\alpha}$. Then $x^{\alpha} \in R_F$ and by [7, Lemma 1.1] we have

$$\Delta_{\alpha}(I) = \{F \subseteq [r] \setminus G_{\alpha} \mid x^{\alpha} \notin I_F\}. \quad (1)$$

Lemma 1.4. (See [24, Theorem 1].) $\dim_K H_{\mathfrak{m}}^i(R/I)_{\alpha} = \dim_K \tilde{H}_{i-|G_{\alpha}|-1}(\Delta_{\alpha}(I); K)$.

Let $\mathcal{F}(\Delta)$ denote the set of facets of Δ . If $\mathcal{F}(\Delta) = \{F_1, \dots, F_m\}$, we write $\Delta = \langle F_1, \dots, F_m \rangle$. The Stanley–Reisner ideal of Δ can be written as (see [16, Theorem 1.7]):

$$I_{\Delta} = \bigcap_{F \in \mathcal{F}(\Delta)} P_F,$$

where P_F is the prime ideal of R generated by variables x_i with $i \notin F$. For every integer $n \geq 1$, the n -th symbolic power of I_{Δ} is the monomial ideal

$$I_{\Delta}^{(n)} = \bigcap_{F \in \mathcal{F}(\Delta)} P_F^n.$$

Note that $\Delta(I_{\Delta}^{(n)}) = \Delta$. In [17, Lemma 1.3] there was given an useful formula for computing $\Delta_{\alpha}(I_{\Delta}^{(n)})$. We apply it to edge ideals.

An independent set in a graph G is a set of vertices no two of which are adjacent to each other. An independent set S in G is maximal if the addition to S of any other vertex in the graph destroys the independence. Let $\Delta(G)$ be the set of independent sets of G . Then $\Delta(G)$ is a simplicial complex and this complex is the so-called independence complex of G ; and facets of $\Delta(G)$ are just maximal independent sets of G . It is easy to see that $I(G) = I_{\Delta(G)}$.

Now we can compute $\Delta_{\alpha}(I(G)^n)$ for bipartite graphs G .

Lemma 1.5. *Let G be a bipartite graph. Then, for all $\alpha \in \mathbb{N}^r$ and $n \geq 1$, we have*

$$\Delta_{\alpha}(I(G)^n) = \left\langle F \in \mathcal{F}(\Delta(G)) \mid \sum_{i \notin F} \alpha_i \leq n-1 \right\rangle.$$

Proof. Let $\Delta := \Delta(G)$. Then, $I_{\Delta} = I(G)$. By Lemma 1.2, we have $I(G)^n = I(G)^{(n)}$. Therefore, $\Delta_{\alpha}(I(G)^n) = \Delta_{\alpha}(I_{\Delta}^{(n)})$. The lemma now follows from [17, Lemma 1.3]. \square

We conclude this section with some remarks about operations on monomial ideals. Let $A := K[x_1, \dots, x_s]$, $B := K[y_1, \dots, y_t]$ and $R := K[x_1, \dots, x_s, y_1, \dots, y_t]$ be polynomial rings where $\{x_1, \dots, x_s\}$ and $\{y_1, \dots, y_t\}$ are two disjoint sets of variables. Then for monomial ideals I, I_1, I_2 of R we have

$$I \cap (I_1 + I_2) = I \cap I_1 + I \cap I_2. \quad (2)$$

Let I_1, I_2 be monomial ideals in A and let J_1, J_2 be monomial ideals in B . For simplicity, we denote $I_s R$ by I_s and $J_s R$ by J_s for $s = 1, 2$, then by [11, Lemma 1.1] we have

$$I_1 J_1 \cap I_2 J_2 = (I_1 \cap I_2)(J_1 \cap J_2). \quad (3)$$

Lemma 1.6. *Let I be a proper monomial ideal of A and J a proper monomial ideal of B . Then, for all $n \geq 1$ we have*

$$\text{depth } R/(I + J)^n \geq \min\{\text{depth } A/I^m \mid 1 \leq m \leq n\}.$$

Proof. Since the case $I = \mathbf{0}$ or $J = \mathbf{0}$ is obvious, so we may assume that I and J are nonzero ideals. For each $i = 0, \dots, n$, we put:

$$W_i := I^i J^{n-i} + \dots + I^n J^0 \subseteq R,$$

where $I^0 = J^0 = R$. Since $W_0 = (I + J)^n$, in order to prove the lemma it suffices to show that

$$\text{depth } R/W_i \geq \min\{\text{depth } A/I^j \mid \max\{i, 1\} \leq j \leq n\} \text{ for all } i = 0, \dots, n. \quad (4)$$

Indeed, if $i = n$, then $\text{depth } R/W_n = \text{depth } R/I^n = \text{depth } A/I^n + t \geq \text{depth } A/I^n$. Next assume that the claim holds for $i + 1$ with $0 \leq i < n$. By Equations (2) and (3) we have $I^i J^{n-i} \cap W_{i+1} = I^{i+1} J^{n-i}$. Since $W_i = I^i J^{n-i} + W_{i+1}$, we have an exact sequence

$$\mathbf{0} \longrightarrow R/I^{i+1} J^{n-i} \longrightarrow R/I^i J^{n-i} \oplus R/W_{i+1} \longrightarrow R/W_i \longrightarrow \mathbf{0}.$$

By Depth Lemma (see, e.g., [3, Proposition 1.2.9]), we have

$$\text{depth } R/W_i \geq \min\{\text{depth } R/I^{i+1} J^{n-i} - 1, \text{depth } R/I^i J^{n-i}, \text{depth } R/W_{i+1}\}.$$

On the other hand, by [11, Lemma 2.2] we have

$$\text{depth } R/I^{i+1}J^{n-i} - 1 = \text{depth } A/I^{i+1} + \text{depth } B/J^{n-i} \geq \text{depth } A/I^{i+1}.$$

Together with the induction hypothesis we then get

$$\text{depth } R/W_i \geq \min\{\text{depth } R/I^iJ^{n-i}, \text{depth } A/I^j \mid j = i+1, \dots, n\}.$$

If $i \geq 1$, by [11, Lemma 2.2] we have

$$\text{depth } R/I^iJ^{n-i} = \text{depth } A/I^i + \text{depth } B/J^{n-i} + 1 \geq \text{depth } A/I^i,$$

which yields the claim.

If $i = 0$, then $\text{depth } R/W_0 \geq \min\{\text{depth } R/J^n, \text{depth } A/I^j \mid j = 1, \dots, n\}$. Note that $\text{depth } R/J^n = s + \text{depth } B/J^n \geq s \geq \text{depth } A/I$, hence the claim also holds. The proof now is complete. \square

2. Depths of powers of edge ideals of connected nonbipartite graphs

Note that for a graph G we always assume that $V(G) = [r]$; $R = K[x_1, \dots, x_r]$ is a polynomial ring over fields K and $\mathfrak{m} = (x_1, \dots, x_r)$ is the maximal homogeneous ideal of R . In this section we always assume that G is a connected nonbipartite graph.

By Lemma 1.1 we have $\text{dstab}(I(G)) = \min\{n \geq 1 \mid \mathfrak{m} \in \text{Ass } R/I(G)^n\}$. Based on [4], we will determine explicitly when $\mathfrak{m} \in \text{Ass } R/I(G)^n$ for a unicyclic graph G .

Recall that a vertex cover (or a cover) of G is a subset S of $V(G)$ such that every edge of G has at least one endpoint in S . A cover is minimal if none of its proper subsets is itself a cover. It is well-known that $P = (x_{i_1}, \dots, x_{i_t})$ is a minimal prime of the edge ideal $I(G)$ if and only if $\{i_1, \dots, i_t\}$ is a minimal cover of G . For a subset U of $V(G)$, the neighbor set of U is the set

$$N(U) := \{v \in V(G) \mid v \text{ is adjacent to some vertex in } U\}.$$

We now describe the process that builds $\text{Ass } R/I(G)^n$ for a unicyclic graph G . Let C be a cycle of G of length $2k - 1$. Let R_k be the set of vertices of C , $B_k := N(R_k) \setminus R_k$ and a monomial

$$d_k := \prod_{i \in R_k} x_i.$$

We now build recursively sets R_n , B_n and a monomial d_n for $n \geq k$. Suppose that $i \in R_s$ and $j \in R_s \cup B_s$ for some $s \geq k$ such that $\{i, j\}$ is an edge of G . Now if $j \in R_s$, then let $R_{s+1} := R_s$ and $B_{s+1} := B_s$. If $j \in B_s$, then let $R_{s+1} := R_s \cup \{j\}$ and $B_{s+1} := (B_s \cup N(j)) \setminus R_{s+1}$. In either case, let $d_{s+1} := d_s(x_i x_j)$.

Now for such a couple (R_n, B_n) with $n \geq k$, we take V to be any minimal subset of $V(G)$ such that $R_n \cup B_n \cup V$ is a cover of G . Then, $(R_n, B_n, V) := (x_i \mid i \in R_n \cup B_n \cup V)$ is an associated prime of $R/I(G)^n$ by [4, Theorem 3.3]. Let P_n be the set of such all prime ideals. Then, by [4, Theorem 5.6] we have

$$\text{Ass } R/I(G)^n = \text{Min}(R/I(G)) \cup P_n. \quad (5)$$

For unicyclic graphs, we have the following observation.

Remark 2.1. Assume that G is a unicyclic graph with a cycle C such that $G \neq C$. For any $v \in V(G) \setminus V(C)$, there is a unique simple path of the form: v_0, v_1, \dots, v_d , where $v_0 \in V(C)$, $v_1, \dots, v_d \notin V(C)$ and $v_d = v$. We say that this path connects C and v . Moreover,

- (1) $d_G(v, C) = d$.
- (2) This simple path can extend to a simple path connecting C to a leaf, i.e., there are vertices u_1, \dots, u_t such that u_s is a leaf and $v_0, v_1, \dots, v_d, u_1, \dots, u_t$ is a simple path.
- (3) If $d_G(v, C)$ is maximal, i.e., $d_G(v, C) \geq d_G(u, C)$ for any $u \in V(G)$, then v is a leaf. Assume further that $d \geq 2$, then $N_G(v_{d-1})$ contains only one non-leaf v_{d-2} .

We now can determine $\text{dstab}(I(G))$ with unicyclic nonbipartite graphs G .

Lemma 2.2. *Let G be a unicyclic nonbipartite graph. If the length of the unique cycle is $2k - 1$, then $\text{dstab}(I(G)) = v(G) - \varepsilon_0(G) - k + 1$.*

Proof. By [4, Corollaries 3.4 and 4.3] we have

$$\mathfrak{m} \in \text{Ass } R/I(G)^n \text{ for all } n \geq v(G) - \varepsilon_0(G) - k + 1.$$

Therefore,

$$\text{depth } R/I(G)^n = 0 \text{ for all } n \geq v(G) - \varepsilon_0(G) - k + 1,$$

so that $\text{dstab}(I(G)) \leq v(G) - \varepsilon_0(G) - k + 1$.

