



**Definition.** Given two real nonzero sequences  $\lambda = \{\lambda_i\}_{i=1}^\infty$  and  $\mathbf{c} = \{c_i\}_{i=1}^\infty$  and a fixed natural number  $k \geq 1$ , we define the following bidimensional discrete  $k$ -delay dynamical system:

$$u_{m,n+1} = \begin{cases} \lambda_m u_{m,n} & \text{if } n \in \mathbb{N} \text{ and } m = 1, 2, 3, \dots, k \\ \lambda_m u_{m,n} + c_{m-k} u_{m-k,n} & \text{if } m, n \in \mathbb{N} \text{ and } m \geq k + 1. \end{cases} \quad (1)$$

Our interest in this family of systems has several reasons. First, notice that system (1) is basically a linear system and convergence of finite dimension linear systems is determined by their associated eigenvalues. In our case, the convergence of (1) will be determined by the sequence  $\{c_i\}_{i=1}^\infty$ , which introduces the delay and by its eigenvalues, which are given by sequence  $\{\lambda_i\}_{i=1}^\infty$ . Second, a sufficient condition for convergence of finite dimensional systems is that their eigenvalues belong to the unit ball; for system (1) this is not the case, we have to impose extra conditions on the eigenvalues and also on the sequence  $\{c_i\}_{i=1}^\infty$ , which will play an important role as we will discover later on. Third, besides its intrinsic value, it is important for their role in several applications for example by approximating the equation

$$\frac{\partial u(x, t)}{\partial t} + \alpha_1(x) \frac{\partial u(x, t)}{\partial x} + \alpha_2(x) \frac{\partial^k u(x, t)}{\partial x^k} = \beta(x)u(x, t),$$

via finite differences equations. For the higher order derivative one can construct approximations based on  $k$ -delay systems by sampling a number of points to the left or to the right of the center point. In particular if we set  $k = 2$ , then the system (1) can be viewed as a consistent discretization via finite differences on an infinite rectangular grid of the one dimensional advection–diffusion–reaction equation with weak diffusion by setting  $u_{m,n} = u(mh, nk)$ ,  $\epsilon = h$ ,  $c_{m-2} = -\alpha_2(mh) \frac{k}{h^2}$  and  $\lambda_m = -\alpha_1(mh) \frac{k}{h} - \alpha_2(mh) \frac{k}{h^2} + \beta(mh)k$  with  $h$  and  $k$  being the step sizes and the indexes  $m$  and  $n$  correspond to the space variable  $x$  and to the time variable  $t$ , respectively, see [11].

2.1. Decoupling

In many cases, delay ordinary differential equations are equivalent to systems of non-delay ordinary differential equations. In a previous work, [13], a discrete dynamical system with a bidiagonal structure was analyzed, which corresponds to the problem studied here, in the special case of  $k = 1$ . Our next goal here will be to decouple system (1) obtaining the equivalent representation of  $k$  discrete 1-delay dynamical systems as follows: Let  $\ell$  be a fixed natural number satisfying  $1 \leq \ell \leq k$ , then for all  $m, n \in \mathbb{N}$ , we have the system:

$$v_{m,n+1}(\ell) = \begin{cases} \mu_1(\ell) v_{1,n}(\ell) & \text{if } m = 1 \\ \mu_m(\ell) v_{m,n}(\ell) + c_{m-1}(\ell) v_{m-1,n}(\ell) & \text{if } m \geq 2, \end{cases} \quad (2)$$

where  $v_{m,n}(\ell) = u_{\ell+(m-1)k,n}$ ,  $\mu_m(\ell) = \lambda_{\ell+(m-1)k}$  for all  $m, n \in \mathbb{N}$  with  $m \geq 1$  and  $c_m(\ell) = c_{\ell+(m-2)k}$  for all  $m \in \mathbb{N}$  with  $m \geq 2$ . System (1) can also be written in the following matrix form

$$\hat{u}_{n+1} = A_{c,\lambda} \hat{u}_n, \quad (3)$$

where  $A_{c,\lambda}$  is the infinite lower matrix

$$A_{c,\lambda} = \begin{pmatrix} \lambda_1 & 0 & 0 & 0 & 0 & \dots \\ 0 & \lambda_2 & 0 & 0 & 0 & \dots \\ \vdots & 0 & \ddots & 0 & 0 & \dots \\ 0 & 0 & 0 & \ddots & 0 & \dots \\ c_1 & 0 & 0 & 0 & \lambda_{k+1} & \vdots \\ 0 & c_2 & 0 & \ddots & \ddots & \ddots \\ \vdots & 0 & \ddots & \ddots & \ddots & \ddots \end{pmatrix} \quad (4)$$

and  $\hat{u}_n$  is the real column sequence  $\hat{u}_n = (u_{1,n}, u_{2,n}, \dots)$ . The matrix formulation of the problem is of paramount importance because it will allow us to show the development of the delay in the structure of the coefficient matrix for different iterations of the system (1).

