



A semi-iterative method for real spectrum singular linear systems with an arbitrary index

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Received 14 February 1997; received in revised form 16 July 1997

Abstract

In this paper we develop a semi-iterative method for computing the Drazin-inverse solution of a singular linear system $Ax = b$, where the spectrum of A is real, but its index (i.e., the size of its largest Jordan block corresponding to the eigenvalue zero) is arbitrary. The method employs a set of polynomials that satisfy certain normalization conditions and minimize some well-defined least-squares norm. We develop an efficient recursive algorithm for implementing this method that has a fixed length independent of the index of A . Following that, we give a complete theory of convergence, in which we provide rates of convergence as well. We conclude with a numerical application to determine eigenprojections onto generalized eigenspaces. Our treatment extends the work of Hanke and Hochbruck (1993) that considers the case in which the index of A is 1.

Keywords: Singular systems; Iterative methods; Polynomial acceleration

AMS classification: 65F10; 65F20; 15A09

1. Introduction

Consider the linear system

$$Ax = b, \tag{1.1}$$

where $A \in \mathbb{C}^{n,n}$ is singular and $\text{ind}(A) = a$ is arbitrary. Here $\text{ind}(\cdot)$ denotes the *index* of a matrix, namely, the size of the largest Jordan block corresponding to its zero eigenvalue. The purpose of

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¹ The work of this author was done while he was visiting the Department of Mathematics at the University of Connecticut, supported by Spanish DGICYT grant number PR94-205.

this paper is to develop a semi-iterative method for computing the Drazin-inverse solution of (1.1), namely, the vector $A^D b$, where A^D is the Drazin inverse of A , in an *efficient manner*. For the Drazin inverse and its properties, see, e.g., [1] or [2].

We shall assume that

$$\sigma(A) \subseteq \{0\} \cup [c - d, c + d], \quad 0 < d < c, \quad (1.2)$$

where $\sigma(\cdot)$ denotes the spectrum of a matrix.

Our work here extends that of Hanke and Hockbruck [8] which treats the case of $a = 1$ and utilizes the general theory of Eiermann et al. [4] of semi-iterative methods for computing the Drazin-inverse solution to singular systems.

We begin with some essential background. Let x_0 be an arbitrary initial vector and let $r_0 = b - Ax_0$ be the corresponding residual vector. Then, beginning with x_0 , the m th iterate x_m is given by

$$x_m = x_0 + q_{m-1}(A)r_0 = p_m(A)x_0 + q_{m-1}(A)b, \quad (1.3)$$

where $q_{m-1}(\lambda)$ is a polynomial of degree at most $m - 1$ and $p_m(\lambda)$ is a polynomial of degree at most m given by

$$p_m(\lambda) = 1 - \lambda q_{m-1}(\lambda). \quad (1.4)$$

We call $p_m(\lambda)$ the m th *residual polynomial*. Note that

$$p_m(0) = 1. \quad (1.5)$$

As is shown by Eiermann et al. [4, Lemma 2], necessary and sufficient conditions for the convergence of the sequence $\{x_m\}_{m=0}^{\infty}$ are that

$$\lim_{m \rightarrow \infty} p_m^{(i)}(0) = 0, \quad i = 1, \dots, a, \quad (1.6)$$

and

$$\lim_{m \rightarrow \infty} p_m^{(i)}(\lambda_j) = 0, \quad i = 0, \dots, k_j - 1, \quad (1.7)$$

where λ_j are the nonzero eigenvalues of A and $k_j = \text{ind}(A - \lambda_j I)$.

The conditions in (1.6) will, of course, be satisfied if

$$p_m^{(i)}(0) = 0, \quad i = 1, \dots, a, \quad \text{for all } m = 0, 1, \dots \quad (1.8)$$

Polynomials $p_m(\lambda)$ satisfying (1.8) and (1.7) were considered by Hanke and Hockbruck [8] for the case $a = 1$. We mention in passing that the polynomials that arise in connection with the extrapolation methods for the Drazin-inverse solution studied by Sidi [10] satisfy (1.8) and (1.7) for arbitrary a .

The plan of this paper is as follows: In Section 2, using a weight function $w(\lambda)$, we provide an integral norm $\|\cdot\|$ and a set of polynomials $\{p_m(\lambda)\}_{m=0}^{\infty}$ satisfying (1.8) and (1.7) such that the norm of $p_m(\lambda)$ is minimal over the set of all polynomials $p(\lambda)$ of degree at most m which satisfy $p(0) = 1$ and $p^{(i)}(0) = 0$, $i = 1, \dots, a$. We use these polynomials to construct our semi-iterative method. Our work here extends directly the developments of Hanke and Hockbruck [8] to the case $a > 1$.

In Section 3 we develop a recursive algorithm for implementing the semi-iterative method defined by (1.3) and (1.4), $p_m(\lambda)$ being the minimal polynomials of Section 2. This algorithm involves only four successive iterates x_m , independently of the index of A . Here we make use of the fact that the $p_m(\lambda)$ can be expressed in terms of a set of orthogonal polynomials that satisfy the usual 3-term recurrence relation.