We next prove the converse inequality. It suffices to show that if $\mathfrak{m} \in \text{Ass } R/I(G)^n$, then $n \geq v(G) - \varepsilon_0(G) - k + 1$.

By Equation (5) we deduce that $\mathfrak{m} \in P_n$. Thus, $\mathfrak{m} = (R_n, B_n, V)$ where V is a minimal subset of $V(G)$ such that $R_n \cup B_n \cup V$ is a vertex cover of G . In particular, $V(G) = R_n \cup B_n \cup V$.

Claim 1. $V = \emptyset$. *Indeed, if V contains no leaves of G , then every leaf of G is in either R_n or B_n , and so $R_n \cup B_n = V(G)$ by Remark 2.1. This forces $V = \emptyset$.*

Suppose V contains a leaf, say i . Let j be the unique neighbor of i in G . Then, $j \in V(G) = R_n \cup B_n \cup V$. Therefore, $R_n \cup B_n \cup (V \setminus \{i\})$ is also a vertex cover of G . This contradicts the minimality of V . Hence, $V = \emptyset$, as claimed.

Claim 2. $|B_n| \leq \varepsilon_0(G)$. *Indeed, assume on the contrary that $|B_n| > |\varepsilon_0(G)|$, so that B_n contains a non-leaf of G , say i . Let p be a simple path connecting C and a leaf that passes through i . Let j be a vertex of p after i . Then, by Remark 2.1 and the construction of R_n and B_n we deduce that $j \notin R_n \cup B_n$, so $j \notin V(G)$ by Claim 1, a contradiction. Hence, $|B_n| \leq \varepsilon_0(G)$, as claimed.*

We now prove the lemma. Since $|R_k| = 2k - 1$ and $|R_n| \leq |R_k| + (n - k)$, together with Claim 2 we obtain $v(G) = |R_n| + |B_n| \leq |R_k| + (n - k) + \varepsilon_0(G) = n + k - 1 + \varepsilon_0(G)$, so $n \geq v(G) - \varepsilon_0(G) - k + 1$, as required. \square

Lemma 2.3. *Let G be a unicyclic nonbipartite graph. Assume that the unique odd cycle of G is of length $2k - 1$. Let $n := v(G) - \varepsilon_0(G) - k + 1$. Then, there is a monomial f of R such that $\deg f = 2n - 1$ and $fx_i \in I(G)^n$ for all $i = 1, \dots, r$.*

Proof. By Lemma 2.2 and Equation (5) we have $\mathfrak{m} \in P_n$. Thus, $\mathfrak{m} = (R_n, B_n, V)$ where V is a minimal subset of $V(G)$ such that $R_n \cup B_n \cup V$ is a vertex cover of G . In particular, $V(G) = R_n \cup B_n \cup V$. By the same way as in the proof of Claim 1 in Lemma 2.2 we have $V = \emptyset$. Hence, $R_n \cup B_n = \{1, \dots, r\}$.

Let $f := d_n$. Together with [4, Lemma 3.2] we imply that $\deg(f) = 2n - 1$ and $fx_i \in I(G)^n$ for all $i = 1, \dots, r$, as required. \square

Let G be a connected nonbipartite graph and let $2l - 1$ be the minimum length of odd cycles of G . Then $\text{dstab}(G) \leq v(G) - \varepsilon_0(G) - l + 1$ by [4, Corollaries 3.4 and 4.3]. The following result improves this bound a little bit.

Proposition 2.4. *Let G be a connected nonbipartite graph. Let $2k - 1$ be the maximum length of odd cycles of G . Then, $\text{dstab}(I(G)) \leq v(G) - \varepsilon_0(G) - k + 1$.*

Proof. Let C be an odd cycle of G of length $2k - 1$. If C' is another cycle of G , then C' has an edge e not lying on the cycle C . Delete this edge from G , thereby obtaining a connected subgraph G' of G with $V(G') = V(G)$ and C is a cycle of G' . This process continues until we obtain a connected subgraph H of G such that $V(G) = V(H)$ and H has only one cycle, that is C . Let $n := v(H) - \varepsilon_0(H) - k + 1$. By Lemma 2.3, there is a monomial $f \in R$ such that $\deg f = 2n - 1$ and $x_if \in I(H)^n$ for all $i = 1, \dots, r$. Since $I(H) \subseteq I(G)$, we have

$$x_if \in I(G)^n \text{ for all } i = 1, \dots, r. \quad (6)$$

As $I(G)$ is generated by quadratic monomials and $\deg f = 2n - 1$, so $f \notin I(G)^n$. Together with Equation (6) one has $I(G)^n : f = \mathfrak{m}$. Hence, $\text{depth } R/I(G)^n = 0$, which implies $\text{dstab}(I(G)) \leq n$ by Lemma 1.1. Since $v(G) = v(H)$ and $\varepsilon_0(G) \leq \varepsilon_0(H)$,

$$\text{dstab}(I(G)) \leq n \leq v(G) - \varepsilon_0(G) - k + 1,$$

as required. \square

3. Depths of powers of edge ideals of connected bipartite graphs

Let G be a bipartite graph with bipartition (X, Y) . Clearly, X and Y are then facets of $\Delta(G)$. Assume further that G is connected. By Lemma 1.3, one has $\text{dstab}(I(G))$ is the smallest integer n such that $\text{depth } R/I(G)^n = 1$. For such graphs we can find $\text{dstab}(I(G))$ via integer linear programming.

Lemma 3.1. *Let G be a connected bipartite graph with bipartition (X, Y) and n a positive integer. Then, $\text{depth } R/I(G)^n = 1$ if and only if $\Delta_\alpha(I(G)^n) = \langle X, Y \rangle$ for some $\alpha = (\alpha_1, \dots, \alpha_r) \in \mathbb{N}^r$. Moreover, if $n = \text{dstab}(I(G))$, then such α must satisfy*

$$\sum_{i \notin X} \alpha_i = \sum_{i \notin Y} \alpha_i = n - 1.$$

Proof. Since G is bipartite, by Lemma 1.2 one has $I(G)^n = I(G)^{(n)}$. Hence,

$$\text{depth } R/I(G)^n = \text{depth } R/I(G)^{(n)} \geq 1,$$

and hence $\text{depth } R/I(G)^n = 1$ if and only if $H_m^1(R/I(G)^n) \neq \mathbf{0}$. By [17, Corollary 1.2] this is equivalent to the condition $\Delta_\alpha(I(G)^n)$ being disconnected for some $\alpha = (\alpha_1, \dots, \alpha_r) \in \mathbb{N}^r$.

Therefore, in order to prove the lemma it suffices to show that if $\Delta_\alpha(I(G)^n)$ is disconnected, then $\Delta_\alpha(I(G)^n) = \langle X, Y \rangle$. Indeed, since $\Delta_\alpha(I(G)^n)$ is disconnected, there are two facets F and H of it such that $F \cap H = \emptyset$. Hence, $(V(G) \setminus F) \cup (V(G) \setminus H) = V(G)$. Together with the fact that $X \cap Y = \emptyset$ and $X \cup Y = V(G)$ we get

$$\sum_{i \notin X} \alpha_i + \sum_{i \notin Y} \alpha_i = \sum_{i \in V(G)} \alpha_i \leq \sum_{i \notin F} \alpha_i + \sum_{i \notin H} \alpha_i.$$

Since F and H are members of $\mathcal{F}(\Delta_\alpha(I(G)^n))$, by Lemma 1.5 we have

$$\sum_{i \notin F} \alpha_i \leq n - 1, \quad \text{and} \quad \sum_{i \notin H} \alpha_i \leq n - 1.$$

Therefore,

$$\sum_{i \notin X} \alpha_i + \sum_{i \notin Y} \alpha_i \leq 2(n - 1),$$

which yields

$$\sum_{i \notin X} \alpha_i \leq n-1 \quad \text{or} \quad \sum_{i \notin Y} \alpha_i \leq n-1.$$

Thus we may assume that

$$\sum_{i \notin X} \alpha_i \leq n-1,$$

and thus $X \in \Delta_\alpha(I(G)^n)$ by Lemma 1.5. As $\Delta_\alpha(I(G)^n)$ is disconnected, there is a facet L of $\Delta_\alpha(I(G)^n)$ such that $X \cap L = \emptyset$. We then have $L \subseteq V(G) \setminus X = Y$. The maximality of L forces $L = Y$, hence $Y \in \Delta_\alpha(I(G)^n)$. If $\Delta_\alpha(I(G)^n)$ has another facet, say T , that is different from X and Y , then neither X nor Y contains T , and then T meets both X and Y . This is impossible since $\Delta_\alpha(I(G)^n)$ is disconnected. Hence, $\Delta_\alpha(I(G)^n) = \langle X, Y \rangle$, as claimed.

Finally, assume that $n = \text{dstab}(I(G))$. Then, by Lemma 1.3, n is the smallest positive integer such that $\text{depth } R/I(G)^n = 1$.

Assume that $\sum_{i \notin X} \alpha_i < n-1$ and $\sum_{i \notin Y} \alpha_i < n-1$. Then, $n-1 \geq 1$ and

$$\sum_{i \notin X} \alpha_i \leq (n-1) - 1 \quad \text{and} \quad \sum_{i \notin Y} \alpha_i \leq (n-1) - 1.$$

If F is a facet of $\Delta(G)$ that is different from X and Y , then $F \notin \mathcal{F}(\Delta_\alpha(I(G)))$, and then $\sum_{i \notin F} \alpha_i \geq n > n-1$ according to Lemma 1.5. From these equations and Lemma 1.5, we get $\Delta_\alpha(I(G)^{n-1}) = \langle X, Y \rangle$. In particular, $\Delta_\alpha(I(G)^{n-1})$ is disconnected, so $\text{depth } R/I(G)^{n-1} = 1$. This contradicts to the minimality of n . Thus, we may assume that $\sum_{i \notin Y} \alpha_i = n-1$.

Assume now that $\sum_{i \notin X} \alpha_i < n-1$. Since

$$\sum_{i \in X} \alpha_i = \sum_{i \notin Y} \alpha_i = n-1 \geq 1,$$

$\alpha_i \geq 1$ for some $i \in X$. We may assume that $i = 1$. Let $\beta = (\alpha_1 - 1, \alpha_2, \dots, \alpha_r)$, so that $\beta \in \mathbb{N}^r$ as $\alpha_1 \geq 1$. By the same way as in the previous paragraph we get $\Delta_\beta(I(G)^{n-1}) = \langle X, Y \rangle$, which yields $\text{depth } R/I(G)^{n-1} = 1$. This also contradicts to the minimality of n . Hence,

$$\sum_{i \notin X} \alpha_i = \sum_{i \notin Y} \alpha_i = n-1,$$

as required. \square

We now give an explicit solution of the equation $\Delta_\alpha(I(G)^n) = \langle X, Y \rangle$. This solution turns out to be optimal for studying $\text{dstab}(I(G))$.