We can also restate in matrix notation each of the  $k$  1-delay systems by using the following algorithm, which is illustrated in an example below:

**Algorithm 2.1** (Decoupling). Generation of  $k$  1-delay discrete systems.

For each value of  $\ell$  from 1, 2, 3, ...,  $k$  perform the following steps; each value of  $\ell$  will give rise to a 1-delay discrete system.

1. Replace the variable  $u_1$  with  $u_\ell$ . This amounts to interchanging rows 1 and  $\ell$  and columns 1 and  $\ell$  in matrix (4).
2. Move the  $c_\ell$  value that is located at the entry  $k + \ell, 1$  to the entry 2, 1 by interchanging rows  $k + \ell$  and 2 and also columns  $k + \ell$  and 2.

This step alters the lower triangular structure of the matrix, since the value  $c_2$  will be located in the entry  $k + 2, k + \ell$ , which will be corrected in the next step.

3. (Intermediate step) Interchange rows  $k + 2$  and  $k + 3$  and columns  $k + 2$  and  $k + 3$  in order to recover the lower triangular form of the coefficient matrix.
4. Continue moving all subsequent nonzero subdiagonal values as in steps 2 and 3. That is, for  $m = 1, 2, 3, \dots$  move the  $c_{\ell+mk}$  value by interchanging rows  $m + 2$  and  $\ell + mk + 3$  and columns  $m + 2$  and  $\ell + mk + 3$  then perform the intermediate step by exchanging rows  $m + 3$  and  $\ell + mk + 3$  and columns  $m + 3$  and  $\ell + mk + 3$ .

The following example with  $k = \ell = 3$  illustrates the steps of the algorithm.

**Example** (Decoupling algorithm).

$$A_{c,\lambda} = \begin{pmatrix} \lambda_1 & 0 & 0 & 0 & 0 & \dots \\ 0 & \lambda_2 & 0 & 0 & 0 & \dots \\ 0 & 0 & \lambda_3 & 0 & 0 & \dots \\ c_1 & 0 & 0 & \lambda_4 & 0 & \dots \\ 0 & c_2 & 0 & 0 & \lambda_5 & \vdots \\ 0 & 0 & c_3 & \ddots & \ddots & \ddots \\ \vdots & 0 & \ddots & \ddots & \ddots & \ddots \end{pmatrix}. \tag{5}$$

After the first step (interchanging row 1 and row 3 and column 1 by column 3) matrix (5) becomes

$$A_{c,\lambda} = \begin{pmatrix} \lambda_3 & 0 & 0 & 0 & 0 & \dots \\ 0 & \lambda_2 & 0 & 0 & 0 & \dots \\ 0 & 0 & \lambda_1 & 0 & 0 & \dots \\ 0 & 0 & c_1 & \lambda_4 & 0 & \dots \\ 0 & c_2 & 0 & 0 & \lambda_5 & \vdots \\ c_3 & 0 & 0 & \ddots & \ddots & \ddots \\ \vdots & 0 & \ddots & \ddots & \ddots & \ddots \end{pmatrix}. \tag{6}$$

After the second step (interchanging rows  $k + \ell = 6$  and 2) matrix (6) becomes

$$A_{c,\lambda} = \begin{pmatrix} \lambda_3 & 0 & 0 & 0 & 0 & 0 & \dots \\ c_3 & \lambda_6 & 0 & 0 & 0 & 0 & \dots \\ 0 & 0 & \lambda_1 & 0 & 0 & 0 & \dots \\ 0 & 0 & c_1 & \lambda_4 & 0 & 0 & \dots \\ 0 & 0 & 0 & 0 & \lambda_5 & c_2 & 0 \\ 0 & 0 & 0 & \ddots & 0 & \lambda_2 & 0 \\ \vdots & 0 & \vdots & c_4 & \ddots & 0 & \ddots \\ \vdots & 0 & \vdots & 0 & c_5 & \ddots & \ddots \\ \vdots & c_6 & \vdots & 0 & \ddots & \ddots & \ddots \end{pmatrix}. \tag{7}$$