In Section 4 we prove the convergence of the method and provide error bounds and the corresponding rates of convergence for the case in which

$$w(\lambda) = \frac{1}{\sqrt{(\lambda - c + d)(c + d - \lambda)}}. \tag{1.9}$$

In particular, we show that if A satisfies (1.2), then

$$\|x_m - A^D b - \tilde{x}_0\| = O(m^{a+s}\kappa^m) \quad \text{as } m \rightarrow \infty,$$

where \tilde{x}_0 is that part of x_0 that lies in the null space of A^a , s is a nonnegative integer, and

$$\kappa = \frac{c - \sqrt{c^2 - d^2}}{d} < 1. \tag{1.10}$$

The asymptotic estimates that we give for the bounds on our residual polynomials in equation (4.17) of Theorem 4.5 do not reach, except in the case of the index A being equal to 1, the near optimal rate achieved by the residuals of Bernstein (see [6]) which is displayed here in (4.24). But we believe that our short recurrence relation for computing the residuals makes up for this deficiency.

In Section 5 we present several numerical examples in which we compute the *projections onto the generalized eigenspaces of matrices* whose spectrum is real and satisfies the condition of (1.2). The algorithm does well when the transforming matrix of A to its Jordan canonical form has a relatively low condition number.

2. Minimal polynomials

Let $w(\lambda)$ be a nonnegative weight function over the interval $[c - d, c + d]$ and let f and g be functions defined on $[c - d, c + d]$. Define the inner product $\langle \cdot, \cdot \rangle$ on $[c - d, c + d]$ by

$$\langle f, g \rangle = \int_{c-d}^{c+d} w(\lambda) f(\lambda) g(\lambda) d\lambda.$$

Next define the norm $||| \cdot |||$ via

$$|||f|||^2 = \left\langle f, \frac{1}{\lambda^a} f \right\rangle.$$

Let Π_m denote the set of all real polynomials of degree at most m and define

$$\Pi_m^0 = \{ p \in \Pi_m : p(0) = 1, p^{(i)}(0) = 0, i = 1, \dots, a \}.$$

Note that $p_m(\lambda) = 1$ is the only member of the set Π_m^0 for $m = 0, 1, \dots, a$.

Theorem 2.1. *Let a be a positive integer. Then, for $m \geq a + 1$, the minimization problem*

$$\min_{p \in \Pi_m^0} |||p||| \tag{2.1}$$

admits a unique solution $p_m(\lambda)$ which is characterized by

$$\langle p_m, \lambda^j \rangle = 0, \quad j = 1, \dots, m - a. \tag{2.2}$$

Proof. We start by noting that $p \in \Pi_m^0$ implies that $p(\lambda) = 1 - \lambda^{a+1}u(\lambda)$ with $u \in \Pi_{m-a-1}$. Thus, we have

$$|||p|||^2 = \int_{c-d}^{c+d} w(\lambda)\lambda^{a+2} [\lambda^{-a-1} - u(\lambda)]^2 d\lambda. \tag{2.3}$$

As $w(\lambda)\lambda^{a+2}$ is nonnegative on $[c - d, c + d]$, there exists a unique polynomial $u^*(\lambda)$ in Π_{m-a-1} that minimizes the integral on the right-hand side of (2.3); $u^*(\lambda)$ is the best approximation from Π_{m-a-1} to λ^{-a-1} in the norm induced by the inner product $\langle \cdot, \lambda^{a+2} \cdot \rangle$, see, e.g., [3]. Consequently, $p_m(\lambda) = 1 - \lambda^{a+1}u^*(\lambda)$ is the unique solution to (2.1).

Consider the polynomial $p(\lambda) = p_m(\lambda) + \alpha\lambda^{a+j}$, where $\alpha \in \mathbb{R}$ and $j = 1, \dots, m - a$. Obviously, $p \in \Pi_m^0$. Hence,

$$|||p_m|||^2 \leq |||p|||^2 = |||p_m|||^2 + 2\alpha\langle p_m, \lambda^j \rangle + \alpha^2|||\lambda^{a+j}|||^2.$$

Since α has an arbitrary sign, this inequality holds if and only if (2.2) holds.

Conversely, assume that (2.2) holds for some $p_m \in \Pi_m^0$. Let $p \in \Pi_m^0$. Then $p(\lambda) - p_m(\lambda)$ has a zero of multiplicity at least $a + 1$ at $\lambda = 0$. Thus,

$$v(\lambda) = \frac{p(\lambda) - p_m(\lambda)}{\lambda^a} \in \text{span}\{\lambda, \lambda^2, \dots, \lambda^{m-a}\}.$$

Consequently,

$$|||p|||^2 = |||p_m|||^2 + 2\langle p_m, v \rangle + |||\lambda^a v|||^2 = |||p_m|||^2 + |||\lambda^a v|||^2 \geq |||p_m|||^2,$$

since $\langle p_m, v \rangle = 0$ by (2.2). This means that $p_m(\lambda)$ is the unique solution to (2.1). \square

Now, let us define

$$u_m(\lambda) = \frac{p_m(\lambda) - p_{m+1}(\lambda)}{\lambda}, \quad m \geq a. \tag{2.4}$$

Clearly, $u_m(\lambda)$ is a polynomial in $\text{span}\{\lambda^a, \lambda^{a+1}, \dots, \lambda^m\}$.