Definition 3.2. Let G be a graph. We define:

- (1) For each $i \in V(G)$, denote $\mu_G(i)$ to be the number of non-leaf edges of G that are incident with i ,
- (2) $\mu(G) := (\mu_G(1), \dots, \mu_G(r)) \in \mathbb{N}^r$.

Lemma 3.3. Let G be a connected bipartite graph with bipartition (X, Y) . Let $\alpha := \mu(G)$ and $n := \varepsilon(G) - \varepsilon_0(G) + 1$. Then,

$$\Delta_\alpha(I(G)^n) = \langle X, Y \rangle, \text{ and } \sum_{i \notin X} \alpha_i = \sum_{i \notin Y} \alpha_i = \varepsilon(G) - \varepsilon_0(G).$$

Proof. Clearly, X and Y are facets of $\Delta(G)$. If $v(G) = 2$, i.e., G is exactly an edge $\{1, 2\}$, then $n = 1$ and $\alpha = (0, 0)$. We may assume that $X = \{1\}$ and $Y = \{2\}$. Then, $\Delta_\alpha(I(G)^n) = \Delta(I(G)) = \Delta(G) = \langle \{1\}, \{2\} \rangle$, so the lemma holds for this case.

Assume that $v(G) \geq 3$. Let $S := \{i \in X \mid \deg i = 1\}$ and $T := \{j \in Y \mid \deg j = 1\}$, so that

$$|S| + |T| = \varepsilon_0(G). \quad (7)$$

From [1, Theorem 1.1 and Exercise 1.1.9] we have

$$\sum_{i \in X} \deg i = \sum_{j \in Y} \deg j = \varepsilon(G). \quad (8)$$

Note that the unique neighbor of each leaf of G in X is a non-leaf of G in Y . Together with Formulas (7)–(8), this fact gives

$$\sum_{i \in X} \mu_G(i) = \sum_{i \in X} \deg i - |S| - |T| = \varepsilon(G) - \varepsilon_0(G) = n - 1.$$

Similarly,

$$\sum_{j \in Y} \mu_G(j) = \sum_{j \in Y} \deg j - |S| - |T| = \varepsilon(G) - \varepsilon_0(G) = n - 1.$$

Hence, $X, Y \in \mathcal{F}(\Delta_\alpha(I(G)^n))$ by Lemma 1.5. So in order to prove the lemma it remains to prove that $\Delta_\alpha(I(G)^n) = \langle X, Y \rangle$, or equivalently, if $F \in \mathcal{F}(\Delta(G)) \setminus \{X, Y\}$ then $F \notin \mathcal{F}(\Delta_\alpha(I(G)^n))$.

Indeed, by the maximality of F , we can partition F into $F = U \cup V$, where U and V are nonempty proper subsets of X and Y , respectively, such that every vertex in $X \setminus U$ (resp. in $Y \setminus V$) is adjacent to at least one vertex in V (resp. in U), and no vertex in U is adjacent to a vertex in V . Then, we have

$$\begin{aligned}\sum_{i \in X \setminus U} \mu_G(i) &= \sum_{i \in X \setminus U} \deg i - |S \cap (X \setminus U)| - |T \cap V|, \\ \sum_{j \in Y \setminus V} \mu_G(j) &= \sum_{j \in Y \setminus V} \deg j - |T \cap (Y \setminus V)| - |S \cap U|,\end{aligned}$$

and

$$\sum_{j \in Y \setminus V} \deg j = \sum_{i \in U} \deg i + \sum_{j \in Y \setminus V} |N_G(j) \cap (X \setminus U)|.$$

Combining these Equations with Formulas (7)–(8) we obtain

$$\begin{aligned}\sum_{u \notin F} \mu_G(u) &= \sum_{i \in X \setminus U} \mu_G(i) + \sum_{j \in Y \setminus V} \mu_G(j) \\ &= \sum_{i \in X \setminus U} \deg i - |S \cap (X \setminus U)| - |T \cap V| + \sum_{j \in Y \setminus V} \deg j - |T \cap (Y \setminus V)| - |S \cap U| \\ &= \sum_{i \in X \setminus U} \deg i + \sum_{j \in Y \setminus V} \deg j - (|S \cap U| + |S \cap (X \setminus U)| + |T \cap V| + |T \cap (Y \setminus V)|) \\ &= \sum_{i \in X \setminus U} \deg i + \sum_{j \in Y \setminus V} \deg j - (|S| + |T|) \\ &= \sum_{i \in X \setminus U} \deg i + \sum_{i \in U} \deg i + \sum_{j \in Y \setminus V} |N_G(j) \cap (X \setminus U)| - \varepsilon_0(G) \\ &= \sum_{i \in X} \deg i - \varepsilon_0(G) + \sum_{j \in Y \setminus V} |N_G(j) \cap (X \setminus U)| \\ &= \varepsilon(G) - \varepsilon_0(G) + \sum_{j \in Y \setminus V} |N_G(j) \cap (X \setminus U)|,\end{aligned}$$

or equivalently,

$$\sum_{u \notin F} \mu_G(u) = \varepsilon(G) - \varepsilon_0(G) + |P| = n - 1 + |P|,$$

where $P = \{(a, b) \mid a \in X \setminus U, b \in Y \setminus V \text{ and } ab \in E(G)\}$. Therefore, by Lemma 1.5 we have $F \notin \Delta_\alpha(I(G)^n)$ whenever $|P| \geq 1$, i.e., $P \neq \emptyset$.

In order to prove $P \neq \emptyset$, let $\ell := \min\{d_G(i, j) \mid i \in U \text{ and } j \in V\}$. Then, ℓ is finite because G is connected. Let $a \in U$ and $b \in V$ such that there is a path of length ℓ connects a and b . Suppose

$$a = a_1, b_1, a_2, b_2, \dots, a_s, b_s = b$$

is such a path, where $a_1, \dots, a_s \in X$ and $b_1, \dots, b_s \in Y$. Then, $b_1 \in Y \setminus V$ because $a_1 = a \in U$. Now if $a_2 \in U$, then we would have the path $a_2, b_2, \dots, a_s, b_s = b$ that

connects $a_2 \in U$ and $b \in V$ of length $\ell - 2$. This contradicts to the minimality of ℓ . Thus, $a_2 \in X \setminus U$. This implies $(a_2, b_1) \in P$, so $P \neq \emptyset$, as required. \square

Let G be a graph and C be a cycle of G . For any vertex v of G , we define the distance from v to C to be:

$$d_G(v, C) = \{d_G(v, u) \mid u \in V(C)\}.$$

Proposition 3.4. *Let G be a connected bipartite graph and let $2k$ be the maximum length of cycle of G ($k := 1$ if G is a tree). Then, $\text{dstab}(I(G)) \leq v(G) - \varepsilon_0(G) - k + 1$.*

Proof. Let (X, Y) be a bipartition of G .

If G is a tree, then $\varepsilon(G) = v(G) - 1$ by [1, Theorem 4.3]. Let $\alpha := \mu(G)$ and $n := \varepsilon(G) - \varepsilon_0(G) + 1$. Then, $\Delta_\alpha(I(G)^n) = \langle X, Y \rangle$ by Lemma 3.3. Hence, by Lemma 3.1, we have

$$\text{dstab}(I(G)) \leq n = \varepsilon(G) - \varepsilon_0(G) + 1 = v(G) - \varepsilon_0(G),$$

and the proposition follows.

Assume that G has a cycle, say C_{2k} , of length $2k$ where $k \geq 2$. If C is another cycle of G , then C has an edge e not lying in the cycle C_{2k} . Delete this edge from G , thereby obtaining a connected subgraph G' of G with $V(G') = V(G)$ and C_{2k} is a cycle of G' . This process continues until we obtain a connected subgraph H of G such that $V(G) = V(H)$ and H has only one cycle, that is C_{2k} . Note that H is also a bipartite graph with bipartition (X, Y) . Assume that the cycle C_{2k} is:

$$1, 2, \dots, 2k - 1, 2k, 1.$$

Let $n := v(H) - \varepsilon_0(H) - k + 1$ and define $\alpha = (\alpha_1, \dots, \alpha_r) \in \mathbb{N}^r$ by

$$\alpha_j := \begin{cases} \mu_H(j) - 1 & \text{if } 1 \leq j \leq 2k + 2, \\ \mu_H(j) & \text{otherwise.} \end{cases}$$

Claim 1.

$$\Delta_\alpha(I(H)^n) = \langle X, Y \rangle \text{ and } \sum_{i \notin X} \alpha_i = \sum_{i \notin Y} \alpha_i = n - 1.$$

Proof. We will prove this claim by induction on $v(H)$. If $v(H) = 2k$, then $H = C_{2k}$, $r = 2k$ and $n = k + 1$. We may assume also that $X = \{1, 3, \dots, 2k - 1\}$ and $Y = \{2, 4, \dots, 2k\}$. By noticing that $\alpha = (1, 1, \dots, 1) \in \mathbb{N}^r$, we have

$$\sum_{i \notin X} \alpha_i = \sum_{i \notin Y} \alpha_i = k = n - 1,$$

and therefore X and Y are facets of $\Delta_{\alpha}(I(H)^n)$. Hence, it remains to show that $\Delta_{\alpha}(I(H)^n) = \langle X, Y \rangle$. Let F be a facet of $\Delta(H)$ that is different from X and Y . Since all facets of $\Delta(C_{2k})$ have at most k elements; and only X and Y have exactly k elements, we must have $|F| < k$. Hence,

$$\sum_{i \notin F} \alpha_i \geq k + 1 = n,$$

and hence $F \notin \Delta_{\alpha}(I(H)^n)$. Therefore, $\Delta_{\alpha}(I(H)^n) = \langle X, Y \rangle$, and the claim follows.

Assume that $v(H) > 2k$. Clearly, r is not in C_{2k} , so we may assume that $d_G(r, C_{2k}) \geq d_G(v, C_{2v})$ for any vertex v of G . Then, r is a leaf by [Remark 2.1](#). Let t be the unique neighbor of r in G .

Let $T := H \setminus \{r\}$. Then, T is also a connected bipartite graph with only cycle C_{2k} and $v(T) = v(H) - 1$. We may assume that $r \in X$, so that $(X \setminus \{r\}, Y)$ is a bipartition of T . Let $s := v(T) - \varepsilon_0(T) - k + 1$ and define $\beta = (\beta_1, \dots, \beta_{r-1}) \in \mathbb{N}^{r-1}$ by

$$\beta_j := \begin{cases} \mu_T(j) - 1 & \text{if } 1 \leq j \leq 2k, \\ \mu_T(j) & \text{otherwise.} \end{cases}$$

We now distinguish two cases:

Case 1. $d_G(r, C_{2k}) = 1$. In this case $V(G) \setminus V(C_{2k})$ is the set of all leaves of G . Thus, $\beta = (\alpha_1, \dots, \alpha_{r-1})$ and $\varepsilon_0(T) = \varepsilon_0(H) - 1$, and thus $s = n$.