Notice that the value  $c_2$  is the upper half position of the matrix breaking the lower structure of the matrix. This is where the intermediate step corrects the structure of the matrix, thus obtaining

$$A_{c,\lambda} = \begin{pmatrix} \lambda_3 & 0 & 0 & 0 & 0 & 0 & \dots \\ c_3 & \lambda_6 & 0 & 0 & 0 & 0 & \dots \\ 0 & 0 & \lambda_1 & 0 & 0 & 0 & \dots \\ 0 & 0 & c_1 & \lambda_4 & 0 & 0 & \dots \\ 0 & 0 & 0 & 0 & \lambda_2 & 0 & 0 \\ 0 & 0 & 0 & \ddots & c_2 & \lambda_5 & 0 \\ \vdots & 0 & \vdots & c_4 & \ddots & 0 & \ddots \\ \vdots & 0 & \vdots & 0 & c_5 & \ddots & \ddots \\ \vdots & c_6 & \vdots & 0 & \ddots & \ddots & \ddots \end{pmatrix}. \tag{8}$$

Finally, step 4 moves all the values  $c_6, c_9, c_{12}$  to its respective position giving rise to the coefficient matrix sought, namely

$$A_{c,\lambda} = \begin{pmatrix} \lambda_3 & 0 & 0 & 0 & 0 & 0 & \dots \\ c_3 & \lambda_6 & 0 & 0 & 0 & 0 & \dots \\ 0 & c_6 & \lambda_9 & 0 & 0 & 0 & \dots \\ 0 & 0 & c_9 & \lambda_{12} & 0 & 0 & \dots \\ 0 & 0 & 0 & c_{12} & \lambda_{15} & 0 & 0 \\ 0 & 0 & 0 & \ddots & c_{15} & \lambda_{18} & 0 \\ \vdots & 0 & \vdots & 0 & \ddots & c_{18} & \ddots \\ \vdots & 0 & \vdots & 0 & 0 & \ddots & \ddots \\ \vdots & 0 & \vdots & 0 & \ddots & \ddots & \ddots \end{pmatrix}.$$

**Algorithm 2.1** resembles the Gaussian elimination algorithm [14] in the sense that both deal with the entries below the diagonal; in this algorithm by moving the nonzero values to its higher position below the diagonal and in the Gaussian elimination by generating zeros under the diagonal. The main difference between both algorithms is in the intermediate step in **Algorithm 2.1**. It is important to remark that there are several techniques to move the nonzero entries below the diagonal to their right position but most of them destroy the geometry of the lower triangular matrix.

**Remark.** Although the algorithm was formulated for a matrix with the special structure as in (4), it is clear that an analogous algorithm can be applied to decouple any system whose coefficient matrix is infinite lower triangular and has at most one nonzero entry in each column and in each row of its main diagonal.

We can replace **Algorithm 2.1** by a more precise procedure by performing the following change of variables in (3):

$$\hat{u}_n = F_\ell E \hat{v}_n(\ell), \quad \ell = 1, 2, \dots, k, \tag{9}$$

where

$$F_\ell = \prod_{m=0}^{\infty} E_{m k+1, m k+\ell} \quad \text{and} \quad E = \prod_{m=2}^{\infty} (E_{m, m+k-1} E_{m+k, m k+1})$$

here  $E_{s,t}$  denotes the identity infinity matrix with the columns  $s$  and  $t$  permuted. The reason for having a double row–column interchange in each step (see the double factors in the product of the matrix  $E$ ) is to keep the lower triangular structure on the coefficient matrix. On the other hand, the matrix  $F_\ell$  basically is used to change the ordered set of variables

$$(\hat{u}_1, \hat{u}_2, \hat{u}_3, \hat{u}_4, \dots)$$

into the ordered set

$$(\hat{u}_\ell, \hat{u}_{\ell+k}, \hat{u}_{\ell+2k}, \hat{u}_{\ell+3k}, \hat{u}_{\ell+4k}, \dots) \quad \text{for } \ell = 1, 2, \dots, k.$$