Theorem 2.2. *The polynomials $\lambda^{-a}u_m(\lambda)$, $m = a, a + 1, \dots$, are orthogonal with respect to the inner product $\langle \cdot, \lambda^{a+2} \cdot \rangle$.*

Proof. First, $\lambda^{-a}u_m(\lambda)$ is of degree precisely $m - a$. For any polynomial $p(\lambda)$ in Π_{m-a-1} , with $m \geq a + 1$, we then have that

$$\langle \lambda^{-a}u_m, \lambda^{a+2}p \rangle = \langle u_m, \lambda^2p \rangle = \langle p_m - p_{m+1}, \lambda p \rangle = 0$$

by (2.2). \square

From Theorem 2.2 we now have that the $u_m(\lambda)$ satisfy a 3-term recursion relation of the form

$$u_m(\lambda) = (\omega_m \lambda + \mu_m)u_{m-1}(\lambda) + v_m u_{m-2}(\lambda), \quad m \geq a + 1 \tag{2.5}$$

for some constants ω_m , μ_m , and v_m , with $v_{a+1} = 0$.

Let us denote by $t_m(\lambda)$ the orthogonal polynomial of degree m with respect to the inner product $\langle \cdot, \cdot \rangle$ and normalized such that $t_m(0) = 1$. As a result of this normalization, the $t_m(\lambda)$ satisfy a 3-term recursion relation of the form

$$t_{m+1}(\lambda) = -\alpha_m \lambda t_m(\lambda) + (1 + \beta_m)t_m(\lambda) - \beta_m t_{m-1}(\lambda), \quad m \geq 0 \tag{2.6}$$

with

$$t_{-1}(\lambda) = 0 \quad \text{and} \quad t_0(\lambda) = 1,$$

for some constants α_m and β_m .

Theorem 2.3. *For $m \geq a$, the polynomials $p_m(\lambda)$ can be expressed in terms of the polynomials $t_j(\lambda)$ as*

$$\lambda p_m(\lambda) = \sum_{j=m-a}^{m+1} \pi_{m,j} t_j(\lambda) \tag{2.7}$$

for some constants $\pi_{m,j}$ which satisfy the linear system

$$\begin{aligned} \sum_{j=m-a}^{m+1} \pi_{m,j} &= 0, \\ \sum_{j=m-a}^{m+1} \pi_{m,j} \tau_j^{(1)} &= 1, \\ \sum_{j=m-a}^{m+1} \pi_{m,j} \tau_j^{(i)} &= 0, \quad i = 2, 3, \dots, a + 1, \end{aligned} \tag{2.8}$$

where $\tau_j^{(i)} = t_j^{(i)}(0)$.

Proof. Since $\lambda p_m(\lambda)$ is in Π_{m+1} , we have that

$$\lambda p_m(\lambda) = \sum_{j=0}^{m+1} \pi_{m,j} t_j(\lambda).$$

But

$$\langle \lambda p_m, t_j \rangle = \langle p_m, \lambda t_j \rangle = 0, \quad j = 0, 1, \dots, m - a - 1,$$

by (2.2). Therefore, $\pi_{m,j} = 0$ for $j = 0, 1, \dots, m - a - 1$. This proves (2.7). The first of the equations in (2.8) follows by letting $\lambda = 0$ in (2.7). The remaining equations in (2.8) can be obtained by taking

the i th derivative of both sides of (2.7), $i = 1, 2, \dots, a + 1$, noting that

$$[\lambda p_m(\lambda)]^{(i)} = \lambda p_m^{(i)}(\lambda) + i p_m^{(i-1)}(\lambda), \quad i = 1, 2, \dots,$$

and letting $\lambda = 0$ on both sides, and recalling that $p_m \in \Pi_m^0$. \square

We now show how the coefficients ω_m , μ_m , and ν_m in (2.5) can be determined from the constants $\pi_{m,j}$ in (2.7) and the coefficients α_m and β_m in (2.6).

Theorem 2.4. *Let*

$$\gamma_m := \pi_{m,m+1}, \quad \delta_m := \pi_{m,m} \quad \text{and} \quad \varepsilon_m := \pi_{m,m-a}. \tag{2.9}$$

Then the coefficients ω_m , μ_m , and ν_m of (2.5) can be computed from

$$\begin{aligned} \omega_m &= -\frac{\gamma_{m+1}}{\gamma_m} \alpha_{m+1}, \\ \mu_m &= -\frac{1}{\gamma_m} \left[\gamma_m - \delta_{m+1} + \frac{\omega_m(\gamma_{m-1} - \delta_m)}{\alpha_m} - \gamma_{m+1}(1 + \beta_{m+1}) \right], \\ \nu_m &= \frac{\omega_m \varepsilon_{m-1} \beta_{m-a-1}}{\alpha_{m-a-1} \varepsilon_{m-2}}. \end{aligned} \tag{2.10}$$

Proof. Using (2.7) in (2.4) we see that

$$u_m = \frac{1}{\lambda^2} \left[\pi_{m,m-a} t_{m-a} + \sum_{j=m-a+1}^{m+1} (\pi_{m,j} - \pi_{m+1,j}) t_j - \pi_{m+1,m+2} t_{m+2} \right], \quad m \geq a, \tag{2.11}$$

where we have written t_j instead of $t_j(\lambda)$ for short.