Since $v(T) = v(H) - 1$ and $\alpha_r = 0$, by the induction hypothesis we have $\Delta_{\beta}(I(T)^n) = \langle X \setminus \{r\}, Y \rangle$, and

$$\sum_{i \notin X} \alpha_i = \sum_{i \notin X \setminus \{r\}} \beta_i = n - 1, \text{ and } \sum_{i \notin Y} \alpha_i = \sum_{i \notin Y} \beta_i + \alpha_r = n - 1. \quad (9)$$

In particular, $X \in \Delta_{\alpha}(I(H)^n)$ and $Y \in \Delta_{\alpha}(I(H)^n)$. Thus it remains to show that $\Delta_{\alpha}(I(H)^n) = \langle X, Y \rangle$. Let F be any facet of $\Delta(H)$ that is different from X and Y .

Assume that $t \in F$. Then, F is also a facet of $\Delta(T)$ that is different from $X \setminus \{r\}$ and Y . Therefore,

$$\sum_{i \notin F} \alpha_i = \sum_{i \notin F} \beta_i + \alpha_r = \sum_{i \notin F} \beta_i \geq n.$$

Therefore, $F \notin \Delta_{\alpha}(I(H)^n)$.

Assume that $t \notin F$. Then, $r \in F$ and $F \setminus \{r\}$ is a subset of neither $X \setminus \{r\}$ nor Y . Since $F \setminus \{r\} \in \Delta(T)$, there is a facet F' of $\Delta(T)$ containing $F \setminus \{r\}$ and being different from $X \setminus \{r\}$ and Y . Therefore,

$$\sum_{i \notin F} \alpha_i = \sum_{i \notin F \setminus \{r\}} \beta_i \geq \sum_{i \notin F'} \beta_i \geq n.$$

Which implies $F \notin \Delta_{\alpha}(I(H)^n)$. The claim holds for this case.

Case 2. $d_G(r, C_{2k}) \geq 2$. By Remark 2.1 we can assume that $N_G(t) = \{t-1, t+1, \dots, r\}$ where $t-1$ is a non-leaf and $t+1, \dots, r$ are leaves. We now distinguish two subcases:

Case 2a. $t+1 = r$. Then, $\varepsilon_0(T) = \varepsilon_0(H)$ and $s = n-1$. Since $v(T) = v(H) - 1$, $\alpha_r = 0$ and

$$\beta_j = \begin{cases} \alpha_j - 1 & \text{if } j = t-1 \text{ or } j = t, \\ \alpha_j & \text{otherwise,} \end{cases}$$

by the induction hypothesis we have $\Delta_\beta(I(T)^{n-1}) = \langle X \setminus \{r\}, Y \rangle$, and

$$\sum_{i \notin X} \alpha_i = \sum_{i \notin X \setminus \{r\}} \beta_i + 1 = n-1, \text{ and } \sum_{i \notin Y} \alpha_i = \sum_{i \notin Y} \beta_i + \alpha_r + 1 = n-1. \quad (10)$$

In particular, $X \in \Delta_\alpha(I(H)^n)$ and $Y \in \Delta_\alpha(I(H)^n)$. Thus it remains to show that $\Delta_\alpha(I(H)^n) = \langle X, Y \rangle$. Let F be any facet of $\Delta(H)$ that is different from X and Y .

Assume that $t \in F$. Then, F is also a facet of $\Delta(T)$ that is different from $X \setminus \{r\}$ and Y . Since $t-1 \notin F$ and $\alpha_{t-1} = \beta_{t-1} + 1$, we have

$$\sum_{i \notin F} \alpha_i = \sum_{i \notin F} \beta_i + 1 + \alpha_r = \sum_{i \notin F} \beta_i + 1 \geq s+1 = n.$$

Therefore, $F \notin \Delta_\alpha(I(H)^n)$.

Assume that $t \notin F$. Then, $r \in F$. If $t-1 \in F$, then $F \setminus \{r\}$ is a subset of neither $X \setminus \{r\}$ nor Y . Hence, there is a facet F' of $\Delta(T)$ containing F and being different from $X \setminus \{r\}$ and Y . Therefore,

$$\sum_{i \notin F} \alpha_i \geq \sum_{i \notin F} \beta_i + 1 \geq \sum_{i \notin F'} \beta_i + 1 \geq s+1 = n.$$

Which implies $F \notin \Delta_\alpha(I(H)^n)$.

If $t-1 \notin F$, then $(F \cup \{t\}) \setminus \{r\}$ is a facet of $\Delta(T)$. Noticing that $\alpha_{t-1} = \beta_{t-1} + 1$ and $\alpha_t = 1$, we get

$$\sum_{i \notin F} \alpha_i = \sum_{i \notin (F \cup \{t\}) \setminus \{r\}} \beta_i + 1 + \alpha_t \geq (s-1) + 2 = n.$$

Which again implies $F \notin \Delta_\alpha(I(H)^n)$.

Case 2. $t+1 < r$. Thus $\beta = (\alpha_1, \dots, \alpha_{r-1})$, and thus $s = n$. Now we can proceed as in Case 1. This completes the proof of Claim 1.

Claim 2. $\Delta_\alpha(I(G)^n) = \langle X, Y \rangle$.

Proof. by Claim 1 and Lemma 1.5, X and Y are facets of $\Delta_\alpha(I(G)^n)$. It remains to show that for any facet F of $\Delta(G)$ being different from X and Y , then $F \notin \Delta_\alpha(I(G)^n)$. Since F is a face of H , we have $F \subseteq F'$ for some facet F' of $\Delta(H)$. Then, F' is different from X and Y , and then $F' \notin \Delta_\alpha(I(H)^n)$. Thus, by Lemma 1.5 we have

$$\sum_{i \notin F} \alpha_i \geq \sum_{i \notin F'} \alpha_i \geq n$$

and thus $F \notin \Delta_\alpha(I(G)^n)$, as claimed.

Now we return to the proof of the proposition. Claim 2 and Lemma 3.1 give $\text{dstab}(I(G)) \leq n$, or equivalently

$$\text{dstab}(I(G)) \leq \varepsilon(H) - \varepsilon_0(H) - k + 1.$$

Let e be an edge of the cycle C_{2k} . Then, $H \setminus e$ is a tree. Hence, by [1, Theorem 4.3] we have $\varepsilon(H) = \varepsilon(H \setminus e) + 1 = (v(H \setminus e) - 1) + 1 = v(H \setminus e) = v(H) = v(G)$. Clearly, $\varepsilon_0(G) \leq \varepsilon_0(H)$. Therefore,

$$\text{dstab}(I(G)) \leq \varepsilon(H) - \varepsilon_0(H) - k + 1 \leq v(G) - \varepsilon_0(G) - k + 1,$$

as required. \square

4. Depths of powers of edge ideals

In this section we study the stability of $\text{depth } R/I(G)^n$ for any graph G . First we need some basic facts of homological modules of simplicial complexes.

A tool which will be of much use is the Mayer–Vietoris sequence, see [21, Theorem 25.1] or [23, in Page 21]. For two simplicial complexes Δ_1 and Δ_2 , we have the long exact sequence of reduced homology modules

$$\cdots \rightarrow \tilde{H}_i(\Delta_1) \oplus \tilde{H}_i(\Delta_2) \rightarrow \tilde{H}_i(\Delta) \rightarrow \tilde{H}_{i-1}(\Delta_1 \cap \Delta_2) \rightarrow \tilde{H}_{i-1}(\Delta_1) \oplus \tilde{H}_{i-1}(\Delta_2) \rightarrow \cdots$$

where $\Delta = \Delta_1 \cup \Delta_2$.

A simplicial complex Δ is a cone if there is a vertex v such that $\{v\} \cup F \in \Delta$ for every $F \in \Delta$. If Δ is a cone, then it is acyclic (see [21, Theorem 8.2]), i.e.,

$$\tilde{H}_i(\Delta; K) = \mathbf{0} \text{ for every } i \in \mathbb{Z}.$$

Finally, for two simplicial complexes Δ and Γ over two disjoint vertex sets, the join of Δ and Γ , denoted by $\Delta * \Gamma$, is defined by

$$\Delta * \Gamma := \{F \cup G \mid F \in \Delta \text{ and } G \in \Gamma\}.$$

Lemma 4.1. *Let G be a bipartite graph with connected components G_1, \dots, G_s and let $n := \sum_{i=1}^s \text{dstab}(I(G_i)) - s + 1$. Then there is $\alpha = (\alpha_1, \dots, \alpha_r) \in \mathbb{N}^r$ such that*

$$\sum_{i \notin F} \alpha_i = n - 1 \quad \text{for all } F \in \mathcal{F}(\Delta_\alpha(I(G)^n)) \quad \text{and} \quad \tilde{H}_{s-1}(\Delta_\alpha(I(G)^n); K) \neq \mathbf{0}.$$

Proof. For each i , let (X_i, Y_i) be a bipartition of G_i and $n_i := \text{dstab}(I(G_i))$, so that

$$n = \sum_{i=1}^s n_i - s + 1.$$

Since the vertex sets of G_1, \dots, G_s are mutually disjoint, by Lemma 3.1 there is $\alpha = (\alpha_1, \dots, \alpha_r) \in \mathbb{N}^r$ such that

$$\sum_{j \in V(G_i) \setminus X_i} \alpha_j = \sum_{j \in V(G_i) \setminus Y_i} \alpha_j = n_i - 1, \quad (11)$$

and

$$\sum_{j \in V(G_i) \setminus F_i} \alpha_j \geq n_i \quad \text{for all } F_i \in \mathcal{F}(\Delta(G_i)) \setminus \{X_i, Y_i\}. \quad (12)$$

For any $F \in \mathcal{F}(\Delta(G))$, we can partition F into $F = \bigcup_{i=1}^s F_i$ where $F_i \in \mathcal{F}(\Delta(G_i))$ for $i = 1, \dots, s$. By Equation (11) and Inequality (12) we get

$$\sum_{j \notin F} \alpha_j = \sum_{i=1}^s \sum_{j \in V(G_i) \setminus F_i} \alpha_j \geq \sum_{i=1}^s (n_i - 1) = n - 1$$

and the equality occurs if and only if

$$\sum_{j \in V(G_i) \setminus F_i} \alpha_j = n_i - 1 \quad \text{for all } i = 1, \dots, s,$$

or equivalently, either $F_i = X_i$ or $F_i = Y_i$ for all $i = 1, \dots, s$. Together with Lemma 1.5 we have

$$\sum_{j \notin F} \alpha_j = n - 1 \quad \text{for all } F \in \mathcal{F}(\Delta_\alpha(I(G)^n)),$$

and

$$\Delta_\alpha(I(G)^n) = \langle X_1, Y_1 \rangle * \dots * \langle X_s, Y_s \rangle.$$

