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Iterative methods for least-square problems based on proper splittings[☆]

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Abstract

For the linear-squares problems $\min_x \|b - Ax\|_2$, where A is large and sparse, straightforward application of Cholesky or QR factorization will lead to catastrophic fill in factor R . We consider handling such problems by a iterative methods based on proper splittings. We establish the convergence, to the least-square solution $y = A^\dagger x$, for the sequential two-stage iterative method and for the parallel stationary iterative method.

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1. Introduction

Consider the linear system

$$Ax = b, \tag{1}$$

where A is an $m \times n$ real, large, and sparse matrix, x is the unknown real n -vector, and b is a given real m -vector.

If (1) is inconsistent, then the least-square solution, that is, the vector y of \mathbb{R}^n such that

$$\|b - Ay\|_2 = \min_x \|b - Ax\|_2$$

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is given by $\mathbf{y} = A^\dagger \mathbf{b}$ where A^\dagger is the Moore–Penrose generalized inverse of matrix A . The vector $A^\dagger \mathbf{b}$ is also the solution of the normal equation $A^T A \mathbf{x} = A^T \mathbf{b}$, but the matrix $A^T A$ has the disadvantage that it is frequently ill-conditioned and influenced greatly by roundoff error (see [5]). To avoid this disadvantage, Berman and Plemmons [2] introduce the following iterative method similar to the iterative method for nonsingular linear systems:

$$\mathbf{x}^{(k+1)} = M^\dagger N \mathbf{x}^{(k)} + M^\dagger \mathbf{b}, \quad k = 0, 1, 2, \dots, \quad (2)$$

where $A = M - N$ is a proper splitting of A ; that is M and N are $m \times n$ matrices such that

$$\mathcal{R}(M) = \mathcal{R}(A) \quad \text{and} \quad \ker(M) = \ker(A).$$

Here, $\mathcal{R}(A)$ is the range or column space of A and $\ker(A)$ is the nullspace of A . Note that for the nonsingular case, a proper splitting is a splitting (see [8]).

Berman and Plemmons [2] introduce the following convergence result.

Theorem 1 (Berman and Plemmons [2, Corollary 1]). *Let A be an $m \times n$ real matrix and let $A = M - N$ be a proper splitting. Then iterative process (2) converges to $A^\dagger \mathbf{b}$ for each initial vector $\mathbf{x}^{(0)}$ if and only if $\rho(M^\dagger N) < 1$.*

We say that matrix $A = [a_{ij}]$ is nonnegative ($A \geq 0$) if $a_{ij} \geq 0$ for $i = 1, 2, \dots, m$ and $j = 1, 2, \dots, n$. Also we say that $A \leq B$ if $B - A \geq 0$. Similar for vectors.

Definition 1. Let A be an $m \times n$ real matrix and let $A = M - N$ be a proper splitting. We say that $A = M - N$ is regular if $M^\dagger \geq 0$ and $N \geq 0$. Furthermore, we say that $A = M - N$ is weak nonnegative of the first type if $M^\dagger \geq 0$ and $M^\dagger N \geq 0$.

Climent et al. [3] extend the convergence results introduced by Berman and Plemmons [2], proving the following theorem that we quote for further references.

Theorem 2 (Climent et al. [3, Theorem 3]). *Let A be an $m \times n$ real matrix with complete rank and let $A = M - N$ be a weak nonnegative proper splitting of the first type. Then the following conditions are equivalent:*

- (i) $A^\dagger \geq 0$,
- (ii) $A^\dagger \geq M^\dagger$,
- (iii) $A^\dagger M \geq 0$,
- (iv) $\rho(M^\dagger N) = (\rho(A^\dagger M) - 1) / \rho(A^\dagger M)$,
- (v) $\rho(M^\dagger N) = \rho(NM^\dagger) < 1$,
- (vi) $(I - M^\dagger N)^{-1} \geq 0$,
- (vii) $A^\dagger N \geq 0$,
- (viii) $A^\dagger N \geq M^\dagger N$,
- (ix) $\rho(M^\dagger N) = \rho(A^\dagger N) / (1 + \rho(A^\dagger N)) < 1$.