Now, let $m \geq a + 2$. Substituting (2.11) in (2.5), and multiplying throughout by λ^2 , we obtain that

$$\begin{aligned} &\pi_{m,m-a} t_{m-a} + \sum_{j=m-a+1}^{m+1} (\pi_{m,j} - \pi_{m+1,j}) t_j - \pi_{m+1,m+2} t_{m+2} \\ &- \omega_m \left[\pi_{m-1,m-1-a} \lambda t_{m-1-a} + \sum_{j=m-a}^m (\pi_{m-1,j} - \pi_{m,j}) \lambda t_j - \pi_{m,m+1} \lambda t_{m+1} \right] \\ &- \mu_m \left[\pi_{m-1,m-1-a} t_{m-1-a} + \sum_{j=m-a}^m (\pi_{m-1,j} - \pi_{m,j}) t_j - \pi_{m,m+1} t_{m+1} \right] \\ &- \nu_m \left[\pi_{m-2,m-2-a} t_{m-2-a} + \sum_{j=m-1-a}^{m-1} (\pi_{m-2,j} - \pi_{m-1,j}) t_j - \pi_{m-1,m} \lambda t_m \right] \\ &= 0. \end{aligned} \tag{2.12}$$

Next, from (2.6) we know that

$$\lambda t_j = -\frac{1}{\alpha_j} t_{j+1} + \frac{1 + \beta_j}{\alpha_j} t_j - \frac{\beta_j}{\alpha_j} t_{j-1}, \quad j = 0, 1, \dots \tag{2.13}$$

Substituting (2.13) in (2.12) yields that

$$\begin{aligned} & \pi_{m,m-a} t_{m-a} + \sum_{j=m-a+1}^{m+1} (\pi_{m,j} - \pi_{m+1,j}) t_j - \pi_{m+1,m+2} t_{m+2} \\ & - \omega_m \left[\pi_{m-1,m-1-a} \left(-\frac{1}{\alpha_{m-1-a}} t_{m-a} + \frac{1 + \beta_{m-1-a}}{\alpha_{m-1-a}} t_{m-1-a} - \frac{\beta_{m-1-a}}{\alpha_{m-1-a}} t_{m-2-a} \right) \right. \\ & + \sum_{j=m-a}^m (\pi_{m-1,j} - \pi_{m,j}) \left(-\frac{1}{\alpha_j} t_{j+1} + \frac{1 + \beta_j}{\alpha_j} t_j - \frac{\beta_j}{\alpha_j} t_{j-1} \right) \\ & \left. - \pi_{m,m+1} \left(-\frac{1}{\alpha_{m+1}} t_{m+2} + \frac{1 + \beta_{m+1}}{\alpha_{m+1}} t_{m+1} - \frac{\beta_{m+1}}{\alpha_{m+1}} t_m \right) \right] \\ & - \mu_m \left[\pi_{m-1,m-1-a} t_{m-1-a} + \sum_{j=m-a}^m (\pi_{m-1,j} - \pi_{m,j}) t_j - \pi_{m,m+1} t_{m+1} \right] \\ & - \nu_m \left[\pi_{m-2,m-2-a} t_{m-2-a} + \sum_{j=m-1-a}^{m-1} (\pi_{m-2,j} - \pi_{m-1,j}) t_j - \pi_{m-1,m} \lambda t_m \right] \\ & = 0, \end{aligned}$$

which is of the form $\sum_{j=m-a-2}^{m+2} \eta_{m,j} t_j = 0$ for some constants $\eta_{m,j}$. Thus, we must have $\eta_{m,j} = 0$ for all $j = m - a - 2, m - a - 1, \dots, m + 2$. Now, from $\eta_{m,m+2} = \eta_{m,m+1} = \eta_{m,m-a-2} = 0$, we obtain the expressions for ω_m , μ_m , and ν_m , respectively, as given in (2.10). \square

3. The algorithm

We now return to the general framework of semi-iterative methods for computing the Drazin-inverse solution of singular linear systems that was discussed in Section 1. We choose the polynomials $p_m(\lambda)$ that appear in (1.3) and (1.4) to be precisely those given in Theorem 2.1, the integer a in the latter being $\text{ind}(A)$. As they are in Π_m^0 , these $p_m(\lambda)$ already satisfy (1.5) and (1.8).

From (1.3), (1.4), and (2.4), the iterates x_m and x_{m+1} of the semi-iterative method satisfy

$$x_{m+1} - x_m = u_m(A) r_0. \tag{3.1}$$

But the $u_m(\lambda)$ satisfy the 3-term recursion relation given in (2.5). Consequently, the x_m satisfy the 4-term recursion relation

$$x_{m+1} = x_m + \omega_m A(x_m - x_{m-1}) + \mu_m(x_m - x_{m-1}) + \nu_m(x_{m-1} - x_{m-2}), \quad m \geq a + 1, \tag{3.2}$$

which is exactly of the form given in [8] for the case $a = 1$. Note that this recursion relation has the same length *independent* of a .