So it remains to prove that $\tilde{H}_{s-1}(\langle X_1, Y_1 \rangle * \dots * \langle X_s, Y_s \rangle; K) \neq \mathbf{0}$. In order to prove this, let $\Delta_i := \langle X_1, Y_1 \rangle * \dots * \langle X_i, Y_i \rangle$ for $i = 1, \dots, s$ and $\Delta_0 := \{\emptyset\}$. Then, for all $i = 1, \dots, s$

we have

$$\Delta_i = \langle X_i \rangle * \Delta_{i-1} \cup \langle Y_i \rangle * \Delta_{i-1} \quad \text{and} \quad \Delta_{i-1} = \langle X_i \rangle * \Delta_{i-1} \cap \langle Y_i \rangle * \Delta_{i-1}.$$

Since $\langle X_i \rangle * \Delta_{i-1}$ and $\langle Y_i \rangle * \Delta_{i-1}$ are cones, by using Mayer–Vietoris sequence, we get an exact sequence $\mathbf{0} \rightarrow \tilde{H}_{s-1}(\Delta_s; K) \rightarrow \tilde{H}_{s-2}(\Delta_{s-1}; K) \rightarrow \mathbf{0}$. Thus,

$$\tilde{H}_{s-1}(\Delta_s; K) \cong \tilde{H}_{s-2}(\Delta_{s-1}; K).$$

By repeating this way we obtain

$$\tilde{H}_{s-1}(\Delta_s; K) \cong \tilde{H}_{s-2}(\Delta_{s-1}; K) \cong \cdots \cong \tilde{H}_1(\Delta_0; K) \cong K,$$

and so $\tilde{H}_{s-1}(\Delta_s; K) \neq \mathbf{0}$, as required. \square

The next lemma gives the limit of the sequence $\text{depth } R/I(G)^n$.

Lemma 4.2. *Let G be a graph. Assume that G_1, \dots, G_s are all connected bipartite components of G and G_{s+1}, \dots, G_{s+t} are all connected nonbipartite components of G . Then*

$$\text{depth } R/I(G)^n = s \quad \text{for all } n \geq \sum_{i=1}^{s+t} \text{dstab}(I(G_i)) - (s+t) + 1.$$

Proof. Let $n_i := \text{dstab}(I(G_i))$ for $i = 1, \dots, s+t$. We divide the proof into three cases:

Case 1. $s = 0$, i.e., every component of G is nonbipartite. This case follows from [Lemma 1.1](#).

Case 2. $t = 0$, i.e., G is bipartite. Let $m := \sum_{i=1}^s n_i - s + 1$. By [Lemmas 1.4 and 4.1](#), there is $\alpha \in \mathbb{N}^r$ such that

$$\dim_K H_m^s(R/I(G)^m)_\alpha = \dim_K \tilde{H}_{s-1}(\Delta_\alpha(I(G)^m); K) \neq 0.$$

Hence, $H_m^s(R/I(G)^m) \neq \mathbf{0}$, which yields $\text{depth } R/I(G)^m \leq s$. On the other hand, by [Lemma 1.3](#) we have $\text{depth } R/I(G)^m \geq s$. Thus, $\text{depth } R/I(G)^m = s$. The lemma now follows from [Lemma 1.3](#).

Case 3. $s \neq 0$ and $t \neq 0$. Let G' and G'' be induced subgraphs of G defined by

$$G' := \bigcup_{i=1}^s G_i \quad \text{and} \quad G'' := \bigcup_{i=1}^t G_{s+i}.$$

We may assume that $V(G') = [p]$ and $V(G'') = \{p+1, \dots, p+q\}$, where $p+q = r$. For simplicity, we set $y_1 := x_{p+1}, \dots, y_q := x_{p+q}$. Then $R = K[x_1, \dots, x_p, y_1, \dots, y_q]$. Let

$R' := K[x_1, \dots, x_p]$, $R'' := K[y_1, \dots, y_q]$, $m := \sum_{i=1}^s n_i - s + 1$ and $n_0 := n - m + 1$. Note that $n_0 \geq \sum_{i=1}^t n_{s+i} - t + 1$, so $(y_1, \dots, y_q) \in \text{Ass}(R''/I(G'')^{n_0})$ by Lemma 1.1. Accordingly, there exists $\beta = (\beta_1, \dots, \beta_q) \in \mathbb{N}^q$ such that $(y_1, \dots, y_q) = I(G'')^{n_0} : \mathbf{y}^\beta$. This implies

$$\mathbf{y}^\beta \in I(G'')^{n_0-1}, \mathbf{y}^\beta \notin I(G'')^{n_0} \text{ and } \mathbf{y}^\beta \in I(G'')_{F'}^{n_0} \text{ whenever } \emptyset \neq F \in \Delta(G''). \quad (13)$$

Next, by Lemma 4.1 there is $\alpha = (\alpha_1, \dots, \alpha_p) \in \mathbb{N}^p$ such that

$$\tilde{H}_{s-1}(\Delta_\alpha(I(G')^m); K) \neq \mathbf{0}, \text{ and } \sum_{i \notin V} \alpha_i = m - 1 \text{ for all } V \in \mathcal{F}(\Delta_\alpha(I(G')^m)). \quad (14)$$

Let $\gamma := (\alpha_1, \dots, \alpha_p, \beta_1, \dots, \beta_q) \in \mathbb{N}^r$. Note that $\mathbf{x}^\gamma = \mathbf{x}^\alpha \mathbf{y}^\beta \in R$. We claim that

$$\Delta_\gamma(I(G)^n) = \Delta_\alpha(I(G')^m). \quad (15)$$

Indeed, for all $H \in \Delta_\gamma(I(G)^n)$ we can partition H into $H = H_1 \cup H_2$ where $H_1 \in \Delta(G')$ and $H_2 \in \Delta(G'')$. By Equation (1) we have

$$\mathbf{x}^\gamma = \mathbf{x}^\alpha \mathbf{y}^\beta \notin I(G)_H^n = (I(G')_{H_1} + I(G'')_{H_2})^n = \sum_{i=0}^n I(G')_{H_1}^i I(G'')_{H_2}^{n-i}. \quad (16)$$

Now, if $H_2 \neq \emptyset$, then by Formula (13) we would have $\mathbf{y}^\beta \in I(G'')_{H_2}^{n_0}$. Then, Formula (16) forces $\mathbf{x}^\alpha \notin I(G')_{H_1}^{n-n_0} = I(G')_{H_1}^{m-1}$, thus $H_1 \in \Delta_\alpha(I(G')^{m-1})$. In particular, $\Delta_\alpha(I(G')^{m-1}) \neq \emptyset$. Let us take arbitrary facet V of $\Delta_\alpha(I(G')^{m-1})$. By Lemma 1.5 we then have $\sum_{i \notin V} \alpha_i \leq m - 2$. By Lemma 1.5 again, V is a facet of $\Delta_\alpha(I(G')^m)$, which contradicts (14). Thus, $H_2 = \emptyset$ and $H = H_1$. Formula (16) now becomes

$$\mathbf{x}^\gamma = \mathbf{x}^\alpha \mathbf{y}^\beta \notin (I(G')_H + I(G''))^n = \sum_{i=0}^n I(G')_H^i I(G'')^{n-i}.$$

Together with Formula (13), this fact implies $\mathbf{x}^\alpha \notin I(G')_H^{n-n_0+1} = I(G')_H^m$, or equivalently, $H \in \Delta_\alpha(I(G')^m)$, so $\Delta_\gamma(I(G)^n) \subseteq \Delta_\alpha(I(G')^m)$.

In order to prove the reverse inclusion, suppose that $H \in \Delta_\alpha(I(G')^m)$. Then, $\mathbf{x}^\alpha \notin I(G')_H^m$ by Equation (1). If $\mathbf{x}^\gamma \in I(G)_H^n$, then

$$\mathbf{x}^\gamma = \mathbf{x}^\alpha \mathbf{y}^\beta \in I(G)_H^n = (I(G')_H + I(G''))^n = \sum_{i=0}^n I(G')_H^i I(G'')^{n-i}.$$

Hence, $\mathbf{x}^\alpha \mathbf{y}^\beta \in I(G')_H^\nu I(G'')^{n-\nu}$ for some nonnegative integer ν . Since $V(G') \cap V(G'') = \emptyset$, it yields $\mathbf{x}^\alpha \in I(G')_H^\nu$ and $\mathbf{y}^\beta \in I(G'')^{n-\nu}$. By Formula (13) we deduce that $n - \nu \leq n_0 - 1$, and so $\nu \geq n - n_0 + 1 = m$. But then $\mathbf{x}^\alpha \in I(G')_H^m$, a contradiction.

Hence, $\mathbf{x}^\gamma \notin I(G)_H^n$, i.e., $H \in \Delta_\gamma(I(G)^n)$, and hence $\Delta_\alpha(I(G')^m) \subseteq \Delta_\gamma(I(G)^n)$, as claimed.

Combining Formulas (14) and (15) with Lemma 1.4, we get

$$\begin{aligned} \dim_K H_{\mathfrak{m}}^s(R/I(G)^n)_\gamma &= \dim_K \tilde{H}_{s-1}(\Delta_\gamma(I(G)^n); K) \\ &= \dim_K \tilde{H}_{s-1}(\Delta_\alpha(I(G')^m); K) \neq 0. \end{aligned}$$

Therefore, $H_{\mathfrak{m}}^s(R/I(G)^n) \neq \mathbf{0}$, so

$$\text{depth } R/I(G)^n \leq s. \quad (17)$$

On the other hand, since G' is bipartite, by Lemmas 1.3 and 1.6 we get

$$\text{depth } R/I(G)^n = \text{depth } R/(I(G') + I(G''))^n \geq \min_{\nu \geq 1} \text{depth } R'/I(G')^\nu = s.$$

Together with Inequality (17), we obtain $\text{depth } R/I(G)^n = s$, as required. \square

Corollary 4.3. *For all graphs G we have $\lim_{n \rightarrow \infty} \text{depth } R/I(G)^n = \dim R - \ell(I(G))$.*

Proof. Let s be the number of bipartite components of G . Then $s = \dim R - \ell(I(G))$ (see [25, Page 50]), so the corollary immediately follows from Lemma 4.2. \square

We are now ready to prove the first main result of the paper.