If in the above theorem we remove the hypothesis “ A has complete rank”, then we only can establish the equivalence between parts (i), (ii) and (v)–(ix), obtaining in this way the extension of [2, Theorem 3]. That hypothesis is necessary as we can see in [3, Example 1].

It is very usual to use, to obtain the least-square solution of system (1), one of the several existing codes for the sparse Cholesky factorization of $A^T A$ or the QR factorization of matrix A , than the iterative method (2). However, when matrix A is sparse except for a few dense rows, it is an open problem that a straightforward application of Cholesky factorization or the QR factorization will lead to catastrophic fill in the factor R . Recently, authors as Adlers and Björck [1] introduce a new technique by using a matrix stretching by row splittings.

In this paper, with the purpose to introduce alternative solution to this problem, we introduce two iterative methods to obtain the least-square solution of system (1), based on the concept of proper splitting. The first one is a two-stage method and the second one is a parallel iterative method.

2. Convergence results

Let $A = M - N$ be a proper splitting of A and let $M = F - G$ be a proper splitting of M . If we consider q inner iterations with splitting $M = F - G$ we obtain the two-stage iterative method (see [6,4] for the nonsingular case)

$$\mathbf{x}^{(k+1)} = (F^\dagger G)^q \mathbf{x}^{(k)} + \sum_{j=0}^{q-1} (F^\dagger G)^j F^\dagger (N \mathbf{x}^{(k)} + \mathbf{b}). \tag{3}$$

Now, using Theorems 1 and 2 we can establish the following convergence result.

Theorem 3. *Let $A = M - N$ be a regular proper splitting, and let $M = F - G$ be a weak nonnegative proper splitting of the first type. Then the stationary two-stage iterative method (3) converges to the least-square solution of system (1) for any initial vector $\mathbf{x}^{(0)}$.*

Proof. Since $MM^\dagger N = N$, $F^\dagger M = F^\dagger F - F^\dagger G$, and $F^\dagger F = M^\dagger M$ we have, from Eq. (3), that

$$\begin{aligned} T_p &= (F^\dagger G)^p + \sum_{j=0}^{p-1} (F^\dagger G)^j F^\dagger N \\ &= (F^\dagger G)^p + \sum_{j=0}^{p-1} (F^\dagger G)^j F^\dagger M M^\dagger N \\ &= (F^\dagger G)^p + \sum_{j=0}^{p-1} (F^\dagger G)^j (F^\dagger F - F^\dagger G) M^\dagger N \\ &= (F^\dagger G)^p + (F^\dagger F - (F^\dagger G)^p) M^\dagger N \\ &= (F^\dagger G)^p + F^\dagger F M^\dagger N - (F^\dagger G)^p M^\dagger N \\ &= (F^\dagger G)^p + M^\dagger M M^\dagger N - (F^\dagger G)^p M^\dagger N \\ &= (F^\dagger G)^p + M^\dagger N - (F^\dagger G)^p M^\dagger N \end{aligned}$$

$$\begin{aligned}
 &= (F^\dagger G)^p + (I - (F^\dagger G)^p)M^\dagger N \\
 &= I - I + (F^\dagger G)^p + (I - (F^\dagger G)^p)M^\dagger N \\
 &= I - (I - (F^\dagger G)^p)(I - M^\dagger N) \\
 &= I - \sum_{j=0}^{p-1} (F^\dagger G)^j (I - F^\dagger G)(I - M^\dagger N).
 \end{aligned}$$

Now, taking into account that $A = M - N$ and $M = F - G$ are convergent, if $y \geq 0$, then $x = (I - M^\dagger N)^{-1}(I - F^\dagger G)^{-1}y \geq 0$ and, consequently,

$$0 \leq T_p x = x - \sum_{j=0}^{p-1} (F^\dagger G)^j y < x,$$

which implies that $\rho(T_p) < 1$. \square

Next, in a similar way that for weak nonnegative splittings of the first type of nonsingular matrices (see [7]), we can obtain a least-square solution of system (1) using a parallel iterative process based on a multisplitting of A .

Definition 2. Let A be an $m \times n$ real matrix. We say that $\{(M_l, N_l, E_l)\}_{l=1}^p$ is a proper multisplitting of A if $A = M_l - N_l$, for $l = 1, \dots, p$, is a proper splitting, E_l , for $l = 1, \dots, p$, is a nonnegative and diagonal $m \times m$ matrix, and $\sum_{l=1}^p E_l = I_m$ where I_m is the $m \times m$ identity matrix.