As the recursion relation above is valid for $m \geq a + 1$ and as $v_{a+1} = 0$, we see that in order to start the algorithm we need x_a and x_{a+1} . Now because $p_a(\lambda) = 1$, we have that $q_{a-1}(\lambda) = 0$ so that $x_a = x_0$. As for x_{a+1} , we proceed as follows: First, we know that

$$p_{a+1}(\lambda) = 1 - \rho\lambda^{a+1}$$

for some constant ρ that can be uniquely determined from the characterization property in (2.2). As $\langle p_{a+1}, \lambda \rangle = 0$, we evidently have that

$$\rho = \frac{\langle 1, \lambda \rangle}{\langle 1, \lambda^{a+2} \rangle}. \tag{3.3}$$

Next, on recalling (2.4), we see that

$$u_a(\lambda) = \rho\lambda^a. \tag{3.4}$$

Finally, from (3.1) and (3.4) we have that

$$x_{a+1} = x_a + \rho A^a r_0 = x_a + \rho A^a (b - Ax_0). \tag{3.5}$$

Assuming that the polynomials $t_m(\lambda)$ and the constants α_m and β_m in (2.6) are known, our algorithm now reads as follows:

- Step 0:* Choose x_0 and set $x_a = x_0$.
 Set $\pi_{a,0} = \frac{-1}{t'_a(0)}$, $\pi_{a,1} = -\pi_{a,0}$, and $\pi_{a,j} = 0$ for $j \neq 0, 1$.
 Determine ρ from (3.3).
 Compute x_{a+1} from (3.5).
- Step 1:* For $m = a + 1, a + 2, \dots$, until convergence, do:
 Solve (2.8) for the $\pi_{m,j}$.
 Compute ω_m, μ_m , and v_m from (2.9) and (2.10).
 Compute x_{m+1} from (3.2).

For the special case in which the weight function $w(\lambda)$ is that defined by (1.9), the polynomials $t_m(\lambda)$ and the corresponding constants α_m and β_m are given by

$$t_m(\lambda) = \frac{T_m(z(\lambda))}{T_m(z(0))}, \quad \text{with } z(\lambda) = \frac{c - \lambda}{d}, \tag{3.6}$$

where $T_m(z)$ are the Chebyshev polynomials of the first kind normalized so that $T_m(1) = 1$, and

$$\begin{aligned} \alpha_0 &= \frac{1}{c}, & \beta_0 &= 0, \\ \alpha_1 &= \frac{2c}{2c^2 - d^2}, & \beta_1 &= c\alpha_1 - 1, \\ \alpha_m &= \frac{1}{c - (\frac{1}{2}d)^2\alpha_{m-1}}, & \beta_m &= c\alpha_m - 1, \quad m \geq 2. \end{aligned}$$

In addition, the constant ρ in (3.3) is now given by

$$\rho = \frac{1}{c^{a+1} \sum_{k=0}^{\lfloor a/2 \rfloor + 1} \binom{a+2}{2k} \binom{2k}{k} \left(\frac{d}{2c}\right)^{2k}}.$$

4. Error bounds and convergence analysis

4.1. General preliminaries and error bounds

Let us denote by $\hat{\mathcal{S}}$ the direct sum of the invariant subspaces of A corresponding to its nonzero eigenvalues λ_j , and by $\tilde{\mathcal{S}}$, its invariant subspace corresponding to its zero eigenvalue. Thus, $\hat{\mathcal{S}} = \mathcal{R}(A^a)$, the range of A^a , and $\tilde{\mathcal{S}} = \mathcal{N}(A^a)$, the nullspace of A^a . Every vector in \mathbb{C}^n can be written as the sum of two unique vectors, one in $\hat{\mathcal{S}}$ and the other in $\tilde{\mathcal{S}}$.

Resolve $b = \hat{b} + \tilde{b}$, where $\hat{b} \in \hat{\mathcal{S}}$ and $\tilde{b} \in \tilde{\mathcal{S}}$. Then $A^D b$, the Drazin-inverse solution of $Ax = b$, is the unique vector in $\hat{\mathcal{S}}$ that satisfies the consistent linear system $Ax = \hat{b}$. From (1.3) and (1.4) we see that

$$\begin{aligned} x_m - A^D b &= p_m(A)x_0 + q_{m-1}(A)(\hat{b} + \tilde{b}) - A^D b \\ &= p_m(A)x_0 + q_{m-1}(A)AA^D b + q_{m-1}(A)\tilde{b} - A^D b \\ &= p_m(A)(x_0 - A^D b) + q_{m-1}(A)\tilde{b}. \end{aligned} \tag{4.1}$$