Theorem 4.4. *Let G be a graph with p connected components G_1, \dots, G_p . Let s be the number of connected bipartite components of G . Then*

- (1) $\min\{\text{depth } R/I(G)^n \mid n \geq 1\} = s$.
- (2) $\text{dstab}(I(G)) = \min\{n \geq 1 \mid \text{depth } R/I(G)^n = s\}$.
- (3) $\text{dstab}(I(G)) = \sum_{i=1}^p \text{dstab}(I(G_i)) - p + 1$.

Proof. We may assume that G_1, \dots, G_s are bipartite.

(1) If $s = 0$ (resp. $s = p$), then the first statement follows from Lemma 1.1 (resp. Lemma 1.3). Assume that $1 \leq s < p$. Let G' be the induced subgraph of G consisting of G_1, \dots, G_s and G'' the induced subgraph of G consisting of G_{s+1}, \dots, G_p . Then, $I(G) = I(G') + I(G'')$. Let $R' := K[x_i \mid i \in V(G')]$. For all $n \geq 1$, since G' is bipartite, by Lemmas 1.3 and 1.6 we have

$$\text{depth } R/I(G)^n \geq \min\{\text{depth } R'/I(G')^m \mid m \geq 1\} = s.$$

Together with Lemma 4.2 we conclude that

$$\min\{\text{depth } R/I(G)^n \mid n \geq 1\} = s,$$

and (1) follows.

We next prove (2) and (3) simultaneously by induction on p . If $p = 1$, then the theorem follows from [Lemmas 1.1 and 1.3](#).

Assume that $p \geq 2$. If $s = 0$, our claim follows from [Lemma 1.1](#). So we may assume that $s \geq 1$. Let H be the induced subgraph of G consisting of components G_2, \dots, G_p . Then, H has $p - 1$ connected components and $s - 1$ connected bipartite components. By [Lemma 4.2](#) we have

$$\text{depth } R/I(G)^n = s \text{ for all } n \geq \sum_{i=1}^p \text{dstab}(I(G_i)) - p + 1.$$

Hence, in order to prove the theorem it suffices to show that if

$$\text{depth } R/I(G)^n = s \tag{18}$$

for a given positive integer n , then $n \geq \sum_{i=1}^p \text{dstab}(I(G_i)) - p + 1$.

In order to prove this assertion let $A := K[x_j \mid j \in V(G_1)]$ and $B := K[x_j \mid j \in V(H)]$. Then, we have $\dim A \geq 2$ and $\dim B \geq s$. For simplicity, we set $I := I(G_1)$ and $J := I(H)$. We now claim that

$$\text{depth } R/I^i J^{n-i} \geq s + 1 \text{ for } i = 0, \dots, n. \tag{19}$$

Indeed, if $i = n$, since $\text{depth } A/I^n \geq 1$ and $\dim B \geq s$, we have

$$\text{depth } R/I^n J^0 = \text{depth } R/I^n = \text{depth } A/I^n + \dim B \geq 1 + s.$$

Since $\text{depth } B/J^n \geq s - 1$ by Part 1, a similar proof also holds for $i = 0$. For all $i = 1, \dots, n - 1$, by [\[11, Lemma 2.2\]](#) we have $\text{depth } R/I^i J^{n-i} = \text{depth } A/I^i + \text{depth } B/J^{n-i} + 1$. Hence, $\text{depth } R/I^i J^{n-i} \geq 1 + (s - 1) + 1 = s + 1$, as claimed.

Let $n_1 := \text{dstab}(G_1)$ and $n_2 := \text{dstab}(H)$. We will prove that $n \geq n_1 + n_2 - 1$. Assume on the contrary that $n \leq n_1 + n_2 - 2$. For each $i = 0, \dots, n$, we put

$$W_i := I^i J^{n-i} + \dots + I^n J^0,$$

where $I^0 = J^0 = R$. We next claim that

$$\text{depth } R/W_i \geq s + 1 \text{ for all } i = 0, \dots, n. \tag{20}$$

Indeed, we prove this by induction on i . If $i = n$, then by Inequality (19) we have

$$\text{depth } R/W_n = \text{depth } R/I^n \geq s + 1.$$

Assume that $\text{depth } R/W_{i+1} \geq s + 1$ for some $0 \leq i < n$. By Equations (2) and (3), we have $I^i J^{n-i} \cap W_{i+1} = I^{i+1} J^{n-i}$. Since $W_i = I^i J^{n-i} + W_{i+1}$, we have an exact sequence

$$0 \longrightarrow R/I^{i+1} J^{n-i} \longrightarrow R/I^i J^{n-i} \oplus R/W_{i+1} \longrightarrow R/W_i \longrightarrow 0.$$

By Depth Lemma, we have

$$\text{depth } R/W_i \geq \min\{\text{depth } R/I^{i+1}J^{n-i} - 1, \text{depth } R/I^iJ^{n-i}, \text{depth } R/W_{i+1}\}.$$

Together with Inequality (19) and the induction hypothesis, this fact yields

$$\text{depth } R/W_i \geq \min\{\text{depth } R/I^{i+1}J^{n-i} - 1, s + 1\}.$$

Therefore, the inequality (20) will follow if $\text{depth } R/I^{i+1}J^{n-i} \geq s + 2$. In order to prove this inequality, note that $(i + 1) + (n - i) = n + 1 \leq n_1 + n_2 - 1$. Hence, either $i + 1 < n_1$ or $n - i < n_2$. Note that $n - i \geq 1$.

If $i + 1 < n_1$, by Part 1 we get $\text{depth } A/I^{i+1} \geq 2$ and $\text{depth } B/J^{n-i} \geq s - 1$. Together with [11, Lemma 2.2] we obtain

$$\text{depth } R/I^{i+1}J^{n-i} = \text{depth } A/I^{i+1} + \text{depth } B/J^{n-i} + 1 \geq 2 + (s - 1) + 1 = s + 2,$$

as claimed.

If $n - i < n_2$, the proof is similar. Thus, the claim (20) is proved.

Notice that $W_0 = (I + J)^n = (I(G_1) + I(H))^n = I(G)^n$. By (20) we have $\text{depth } R/I(G)^n \geq s + 1$. This contradicts (18). Therefore, we must have $n \geq n_1 + n_2 - 1$.

Finally, by the induction hypothesis we have

$$n_2 = \text{dstab}(I(H)) = \sum_{i=2}^p \text{dstab}(I(G_i)) - (p - 1) + 1.$$

Together with $n_1 = \text{dstab}(I(G_1))$, we have

$$n \geq n_1 + n_2 - 1 = \sum_{i=1}^p \text{dstab}(G_i) - p + 1,$$

as required. \square

Remark 4.5. From Theorem 4.4 and Lemmas 1.1 and 3.1 we see that $\text{dstab}(I(G))$ is independent from the characteristic of the base field K , so it depends purely on the structure of G .

We next combine Theorem 4.4 and Propositions 2.4 and 3.4 to get the second main result of the paper, which sets up an upper bound for $\text{dstab}(I(G))$.

Theorem 4.6. Let G be a graph. Let G_1, \dots, G_s be all connected bipartite components of G and let G_{s+1}, \dots, G_{s+t} be all connected nonbipartite components of G . Let $2k_i$ be the

maximum length of cycles of G_i ($k_i := 1$ if G_i is a tree) for all $i = 1, \dots, s$; and let $2k_i - 1$ be the maximum length of odd cycles of G_i for every $i = s + 1, \dots, s + t$. Then

$$\text{dstab}(I(G)) \leq v(G) - \varepsilon_0(G) - \sum_{i=1}^{s+t} k_i + 1.$$

Proof. Since

$$v(G) - \varepsilon_0(G) - \sum_{i=1}^{s+t} k_i + 1 = \sum_{i=1}^{s+t} (v(G_i) - \varepsilon_0(G_i) - k_i + 1) - (s + t) + 1,$$

by [Propositions 2.4 and 3.4](#) we get

$$v(G) - \varepsilon_0(G) - \sum_{i=1}^{s+t} k_i + 1 \geq \sum_{i=1}^{s+t} \text{dstab}(I(G_i)) - (s + t) + 1.$$

Together with [Theorem 4.4](#) we obtain

$$\text{dstab}(I(G)) = \sum_{i=1}^{s+t} \text{dstab}(I(G_i)) - (s + t) + 1 \leq v(G) - \varepsilon_0(G) - \sum_{i=1}^{s+t} k_i + 1,$$

as required. \square

5. The index of depth stability of trees and unicyclic graphs

The aim of this section is to prove that the upper bound of $\text{dstab}(I(G))$ given in [Theorem 4.6](#) is always achieved if G has no cycles of length 4 and every component of G is either a tree or a unicyclic graph. Recall that a connected graph G is a tree if it contains no cycles; and G is a unicyclic graph if it contains exactly one cycle.

If G is a unicyclic graph and C is the unique cycle of G , then for every vertex v of G not lying in C , there is a unique simple path of minimal distance from v to a vertex in C .

Theorem 5.1. *Let G be a graph with p connected components G_1, \dots, G_p such that each G_i is either a tree or a unicyclic graph. For each i , if G_i is bipartite, let $2k_i$ be the length of its unique cycle ($k_i := 1$ if G_i is a tree); and if G_i is nonbipartite, let $2k_i - 1$ be the length of its unique cycle. If G has no cycles of length 4, then*

$$\text{dstab}(I(G)) = v(G) - \varepsilon_0(G) - \sum_{i=1}^p k_i + 1.$$

By Theorem 4.6, it suffices to show that $\text{dstab}(G_i) = v(G_i) - \varepsilon_0(G_i) - k_i + 1$ for each $i = 1, \dots, p$. If G_i is nonbipartite, the equality follows from Lemma 2.2. Thus, it remains to prove this equality for the case G_i is bipartite.

We divide the proof into two lemmas. The first lemma deals with unicyclic bipartite graphs and the second one deals with trees.

For a vertex x of G , we denote $L_G(x)$ to be the set of leaves of G that are adjacent to x . We start with the following observation.

Lemma 5.2. *Let G be a graph with $r = v(G)$. Let p be a leaf of G and q the unique neighbor of p in G . Let $\alpha = (\alpha_1, \dots, \alpha_r) \in \mathbb{N}^r$ and we define $\beta = (\beta_1, \dots, \beta_r)$ by*

$$\beta_i := \begin{cases} \alpha_i + 1 & \text{if } i = p \text{ or } i = q, \\ \alpha_i & \text{otherwise.} \end{cases}$$

Then $\Delta_\alpha(I(G)^n) = \Delta_\beta(I(G)^{n+1})$ for all $n \geq 1$.