Thus each 1-delay system, there are  $k$  of them, can be written as:

$$\hat{v}_{n+1}(\ell) = B_{c,\lambda,\ell} \hat{v}_n(\ell), \quad \ell = 1, 2, 3, \dots, k, \tag{10}$$

where  $B_{c,\lambda,\ell}$  is the infinite lower matrix

$$B_{c,\lambda,\ell} = \begin{pmatrix} \lambda_\ell & 0 & 0 & 0 & 0 & \dots \\ c_\ell & \lambda_{\ell+k} & 0 & 0 & 0 & \dots \\ 0 & c_{\ell+k} & \lambda_{\ell+2k} & 0 & 0 & \dots \\ 0 & 0 & c_{\ell+2k} & \ddots & 0 & \dots \\ 0 & 0 & 0 & \ddots & \ddots & \vdots \\ \vdots & \vdots & \vdots & \ddots & \ddots & \ddots \end{pmatrix} \tag{11}$$

and  $\hat{v}_n(\ell)$  is the real column sequence

$$\hat{v}_n(\ell) = (u_{\ell,n}, u_{\ell+k,n}, \dots). \tag{12}$$

Notice that in matrix (11), once that we fixed the values of  $\ell$ , the main diagonal elements form a subsequence of  $\lambda = \{\lambda_j\}_{j=1}^{\infty}$  where the original generic subindex  $j$  has been replaced by  $\ell + (j - 1)k$ . The previous results can be summarized in the following result.

**Theorem 2.1.** The discrete  $k$ -delay dynamical system

$$\hat{u}_{n+1} = A_{c,\lambda} \hat{u}_n, \tag{13}$$

with  $A_{c,\lambda}$  given as in (4), can be written as  $k$  discrete 1-delay dynamical systems of the form:

$$\hat{v}_{n+1}(\ell) = B_{c,\lambda,\ell} \hat{v}_n(\ell), \tag{14}$$

where  $\ell$  is a fixed natural number satisfying  $1 \leq \ell \leq k$ ,  $B_{c,\lambda,\ell}$  is given by (11) and the sequence  $\hat{v}_n(\ell)$  is given by (12).

Once we get written the family of systems (14) as in (10), then the value of each sequence  $\hat{v}_n$  is known since  $\hat{v}_n = B_{c,\lambda,\ell}^n \hat{v}_0$  where  $\hat{v}_0$  is an initial sequence. Therefore we need to compute the powers of the infinite matrix  $B_{c,\lambda,\ell}$ . From [11] we get that:

**Corollary 2.2.** The entries of  $B_{c,\lambda,\ell}^n$  are given by

$$(B_{c,\lambda,\ell}^n)_{i,j} = \begin{cases} 0 & \text{if } i < j \\ \lambda_{\ell+(i-1)k}^n & \text{if } i = j \\ \left( \prod_{m=j}^{i-1} c_{\ell+(m-1)k} \right) \cdot \mathbf{P}_{n-i+j} [\lambda_{\ell+(j-1)k}, \dots, \lambda_{\ell+(i-1)k}] & \text{if } i > j \end{cases}$$

where we have used the following notation (obtained from [12]): Given any three natural numbers  $\ell, m$  and  $n$ , let  $\mathbf{P}_\ell[\lambda_m, \lambda_{m+2}, \lambda_{m+4}, \dots, \lambda_{m+2(n-1)}]$  denote the homogeneous polynomial of degree  $\ell$  consisting of the sum of all possible monomials with unitary coefficients in the  $n$  variables  $\lambda_m, \lambda_{m+2}, \dots, \lambda_{m+2(n-1)}$ . By convention  $\mathbf{P}_0[\lambda_m, \lambda_{m+2}, \dots, \lambda_{m+2(n-1)}] = 1$  and  $\mathbf{P}_{-\ell}[\lambda_m, \lambda_{m+2}, \dots, \lambda_{m+2(n-1)}] = 0$ , for all  $\ell \in \mathbb{N}$ .

So in essence we have decomposed the system into  $k$  decoupled 1-delay discrete dynamical systems and theoretically the solutions to each system can be computed and then integrate them into a general solution. When the value of  $k$  is small the previous procedure is very practical but for larger values of  $k$  it is impractical. Thus, we have to find a different strategy to compute the solutions for (14) as a complete set.

2.2. Entries of the operator  $(A_{c,\lambda}^n)$

The previous result gives us an explicit suggestion for the functional form for the entries of the operator  $(A_{c,\lambda}^n)$ , but we need to take in consideration the delay. Thus we get the following result, which generalizes Theorem 3.5 of [13]. That result was stated and proved only for the special case  $k = 1$ .