Decompose $x_0 = \hat{x}_0 + \tilde{x}_0$, where $\hat{x}_0 \in \hat{\mathcal{S}}$ and $\tilde{x}_0 \in \tilde{\mathcal{S}}$. Then (4.1) becomes

$$x_m - A^D b = p_m(A)(\hat{x}_0 - A^D b) + p_m(A)\tilde{x}_0 + q_{m-1}(A)\tilde{b}. \tag{4.2}$$

Because

$$p_m(\lambda) = 1 - \lambda^{a+1}u(\lambda) \tag{4.3}$$

for some $u \in \Pi_{m-a-1}$, we have that

$$p_m(A)\tilde{x}_0 = \tilde{x}_0 - u(A)A^{a+1}\tilde{x}_0 = \tilde{x}_0 \tag{4.4}$$

as $\tilde{x}_0 \in \tilde{\mathcal{S}} = \mathcal{N}(A^a)$. Similarly, $q_{m-1}(\lambda) = \lambda^a u(\lambda)$ by (1.4) and (4.3), so that

$$q_{m-1}(A)\tilde{b} = u(A)A^a\tilde{b} = 0 \tag{4.5}$$

as $\tilde{b} \in \tilde{\mathcal{S}} = \mathcal{N}(A^a)$.

Combining (4.4) and (4.5) in (4.2), we deduce the following result.

Theorem 4.1. *Let $x_0 = \hat{x}_0 + \tilde{x}_0$, where $\hat{x}_0 \in \hat{\mathcal{S}}$ and $\tilde{x}_0 \in \tilde{\mathcal{S}}$. Then*

$$x_m - A^D b = p_m(A)(\hat{x}_0 - A^D b) + \tilde{x}_0. \tag{4.6}$$

Now, as the vector $\hat{x}_0 - A^D b$ is in $\hat{\mathcal{S}}$, we observe that the behavior of $x_m - A^D b$ is determined by the action of $p_m(A)$ on $\hat{\mathcal{S}}$.

Recall that, by $\text{ind}(A - \lambda_j I) = k_j$, $\lambda_j \in \sigma(A) \setminus \{0\}$, the fact that $\hat{x}_0 - A^D b \in \hat{\mathcal{S}}$ implies that

$$p_m(A)(\hat{x}_0 - A^D b) = \sum_{\lambda_j \in \sigma(A) \setminus \{0\}} \sum_{i=0}^{k_j-1} r_{ji} p_m^{(i)}(\lambda_j) \tag{4.7}$$

for some vectors r_{ji} that lie in the invariant subspace of A corresponding to λ_j . Thus, from (4.6) and (4.7),

$$\|x_m - A^D b - \tilde{x}_0\| = \|p_m(A)(\hat{x}_0 - A^D b)\| \leq C \left(\max_{\lambda_j \in \sigma(A) \setminus \{0\}} \max_{0 \leq i \leq k_j-1} |p_m^{(i)}(\lambda_j)| \right) \tag{4.8}$$

for some positive constant C . Replacing the maximum over the $\lambda_j \in \sigma(A) \setminus \{0\}$ by the maximum over the interval $[c - d, c + d]$, and using

$$\max_{\lambda \in [c-d, c+d]} |p_m^{(i)}(\lambda)| \leq D_i m^{2i} \left(\max_{\lambda \in [c-d, c+d]} |p_m(\lambda)| \right) \quad \text{for some } D_i > 0,$$

which follows from one of Markoff's inequalities, see. e.g., Meinardus [9, p. 67], (4.8) becomes

$$\|x_m - A^D b - \tilde{x}_0\| = \|p_m(A)(\hat{x}_0 - A^D b)\| \leq M m^{2\hat{k}-2} \left(\max_{\lambda \in [c-d, c+d]} |p_m(\lambda)| \right), \tag{4.9}$$

where M is a positive constant and

$$\hat{k} = \max\{k_j: \lambda_j \in \sigma(A) \setminus \{0\}\}. \tag{4.10}$$

Hence, all we have to analyze is $\max_{\lambda \in [c-d, c+d]} |p_m(\lambda)|$.

Before we go on, we observe from (4.6) and (4.7) that the conditions in (1.7) ensure the convergence of $\{x_m\}_{m=0}^\infty$ to $A^D b + \tilde{x}_0$, as guaranteed also by Eiermann et al. [4, Lemma 2]. Also, if $\tilde{x}_0 = 0$, which can be enforced by picking $x_0 = 0$, then $\lim_{m \rightarrow \infty} x_m = A^D b$ under (1.7).

4.2. Convergence analysis

In the sequel, we analyze the case in which the weight function $w(\lambda)$ is that defined by (1.9). Obviously, we first need to know the behavior of the $\pi_{m,j}$ in Theorem 2.3 for $m \rightarrow \infty$. For this we have to start with the behavior of the $t_m^{(i)}(0)$ for $m \rightarrow \infty$, as is obvious from (2.8). Recall that in this case $t_m(\lambda)$ are as in (3.6).