Proof. Let F be a facet of $\Delta(G)$. By the maximality of F , it must contain either p or q but not both, so

$$\sum_{i \notin F} \beta_i = \sum_{i \notin F} \alpha_i + 1.$$

Thus, by Lemma 1.5 we get $\Delta_\alpha(I(G)^n) = \Delta_\beta(I(G)^{n+1})$ for all $n \geq 1$. \square

Lemma 5.3. *Let G be a unicyclic bipartite graph. Assume that the unique cycle of G is C_{2k} of length $2k$ with $k \geq 3$. Then, $\text{dstab}(I(G)) = v(G) - \varepsilon_0(G) - k + 1$.*

Proof. Let $n := \text{dstab}(I(G))$. By Theorem 4.6 we have $n \leq v(G) - \varepsilon_0(G) - k + 1$. Thus, in order to prove the theorem it suffices to show $n \geq v(G) - \varepsilon_0(G) - k + 1$.

Let (X, Y) be a bipartition of G . Then, by Lemma 3.1 there is $\alpha = (\alpha_1, \dots, \alpha_r) \in \mathbb{N}^r$ such that

$$\Delta_\alpha(I(G)^n) = \langle X, Y \rangle \text{ and } \sum_{j \in X} \alpha_j = \sum_{j \in Y} \alpha_j = n - 1. \quad (21)$$

Observe that for any face F of $\Delta(G)$ with $F \cap X \neq \emptyset$ and $F \cap Y \neq \emptyset$, we have

$$\sum_{i \notin F} \alpha_i \geq n. \quad (22)$$

Indeed, let L be a facet of $\Delta(G)$ which contains F , so that L meets both X and Y . Since $\Delta_\alpha(I(G)^n) = \langle X, Y \rangle$, $L \notin \Delta_\alpha(I(G)^n)$. By Lemma 1.5 we get

$$\sum_{i \notin F} \alpha_i \geq \sum_{i \notin L} \alpha_i \geq n,$$

and the formula (22) follows.

We now prove $n \geq v(G) - \varepsilon_0(G) - k + 1$ by induction on $v(G)$.

If $v(G) = 2k$, i.e., $G = C_{2k}$, then $v(G) - \varepsilon_0(G) - k + 1 = k + 1$. For each $i \in X$, let $N_G(i) = \{u_i, v_i\}$ and $F_i := \{i\} \cup (Y \setminus \{u_i, v_i\})$. Then, $F_i \in \Delta(G)$. Since $|X| = |Y| = k \geq 3$, $F_i \cap X \neq \emptyset$ and $F_i \cap Y \neq \emptyset$. Together with Formulas (21) and (22), this fact gives

$$n \leq \sum_{j \notin F_i} \alpha_j = \sum_{j \in X} \alpha_j + \alpha_{u_i} + \alpha_{v_i} - \alpha_i = n - 1 + \alpha_{u_i} + \alpha_{v_i} - \alpha_i,$$

whence $\alpha_i + 1 \leq \alpha_{u_i} + \alpha_{v_i}$. Hence,

$$\sum_{i \in X} \alpha_i + k = \sum_{i \in X} (\alpha_i + 1) \leq \sum_{i \in X} (\alpha_{u_i} + \alpha_{v_i}) = 2 \sum_{j \in Y} \alpha_j.$$

Together with Formula (21), this gives $(n - 1) + k \leq 2(n - 1)$. Thus, $n \geq k + 1$, and thus the lemma holds for this case.

Assume that $v(G) > 2k$. We distinguish two cases:

Case 1. $G \setminus V(C_{2k})$ is totally disconnected. For any vertex u lying in C_{2k} with $L_G(u) \neq \emptyset$, we claim that

$$\alpha_u \geq 1, \text{ and } \alpha_i = 0 \text{ for every } i \in L_G(u). \quad (23)$$

Indeed, without loss of generality we may assume that $u \in Y$, so that $L_G(u) \subseteq X$. Let $F := (Y \setminus \{u\}) \cup L_G(u)$. Then, $F \in \Delta(G)$. Since the length of C_{2k} is at least 6, we have $F \cap Y \neq \emptyset$. Notice that $\emptyset \neq L_G(u) \subseteq F \cap X$. Therefore, $F \cap X \neq \emptyset$ and $F \cap Y \neq \emptyset$. By Formula (22) we have

$$\sum_{i \in X} \alpha_i + \alpha_u - \sum_{i \in L_G(u)} \alpha_i = \sum_{i \notin F} \alpha_i \geq n.$$

By (21), this gives

$$n - 1 + \alpha_u - \sum_{i \in L_G(u)} \alpha_i \geq n,$$

so

$$\alpha_u \geq \sum_{i \in L_G(u)} \alpha_i + 1 \geq 1.$$

Hence, it remains to prove that $\alpha_i = 0$ for all $i \in L_G(u)$. Assume that $\alpha_i \geq 1$ for some $i \in L_G(u)$. Define $\beta = (\beta_1, \dots, \beta_r)$ by

$$\beta_j := \begin{cases} \alpha_j - 1 & \text{if } j = u \text{ or } j = i, \\ \alpha_j & \text{otherwise.} \end{cases}$$

Then, $\beta \in \mathbb{N}^r$. Since $u \in Y$ and $\alpha_u \geq 1$, by (21) we have

$$n - 1 = \sum_{j \in Y} \alpha_j \geq \alpha_u \geq 1.$$

By Lemma 5.2 we have $\Delta_\beta(I(G)^{n-1}) = \Delta_\alpha(I(G)^n)$. Consequently, $\Delta_\beta(I(G)^{n-1}) = \langle X, Y \rangle$, which implies $\text{depth } R/I(G)^{n-1} = 1$ by Lemma 3.1, and so $\text{dstab}(I(G)) \leq n - 1$ by Theorem 4.4. This contradicts to $n = \text{dstab}(I(G))$. Thus, $\alpha_i = 0$, as claimed.

We may assume that $V(H) = \{1, \dots, 2k\}$. Let $\beta := (\alpha_1, \dots, \alpha_{2k}) \in \mathbb{N}^{2k}$, $X_0 := X \cap V(C_{2k})$ and $Y_0 := Y \cap V(C_{2k})$. Then, (X_0, Y_0) is a bipartition of C_{2k} . Clearly,

$$X = X_0 \cup \bigcup_{i \in Y_0} L_G(i) \text{ and } Y = Y_0 \cup \bigcup_{i \in X_0} L_G(i).$$

Together with Claim (23) we have

$$\sum_{i \notin X_0} \beta_i = \sum_{i \notin X} \alpha_i = n - 1.$$

Similarly, $\sum_{i \notin Y_0} \beta_i = n - 1$. Therefore, $X_0, Y_0 \in \Delta_\beta(I(C_{2k})^n)$.

For any facet F of $\Delta(C_{2k})$ which is different from X_0 and Y_0 , let

$$F' := F \cup \bigcup_{i \in V(C) \setminus F} L_G(i).$$

Then, F' is a facet of $\Delta(G)$ which is different from X and Y . Together Claim (23) with Lemma 1.5, we have

$$\sum_{i \notin F} \beta_i = \sum_{i \notin F'} \alpha_i \geq n$$

so that $F \notin \Delta_\beta(I(C_{2k})^n)$. Thus, $\Delta_\beta(I(C_{2k})^n) = \langle X_0, Y_0 \rangle$.

This gives $\text{depth } S/I(C_{2k})^n = 1$ where $S = K[x_1, \dots, x_{2k}]$. From the case $v(G) = 2k$ above, we imply that

$$n \geq k + 1 = v(G) - \varepsilon_0(G) - k + 1,$$

and the lemma holds in this case.

Case 2. $G \setminus V(C_{2k})$ is not totally disconnected. Let v be a leaf of G such that $d_G(v, C_{2k})$ is maximal. By Remark 2.1, we deduce that $N_G(v)$ has only one non-leaf, say u , and $N_G(u)$ also has only one non-leaf, say w . Note that $L_G(u) \neq \emptyset$ since $v \in L_G(u)$. We may assume that $u \in Y$, so that $v \in X$. We first claim that

$$\alpha_u \geq 1, \text{ and } \alpha_i = 0 \text{ for every } i \in L_G(u). \quad (24)$$

Indeed, let $F := (Y \setminus \{u\}) \cup L_G(u)$. Then, $F \in \Delta(G)$. Since $|N_G(w)| \geq 2$ and $N_G(w) \subseteq Y$, we have $\emptyset \neq N_G(w) \setminus \{u\} \subseteq Y \setminus \{u\} \subseteq F \cap Y$. Notice that $\emptyset \neq L_G(u) \subseteq F \cap X$. Therefore, $F \cap X \neq \emptyset$ and $F \cap Y \neq \emptyset$. The proof of claim now carries out the same as in Claim (23).

We next claim that

$$\alpha_w \geq 1. \quad (25)$$

Indeed, assume on the contrary that $\alpha_w = 0$. Note that $w \in X$ and $N_G(u) = L_G(u) \cup \{w\}$. Let $F := (X \cup \{u\}) \setminus N_G(u)$. Then, $F \in \Delta(G)$ and $u \in F \cap Y$. Since $N_G(u) \neq X$, $F \cap X \neq \emptyset$. By Formulas (21)–(24) and the assumption $\alpha_w = 0$, these facts give

$$n \leq \sum_{i \notin F} \alpha_i = \sum_{i \in Y} \alpha_i - \alpha_u + \alpha_w + \sum_{i \in L_G(u)} \alpha_i = n - 1 - \alpha_u,$$

and so $\alpha_u < 0$, a contradiction. Thus, $\alpha_w \geq 1$, as claimed.

Let $H := G \setminus L_G(u)$. Clearly, H is a connected bipartite graph with bipartition $(X \setminus L_G(u), Y)$. Moreover, H has only cycle C_{2k} as well. We may assume that $V(H) = \{1, \dots, s\}$. Then $s \geq 2k$ and $L_G(u) = \{s+1, \dots, r\}$. Let $\theta = (\theta_1, \dots, \theta_s) := (\alpha_1, \dots, \alpha_s) \in \mathbb{N}^s$. We now prove that

$$\Delta_\theta(I(H)^n) = \langle X \setminus L_G(u), Y \rangle. \quad (26)$$

Indeed, by (24) we get $\sum_{i \in L_G(u)} \alpha_i = 0$. Together with Formula (21), this fact gives

$$\sum_{i \in V(H), i \notin Y} \theta_i = \sum_{i \in V(H), i \notin Y} \alpha_i = \sum_{i \in X \setminus L_G(u)} \alpha_i + \sum_{i \in L_G(u)} \alpha_i = \sum_{i \in X} \alpha_i = n - 1.$$

Hence, by Lemma 1.5, $Y \in \Delta_\theta(I(H)^n)$. Similarly, $X \setminus L_G(u) \in \Delta_\theta(I(H)^n)$. Now let F' be any facet of $\Delta(H)$ which is different from $X \setminus L_G(u)$ and Y .