**Theorem 2.3.** The entries  $(A_{c,\lambda}^n)_{i,j}$  for  $i, j \in \mathbb{N}$  are given by

$$\begin{cases} 0 & \text{if } i < j \\ \lambda_i^n & \text{if } i = j \\ 0 & \text{if } i > j \text{ and } i \not\equiv j \pmod{k} \\ \left( \prod_{m=0}^{s-1} c_{j+mk} \right) \cdot \mathbf{P}_{n-s} [\lambda_j, \lambda_{j+k}, \dots, \lambda_i] & \text{if } i > j \text{ and } i \equiv j \pmod{k} \end{cases}$$

where  $s = (i - j)/k$ .

**Proof.** The cases when  $i \leq j$  or when  $i > j$  and  $i \not\equiv j \pmod{k}$  are straightforward. For  $i > j$  and  $i \equiv j \pmod{k}$  the proof is achieved by induction over  $n$ . In this case  $i = j + ks$  for some natural number  $s > 1$ . When  $n = 1$  we have that  $(A_{c,\lambda})_{i,j} = c_j$  and since  $\mathbf{P}_{1-s}$  is different than zero only if  $s = 1$  then we get:

$$\left( \prod_{m=0}^{s-1} c_{j+km} \right) \cdot \mathbf{P}_{1-s} [\lambda_j, \dots, \lambda_{j+k}] = \left( \prod_{m=0}^0 c_{j+km} \right) \cdot \mathbf{P}_0 [\lambda_j, \lambda_{j+k}] = c_j.$$

Suppose now that  $(A_{c,\lambda}^n)_{i,j} = \left( \prod_{m=0}^{s-1} c_{j+mk} \right) \cdot \mathbf{P}_{n-s} [\lambda_j, \lambda_{j+k}, \dots, \lambda_i]$  is valid for a fixed  $n \in \mathbb{N}$  for all  $i, j \in \mathbb{N}$ , then

$$\begin{aligned} (A_{c,\lambda}^{n+1})_{i,j} &= \lambda_i (A_{c,\lambda}^n)_{i,j} + c_{i-k} (A_{c,\lambda}^n)_{i-k,j} \\ &= \lambda_i \left[ \left( \prod_{l=0}^{s-1} c_{j+l} \right) \mathbf{P}_{n-s} [\lambda_j, \lambda_{j+k}, \dots, \lambda_i] \right] + c_{i-k} \left[ \left( \prod_{l=0}^{i-k-j-1} c_{j+l} \right) \mathbf{P}_{n-\frac{1}{k}(i-k-j)} [\lambda_j, \lambda_{j+k}, \dots, \lambda_{i-k}] \right] \\ &= \left( \prod_{l=0}^{s-1} c_{j+l} \right) [\lambda_i \mathbf{P}_{n-s} [\lambda_j, \lambda_{j+k}, \dots, \lambda_i] + \mathbf{P}_{n+1-s} [\lambda_j, \lambda_{j+k}, \dots, \lambda_{i-k}]] \\ &= \left( \prod_{l=0}^{s-1} c_{j+l} \right) \mathbf{P}_{n+1-s} [\lambda_j, \lambda_{j+k}, \dots, \lambda_i]. \quad \blacksquare \end{aligned}$$



3.2. Different element case

As a next step we will consider another particular case: we will consider a special case where all the eigenvalues of  $A_{c,\lambda}$  are different, that is, the sequence  $\lambda$  has different elements.

**Theorem 3.1.** Assume that all eigenvalues of  $A_{c,\lambda}$  are different. Let  $(\Phi_{c,\lambda})_{ij}$  denote the  $i, j$  entry of  $\Phi_{c,\lambda}$  with

$$(\Phi_{c,\lambda})_{ij} = \begin{cases} 0 & \text{if } i < j \\ 1 & \text{if } i = j \\ 0 & \text{if } i > j \text{ and } i \not\equiv j \pmod{k} \\ \frac{1}{(\lambda_j - \lambda_i)} \prod_{m=0}^{s-1} \frac{c_{j+mk}}{\lambda_{j+mk} - \lambda_j} & \text{if } i > j \text{ and } i \equiv j \pmod{k}, \end{cases}$$