As a generalization of Definition 1, we say that a proper multisplitting is regular or weak nonnegative of the first type, if each one of the proper splittings is regular or weak nonnegative of the first type, respectively.

If $\{(M_l, N_l, E_l)\}_{l=1}^p$ is a proper multisplitting of A , we can consider the iterative scheme

$$x^{(k+1)} = Hx^{(k)} + Gb, \quad k = 0, 1, 2, \dots, \tag{4}$$

where

$$H = \sum_{l=1}^p E_l M_l^\dagger N_l \quad \text{and} \quad G = \sum_{l=1}^p E_l M_l^\dagger.$$

Before establishing the convergence result given in Theorem 4 for the above iterative scheme, we need the following result.

Lemma 1. Let A be an $m \times n$ real matrix and let $\{(M_l, N_l, E_l)\}_{l=1}^p$ be a proper multisplitting of the first type of A . Then

- (i) $H \geq 0$ and therefore $H^j \geq 0$, for $j = 0, 1, \dots$,
- (ii) $\sum_{l=1}^p E_l M_l^\dagger A = (I - H)A^\dagger A$,
- (iii) $(I + H + H^2 + \dots + H^m)(I - H) = I - H^{m+1}$.

Proof. Part (i) is a consequence from the definition of matrix H and the definition of weak nonnegative proper splitting of the first type. Part (iii) is a consequence of part (i).

To prove part (iii) we proceed as follows. From $M_l M_l^\dagger N_l = N_l = N_l M_l^\dagger M_l$ and $M_l^\dagger M_l = A^\dagger A$ for $l = 1, 2, \dots, p$ we have that

$$A = M_l - N_l = M_l(I - M_l^\dagger N_l), \quad l = 1, 2, \dots, p$$

and then

$$\begin{aligned} \sum_{l=1}^p E_l M_l^\dagger A &= \sum_{l=1}^p E_l M_l^\dagger M_l (I - M_l^\dagger N_l) \\ &= \sum_{l=1}^p E_l (M_l^\dagger M_l - M_l^\dagger M_l M_l^\dagger N_l) \\ &= \sum_{l=1}^p E_l (M_l^\dagger M_l - M_l^\dagger N_l M_l^\dagger M_l) \\ &= \sum_{l=1}^p E_l ((I - M_l^\dagger N_l) M_l^\dagger M_l) \\ &= \left(I - \sum_{l=1}^p E_l M_l^\dagger N_l \right) A^\dagger A \\ &= (I - H) A^\dagger A, \end{aligned}$$

that is, part (ii) holds. \square

Theorem 4. Let A be an $m \times n$ real matrix and let $\{(M_l, N_l, E_l)\}_{l=1}^p$ be a proper multisplitting of the first type of A . Then iterative method (4) converges to the least-square solution of system (1) for any initial vector $\mathbf{x}^{(0)}$.

Proof. Using Lemma 1 and taking into account that $M_l^\dagger M_l M_l^\dagger = M_l^\dagger$, and $M_l M_l^\dagger = A A^\dagger$ for $l = 1, 2, \dots, p$, and $A^\dagger A A^\dagger = A^\dagger$, we obtain

$$\begin{aligned} 0 &\leq (I + H + H^2 + \dots + H^m) \sum_{l=1}^p E_l M_l^\dagger \\ &= (I + H + H^2 + \dots + H^m) \sum_{l=1}^p E_l M_l^\dagger M_l M_l^\dagger \\ &= (I + H + H^2 + \dots + H^m) \sum_{l=1}^p E_l M_l^\dagger A A^\dagger \\ &= (I + H + H^2 + \dots + H^m) (I - H) A^\dagger A A^\dagger \end{aligned}$$

$$\begin{aligned} &= (I + H + H^2 + \cdots + H^m)(I - H)A^\dagger \\ &= (I - H^{m+1})A^\dagger \leq A^\dagger. \end{aligned}$$

Therefore, the elements of H^m must remain bounded, and therefore H is convergent, that is, iterative method (4) converges to the least-square solution of system (1) for any initial vector $\mathbf{x}^{(0)}$. \square

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