If $u \in F'$ then F' is also a facet of $\Delta(G)$. By noticing that F' is different from X and Y and $\sum_{i \in L_G(u)} \alpha_i = 0$, so by (22) we have

$$\sum_{i \in V(H), i \notin F'} \theta_i = \sum_{i \in V(H), i \notin F'} \alpha_i + \sum_{i \in L_G(u)} \alpha_i = \sum_{i \notin F'} \alpha_i \geq n,$$

and so $F' \notin \Delta_\theta(I(H)^n)$.

If $u \notin F'$, then $w \in F'$ since u is a leaf of H , hence $F' \cup L_G(u)$ is a facet of $\Delta(G)$. Similarly, we have $F' \notin \Delta_\theta(I(H)^n)$, and the formula (26) follows.

Define $\gamma = (\gamma_1, \dots, \gamma_s) \in \mathbb{Z}^s$ by

$$\gamma_j := \begin{cases} \theta_j - 1 & \text{if } j = u \text{ or } j = w, \\ \theta_j & \text{otherwise.} \end{cases}$$

From Inequalities (24) and (25), we have $\gamma_u = \theta_u - 1 = \alpha_u - 1 \geq 0$ and $\gamma_w = \theta_w - 1 = \alpha_w - 1 \geq 0$, so $\gamma \in \mathbb{N}^s$. Note that

$$n - 1 = \sum_{i \in X} \alpha_i \geq \alpha_u \geq 1.$$

Therefore, by Lemma 5.2 we have $\Delta_\gamma(I(H)^{n-1}) = \Delta_\theta(I(H)^n)$. Together with (26) we get

$$\Delta_\gamma(I(H)^{n-1}) = \langle X \setminus L_G(1), Y \rangle.$$

Hence, by Lemma 3.1 we have $\text{depth } S/I(H)^{n-1} = 1$, where $S = K[x_1, \dots, x_s]$. By Theorem 4.4 we have $\text{dstab}(I(H)) \leq n - 1$. On the other hand, since $v(H) = v(G) - |L_G(u)| < v(G)$, by the induction hypothesis we have $\text{dstab}(H) \leq v(H) - \varepsilon_0(H) - k + 1$. As $\{w, u\}$ is not a leaf edge of G and recall that $H = G \setminus L_G(u)$, we conclude that $\varepsilon_0(G) = \varepsilon_0(H) + |L_G(u)| - 1$. Thus,

$$v(G) - \varepsilon_0(G) - k + 1 = v(H) + |L_G(u)| - (\varepsilon_0(H) + |L_G(u)| - 1) - k + 1 = v(H) - \varepsilon_0(H) - k.$$

Hence, $n - 1 \geq \text{dstab}(I(H)) \geq v(G) - \varepsilon_0(G) - k$, and hence $n \geq v(G) - \varepsilon_0(G) - k + 1$. Thus, the proof now is complete. \square

Finally, we compute $\text{dstab}(I(G))$ for trees G . If a tree G has a vertex x being adjacent to every other vertex, then G is called a star with a center x . Note that G is a star if and only if $\text{diam}(G) \leq 2$ where $\text{diam}(G)$ stands for the diameter of G . If $\text{diam}(G) = d$, then there is a path $x_1 x_2 \dots x_d x_{d+1}$ of length d in G . Such a path will be referred to as a path realizing the diameter of G .

Lemma 5.4. $\text{dstab}(I(G)) = v(G) - \varepsilon_0(G)$ for all trees G .

Proof. Let $n := \text{dstab}(I(G))$. By Theorem 4.6 we have $n \leq v(G) - \varepsilon_0(G)$. So it remains to show $n \geq v(G) - \varepsilon_0(G)$.

If G is a star, then $\varepsilon_0(G) = \varepsilon(G) = v(G) - 1$, and then $v(G) - \varepsilon_0(G) = 1 \leq n$. Thus, the lemma holds for this case.

We will prove by induction on $v(G) = r$. If $v(G) = 2$, then G is one edge, and then it is a star. This case is already proved.

If $v(G) \geq 3$. We may assume that G is not a star so that $\text{diam } G \geq 3$. Since $\text{depth } R/I(G)^n = 1$, there is $\alpha = (\alpha_1, \dots, \alpha_r) \in \mathbb{N}^r$ such that $\Delta_\alpha(I(G)^n) = \langle X, Y \rangle$ where (X, Y) is a bipartition of G .

Let $vuw \dots z$ be a path realizing the diameter of G . Then v is a leaf, u and w both are not leaves. By [20, Lemma 3.3] we have $N_G(u) = \{w\} \cup L_G(u)$. And now we prove $n \geq v(G) - \varepsilon_0(G)$ by the same way as in Case 2 in the proof of Lemma 5.3. Thus we only sketch the proof here:

First, we show that $\alpha_u \geq 1, \alpha_w \geq 1$ and $\alpha_i = 0$ for every $i \in L_G(u)$. Then, let $T := G \setminus L_G(u)$. Note that T is also a tree and $v(G) - \varepsilon_0(G) = v(T) - \varepsilon_0(T) + 1$. We may assume that $u \in Y$, $w = s - 1$, $u = s$ and $L_G(u) = \{s + 1, \dots, r\}$. Let $\theta := (\alpha_1, \dots, \alpha_{s-2}, \alpha_{s-1} - 1, \alpha_s - 1) \in \mathbb{N}^s$. Then, we show that

$$\Delta_\theta(I(T)^{n-1}) = \langle X \setminus L_G(u), Y \rangle.$$

This gives $\text{depth } S/I(T)^{n-1} = 1$ where $S = K[x_1, \dots, x_s]$. By the induction hypothesis we have $n - 1 \geq v(T) - \varepsilon_0(T)$. From that we obtain $n \geq v(G) - \varepsilon_0(G)$. \square

Remark 5.5. Let G be a unicyclic bipartite graph. If the unique circle of G is C_4 of length 4, by the same argument as in the proof of [Lemma 5.3](#) we have the following situations:

- (1) If $G = C_4$, then $\text{dstab}(I(G)) = 1$.
- (2) If $G \neq C_4$ and C_4 has at least two adjacent vertices of degree 2 in G , then $\text{dstab}(I(G)) = v(G) - \varepsilon_0(G) - 2$.
- (3) In the remaining cases, $\text{dstab}(I(G)) = v(G) - \varepsilon_0(G) - 1$.

Thus if every connected component of G is either a tree or a unicyclic graph, then we can compute $\text{dstab}(I(G))$ by using [Theorem 4.4](#), [Lemmas 2.2, 5.3, 5.4](#) and [Remark 5.5](#).

Acknowledgments

I would like to thank Professors L.T. Hoa and N.V. Trung for helpful comments. I would like to thank the referee for his/her careful reading and many useful suggestions. A part of this work was carried out while I visited Genoa University under the support from EMMA in the framework of the EU Erasmus Mundus Action 2. I would like to thank them and Professor A. Conca for support and hospitality. This work is also partially supported by NAFOSTED (Vietnam), Project 101.01-2011.48.

References

- [1] J.A. Bondy, U.S.R. Murty, *Graph Theory*, Springer, 2008.
- [2] M. Brodmann, The asymptotic nature of the analytic spread, *Math. Proc. Cambridge Philos. Soc.* 86 (1979) 35–39.
- [3] W. Bruns, J. Herzog, *Cohen–Macaulay Rings*, Cambridge Studies in Advanced Mathematics, vol. 39, Cambridge University Press, Cambridge, 1993.
- [4] J. Chen, S. Morey, A. Sung, The stable set of associated primes of the ideal of a graph, *Rocky Mountain J. Math.* 32 (2002) 71–89.
- [5] R.C. Cowsik, M.V. Nori, Fibers of blowing up, *J. Indian Math. Soc.* 40 (1976) 217–222.
- [6] D. Eisenbud, C. Huneke, Cohen–Macaulay Rees algebras and their specializations, *J. Algebra* 81 (1983) 202–224.
- [7] D.H. Giang, L.T. Hoa, On local cohomology of a tetrahedral curve, *Acta Math. Vietnam.* 35 (2) (2010) 229–241.
- [8] J. Herzog, T. Hibi, The depth of powers of an ideal, *J. Algebra* 291 (2005) 534–550.
- [9] J. Herzog, A. Rauf, M. Vladioiu, The stable set of associated prime ideals of a polymatroidal ideal, *J. Algebraic Combin.* 37 (2) (2013) 289–312.

- [10] J. Herzog, A.A. Qureshi, Persistence and stability properties of powers of ideals, *J. Pure Appl. Algebra* 219 (2015) 530–542.
- [11] L.T. Hoa, N.D. Tam, On some invariants of a mixed product of ideals, *Arch. Math. (Basel)* 94 (4) (2010) 327–337.
- [12] L.T. Hoa, T.N. Trung, Partial Castelnuovo–Mumford regularities of sums and intersections of powers of monomial ideals, *Math. Proc. Cambridge Philos. Soc.* 149 (2010) 1–18.
- [13] M. Hochster, Rings of invariants of tori, Cohen–Macaulay rings generated by monomials, and polytopes, *Ann. of Math.* 96 (1972) 318–337.
- [14] C. Huneke, On the associated graded ring of an ideal, *Illinois J. Math.* 26 (1982) 121–137.
- [15] J. Martinez-Bernal, S. Morey, R.H. Villarreal, Associated primes of powers of edge ideals, *Collect. Math.* 63 (3) (2012) 361–374.
- [16] E. Miller, B. Sturmfels, *Combinatorial Commutative Algebra*, Springer, 2005.
- [17] N.C. Minh, N.V. Trung, Cohen–Macaulayness of powers of two-dimensional squarefree monomial ideals, *J. Algebra* 322 (2009) 4219–4227.
- [18] N.C. Minh, N.V. Trung, Cohen–Macaulayness of monomial ideals and symbolic powers of Stanley–Reisner ideals, *Adv. Math.* 226 (2) (2011) 1285–1306.
- [19] N. Terai, N.V. Trung, Cohen–Macaulayness of large powers of Stanley–Reisner ideals, *Adv. Math.* 229 (2012) 711–730.
- [20] S. Morey, Depths of powers of the edge ideal of a tree, *Comm. Algebra* 38 (11) (2010) 4042–4055.
- [21] J.R. Munkres, *Elements of Algebraic Topology*, Addison–Wesley, 1984.
- [22] A. Simis, W.V. Vasconcelos, R.H. Villarreal, On the ideal theory of graphs, *J. Algebra* 167 (1994) 389–416.
- [23] R.P. Stanley, *Combinatorics and Commutative Algebra*, second edition, Birkhäuser, Boston, MA, 1996.
- [24] Y. Takayama, Combinatorial characterizations of generalized Cohen–Macaulay monomial ideals, *Bull. Math. Soc. Sci. Math. Roumanie (N.S.)* 48 (2005) 327–344.
- [25] W. Vasconcelos, *Integral Closure, Rees Algebras, Multiplicities, Algorithms*, Springer Monographs in Mathematics, Springer-Verlag, Berlin, 2005.