where  $s = \frac{i-j}{k}$ . Let  $\Omega_w$  be an infinite diagonal matrix with  $(\Omega_w)_{ii} = \omega(\lambda_i)$ , where  $\omega : \mathbb{R} \rightarrow \mathbb{R}$  is an arbitrary but fixed function. Then the matrix  $\Phi_{c,\lambda}$  is an invertible infinite lower triangular matrix such that  $(A_{c,\lambda}^n)_{ij} = (\Phi_{c,\lambda} \circ \Omega_w^n \circ \Phi_{c,\lambda}^{-1})_{ij}$  for any natural number  $n \geq 0$ , where the entries  $(\Phi_{c,\lambda} \circ \Omega_w \circ \Phi_{c,\lambda}^{-1})_{ij}$  are given by

$$\begin{cases} 0 & \text{if } i < j \\ \omega(\lambda_i) & \text{if } i = j \\ 0 & \text{if } i > j \text{ and } i \not\equiv j \pmod{k} \\ \left( \frac{1}{s} \prod_{m=0}^{s-1} c_{j+mk} \right) \cdot \frac{d^s}{d\mu^s} \mathcal{L}_\omega[\lambda_j, \lambda_{j+k}, \dots, \lambda_i](\mu) & \text{if } i > j \text{ and } i \equiv j \pmod{k} \end{cases}$$

where  $\mathcal{L}_\omega[\lambda_j, \lambda_{j+k}, \dots, \lambda_i](\lambda)$  is the Lagrange polynomial in the variable  $\mu$  of degree  $s$  at the nodes  $\{\lambda_j, \lambda_{j+k}, \dots, \lambda_i\}$  for the function  $\omega(\mu)$  [15].

**Proof.** It is straightforward to verify that the matrix  $\Phi_{c,\lambda}^{-1}$  exists. Moreover its entries are given by

$$(\Phi_{c,\lambda}^{-1})_{ij} = \begin{cases} 0 & \text{if } i < j \\ 1 & \text{if } i = j \\ 0 & \text{if } i > j \text{ and } i \not\equiv j \pmod{k} \\ \prod_{m=0}^{\frac{i-j}{k}-1} \frac{c_{j+mk}}{\lambda_i - \lambda_{j+mk}} & \text{if } i > j \text{ and } i \equiv j \pmod{k}. \end{cases}$$

Now since all three matrices,  $\Phi$ ,  $\Omega$  and  $\Phi^{-1}$  are lower triangular, their product is also lower triangular. Thus there are several cases to show, but the only one that is not straightforward is if  $i > j$  and  $i \equiv j \pmod{k}$  then  $i = j + ks$  for some  $s \in \mathbb{N}$ . Thus the  $(\Phi_{c,\lambda} \circ \Omega_w \circ \Phi_{c,\lambda}^{-1})_{ij}$  entries are given by

$$\begin{aligned} \sum_{m=j}^i (\Phi_{c,\lambda} \circ \Omega_w)_{i,m} (\Phi_{c,\lambda}^{-1})_{m,j} &= \sum_{m=0}^s (\Phi_{c,\lambda} \circ \Omega_w)_{i,j+mk} (\Phi_{c,\lambda}^{-1})_{j+mk,j} \\ &= \omega(\lambda_i) \prod_{\ell=0}^{s-1} \frac{c_{j+k\ell}}{\lambda_i - \lambda_{j+k\ell}} + \frac{\omega(\lambda_j)}{\lambda_j - \lambda_i} \prod_{\ell=0}^{s-1} \frac{c_{j+k\ell}}{\lambda_j - \lambda_{j+k\ell}} \\ &\quad + \sum_{m=1}^{s-1} \frac{\omega(\lambda_{j+km})}{\lambda_{j+km} - \lambda_i} \left( \prod_{\ell=0}^{s-m-1} \frac{c_{j+km+k\ell}}{\lambda_{j+km} - \lambda_{j+km+k\ell}} \right) \left( \prod_{\ell=0}^{m-1} \frac{c_{j+k\ell}}{\lambda_{j+km} - \lambda_{j+k\ell}} \right) \\ &= \left( \prod_{\ell=0}^{s-1} c_{j+k\ell} \right) \left( \omega(\lambda_i) \prod_{\ell=0}^{s-1} \frac{1}{\lambda_i - \lambda_{j+k\ell}} + \frac{\omega(\lambda_j)}{\lambda_j - \lambda_i} \prod_{\ell=1}^{s-1} \frac{1}{\lambda_j - \lambda_{j+k\ell}} \right) \\ &\quad + \sum_{m=1}^{s-1} \frac{\omega(\lambda_{j+km}) c_{j+km}}{\lambda_{j+km} - \lambda_i} \left( \prod_{l=m+1}^{s-1} \frac{c_{j+k\ell}}{\lambda_{j+km} - \lambda_{j+k\ell}} \right) \left( \prod_{l=0}^{m-1} \frac{c_{j+k\ell}}{\lambda_{j+km} - \lambda_{j+k\ell}} \right) \\ &= \left( \prod_{\ell=0}^{s-1} c_{j+k\ell} \right) \left[ \sum_{m=0}^s \omega(\lambda_{j+km}) \left( \prod_{\ell=0, \ell \neq m}^s \frac{1}{\lambda_{j+km} - \lambda_{j+k\ell}} \right) \right] \\ &= \left( \prod_{\ell=0}^{s-1} c_{j+k\ell} \right) \frac{d^s}{d\mu^s} \mathcal{L}_\omega[\lambda_j, \lambda_{j+k}, \dots, \lambda_i](\mu). \end{aligned}$$