Lemma 4.2. *Suppose that $t_m(\lambda)$ are the polynomials given in (3.6). If $\lambda \in [c - d, c + d]$, then, for $i = 0, 1, 2, \dots$,*

$$t_m^{(i)}(\lambda) = P_i(\lambda, m) \frac{e^{m \cosh^{-1} z(\lambda)} + e^{-m \cosh^{-1} z(\lambda)}}{e^{m \cosh^{-1} z(0)} + e^{-m \cosh^{-1} z(0)}} + N_i(\lambda, m) \frac{e^{m \cosh^{-1} z(\lambda)} - e^{-m \cosh^{-1} z(\lambda)}}{e^{m \cosh^{-1} z(0)} + e^{-m \cosh^{-1} z(0)}}, \tag{4.11}$$

where $P_i(\lambda, m)$ and $N_i(\lambda, m)$ are polynomials in m , whose coefficients are functions of λ and whose degree is dependent on the parity of i , given by

$$\begin{aligned}
 P_{2r}(\lambda, m) &= \left[-\frac{1}{\sqrt{(c-\lambda)^2 - d^2}} \right]^{2r} m^{2r} + O(m^{2r-2}), \\
 N_{2r}(\lambda, m) &= \frac{(2r-1)2r}{2}(c-\lambda) \left[-\frac{1}{\sqrt{(c-\lambda)^2 - d^2}} \right]^{2r+1} m^{2r-1} + O(m^{2r-3}),
 \end{aligned}
 \tag{4.12}$$

with the terms $O(m^{2r-2})$ and $O(m^{2r-3})$ missing for $r = 0, 1$, and

$$\begin{aligned}
 P_{2r+1}(\lambda, m) &= \frac{2r(2r+1)}{2}(c-\lambda) \left[-\frac{1}{\sqrt{(c-\lambda)^2 - d^2}} \right]^{2r+2} m^{2r} + O(m^{2r-2}), \\
 N_{2r+1}(\lambda, m) &= \left[-\frac{1}{\sqrt{(c-\lambda)^2 - d^2}} \right]^{2r+1} m^{2r+1} + O(m^{2r-1}),
 \end{aligned}
 \tag{4.13}$$

with the terms $O(m^{2r-2})$ and $O(m^{2r-1})$ missing for $r = 0, 1$, and $r = 0$, respectively.

Proof. The proof is straightforward and proceeds by induction on i . \square

Taking $\lambda = 0$, (4.11) becomes

$$t_m^{(i)}(0) = P_i(0, m) + N_i(0, m) \left(1 - \frac{2}{\kappa^{-2m} + 1} \right),
 \tag{4.14}$$

where κ , defined by (1.10), satisfies also

$$\kappa = e^{-\cosh^{-1} z(0)}.$$

Upon substituting (4.12) and (4.13) in (4.14), we now have the following result.

Theorem 4.3. *Suppose that $t_m(\lambda)$ are the polynomials defined in (3.6). Then, for $i = 0, 1, 2, \dots$,*

$$t_m^{(i)}(0) = \begin{cases} 1, & i = 0, \\ -\frac{1}{(c^2 - d^2)^{1/2}} m + O(m\kappa^{2m}), & i = 1, \\ (-1)^i \frac{1}{(c^2 - d^2)^{i/2}} m^i + \sum_{k=0}^{i-1} \eta_{i,k} m^k + O(m^l \kappa^{2m}), & i \geq 2 \end{cases}$$

as $m \rightarrow \infty$, where $\eta_{i,k}$ are some constants and $l = i - 1$ if i is even, and $l = i$ if i is odd.

Theorem 4.3 has the following implication. In order to solve the system in (2.8) for the $\pi_{m,j}$, we first introduce the matrices B and E in $\mathbb{C}^{a+2, a+2}$ and the vector h in \mathbb{C}^{a+2} as follow:

- For $i, j = 0, 1, 2, \dots, a + 1$, set

$$b_{i+1, j+1} = \begin{cases} 1, & i = 0, \\ -\frac{1}{(c^2 - d^2)^{1/2}}(m - a + j), & i = 1, \\ (-1)^i \frac{1}{(c^2 - d^2)^{i/2}}(m - a + j)^i + \sum_{k=0}^{i-1} \eta_{i,k}(m - a + j)^k, & 2 \leq i \leq a + 1, \end{cases}$$

with the $\eta_{i,k}$ as in Theorem 4.3.

- For $j = 0, 1, 2, \dots, a + 1$, set

$$e_{i+1, j+1} = \begin{cases} 0, & i = 0, \\ O(m^l \kappa^{2m}), & 1 \leq i \leq a + 1, \end{cases}$$

where l is defined in Theorem 4.3, and observe that $E \rightarrow O$ as $m \rightarrow \infty$.

- Introduce the vector h via

$$h_{i+1} = \begin{cases} 0, & i = 0, \\ 1, & i = 1, \\ 0, & 2 \leq i \leq a + 1. \end{cases}$$

Then the linear system in (2.8) can be written as

$$(B + E)\pi = h, \tag{4.15}$$

where $\pi \in \mathbb{C}^{a+2}$ is the unknown vector whose $(j + 1)$ th entry, $0 \leq j \leq a + 1$, is $\pi_{m, m-a+j}$.