Now let us prove that  $\Phi_{c,\lambda} \circ \Omega_w \circ \Phi_{c,\lambda}^{-1} = (A_{c,\lambda})$ . We will prove such equality proving that equality is achieved entry-wise. Let us notice that we only have to prove the equality for the entries  $i > j$  with  $i \equiv j \pmod{k}$ .

Since the polynomial  $p(x) = x$  interpolates at the nodes given by the eigenvalues of  $A_{c,\lambda}$  and by uniqueness of the Lagrange polynomial, we conclude that  $\mathcal{L}_w[\lambda_j, \lambda_{j+k}, \dots, \lambda_i](\mu)$  is a polynomial of degree one. Therefore its s-derivative is zero except when  $i - j = k$ . For these entries we have

$$\begin{aligned} (\Phi_{c,\lambda} \circ \Omega \circ \Phi_{c,\lambda}^{-1})_{j+k,j} &= \prod_{\ell=0}^0 c_{j+k\ell} \frac{d}{d\mu} \mathcal{L}_w[\lambda_j, \lambda_{j+k}](\mu) \\ &= c_j \frac{d}{d\mu} \left( \frac{\lambda_j(\mu - \lambda_{j+k})}{\lambda_j - \lambda_{j+k}} + \frac{\lambda_{j+k}(\mu - \lambda_j)}{\lambda_{j+k} - \lambda_j} \right) = c_j. \quad \blacksquare \end{aligned}$$

In general the eigenvalues of  $A_{c,\lambda}$  appear in the entries of  $A_{c,\lambda}^n$  as nodes of an interpolation polynomial for the function  $w(x) = x^n$  and the product of the subdiagonal elements can be interpreted as correction coefficients or weights. These two facts show a direct connection with interpolation theory. Notice that we also obtain the following double identity:

**Corollary 3.2.** *If  $\lambda = \{\lambda_j\}_{j=1}^{\infty}$  is a sequence with different elements and  $w(\lambda) = \lambda$ , then  $\forall i, j, n \in \mathbb{N}$  with  $i > j$  and  $i \equiv j \pmod{k}$*

$$\mathbf{sP}_{n-s}[\lambda_j, \lambda_{j+k}, \dots, \lambda_i] = \mathcal{L}_w^{(s)}[\lambda_j, \lambda_{j+k}, \dots, \lambda_i] = \mathbf{sf}_n[\lambda_j, \lambda_{j+k}, \dots, \lambda_i],$$

where  $f_n[\lambda_j, \lambda_{j+k}, \dots, \lambda_i]$  denotes the divided differences (see [16]) of the function  $f_n(x) = x^n$  at the nodes  $\lambda_j, \lambda_{j+k}, \dots, \lambda_i$ , see [12] for details on this last equality.

#### 4. Conclusions

One goal of this study was to analyze the result of introducing a multiple delay in the advection–reaction operators. We discovered that introducing a  $k$  multiple delay in a discrete dynamical system generated by such operators is equivalent to introducing a delay in  $k$  dynamical systems. Moreover, we were able to compute explicitly the iterates of such operator with two different approaches showing their dynamical behavior and their connection with interpolation theory.

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