To solve (4.15) for π we apply elementary row operations to obtain the equivalent system

$$(B^{(2)} + E^{(2)})\pi = h^{(2)},$$

where

$$b_{i+1, j+1}^{(2)} = \begin{cases} 1, & i = 0, \\ (m - a + j)^i, & 1 \leq i \leq a + 1, \end{cases}$$

i.e., $B^{(2)}$ is a Vandermonde matrix, and

$$h_{i+1}^{(2)} = \begin{cases} 0, & i = 0, \\ -(c^2 - d^2)^{1/2}, & i = 1, \\ K_i(c^2 - d^2)^{i/2}, & 2 \leq i \leq a + 1, \end{cases}$$

where K_i is a constant that depends only on the coefficients $\eta_{i,k}$.

Now, using the algorithm to solve Vandermonde systems (see [7, p. 122]), we obtain the equivalent system

$$(I + E^{(3)})\pi = h^{(3)},$$

where

$$h_{i+1}^{(3)} = (-1)^i \binom{a+1}{i} \frac{(c^2 - d^2)^{1/2}}{a!} m^a + O(m^{a-1}) \quad \text{for } 0 \leq i \leq a+1 \tag{4.16}$$

and $E^{(3)} \rightarrow O$, as $m \rightarrow \infty$. Therefore, we have the following result.

Theorem 4.4. *For $i = 0, 1, 2, \dots, a, a+1$, the $\pi_{m,m-a+i}$ in (2.8) are given by $\pi_{m,m-a+i} = h_{i+1}^{(3)}$, where $h_{i+1}^{(3)}$ are defined in (4.16).*

We now combine Theorem 4.4 with the expansion of $p_m(\lambda)$ in Theorem 2.3 to derive an asymptotically optimal upper bound on $|p_m(\lambda)|$ for $\lambda \in [c-d, c+d]$.

Theorem 4.5. *Consider the polynomials $p_m(\lambda)$ of Theorem 2.1 with the weight function $w(\lambda)$ given by (1.9). Then*

$$\max_{\lambda \in [c-d, c+d]} |p_m(\lambda)| = \frac{2(\kappa^{-1} - \kappa)(1 + \kappa)^{a-1}}{a!} m^a \kappa^{m-a+1} + O(m^{a-1} \kappa^m) \quad \text{as } m \rightarrow \infty, \tag{4.17}$$

where κ is given by (1.10).

Proof. For the weight function $w(\lambda)$ given by (1.9) we have that the polynomials $t_m(\lambda)$ are defined by (3.6).

For $\lambda \in [c-d, c+d]$, it is easy to see that

$$t_m(\lambda) = 2\kappa^m \Re(s(\lambda)^m) + O(\kappa^{3m}) \quad \text{as } m \rightarrow \infty, \tag{4.18}$$

where

$$s(\lambda) = e^{i \arccos z(\lambda)} \quad \text{with } |s(\lambda)| = 1$$

and, therefore,

$$\lambda = c - dz(\lambda) = c - d\Re(s(\lambda)). \tag{4.19}$$

Now, from (2.7), (3.6), Theorem 4.4, and (4.18), we obtain that for $m \rightarrow \infty$,

$$\begin{aligned} p_m(\lambda) &= \frac{1}{\lambda} \sum_{j=0}^{a+1} (-1)^j \binom{a+1}{j} \frac{\sqrt{c^2 - d^2}}{a!} m^a 2\kappa^{m-a+j} \Re(s(\lambda)^{m-a+j}) + O(m^{a-1} \kappa^m) \\ &= \frac{2\sqrt{c^2 - d^2}}{a! \lambda} m^a \kappa^{m-a+1} \Re \left(\frac{s(\lambda)^{m-a}}{\kappa} [1 - \kappa s(\lambda)]^{a+1} \right) + O(m^{a-1} \kappa^m). \end{aligned} \tag{4.20}$$

On the other hand, from (1.10) we have that

$$\kappa^{-1} + \kappa = 2\frac{c}{d} \quad \text{and} \quad \kappa^{-1} - \kappa = \frac{2\sqrt{c^2 - d^2}}{d}. \tag{4.21}$$

Now, using (4.19) and the fact that $|s(\lambda)| = 1$, we conclude that

$$\lambda = \frac{1}{2}d\kappa^{-1}[1 - \kappa s(\lambda)][1 - \overline{\kappa s(\lambda)}].$$

Inserting this expression into (4.20) and using (4.21), we obtain that

$$p_m(\lambda) = \frac{2}{a!}(\kappa^{-1} - \kappa)m^a\kappa^{m-a+1}\Re\left(\frac{s(\lambda)^{m-a}[1 - \kappa s(\lambda)]^a}{1 - \overline{\kappa s(\lambda)}}\right) + O(m^{a-1}\kappa^m) \quad \text{as } m \rightarrow \infty.$$

Then

$$\frac{1}{m^a\kappa^{m-a+1}}p_m(\lambda) = \frac{2}{a!}(\kappa^{-1} - \kappa)\Re\left(\frac{s(\lambda)^{m-a}[1 - \kappa s(\lambda)]^a}{1 - \overline{\kappa s(\lambda)}}\right) + O(m^{-1}) \quad \text{as } m \rightarrow \infty. \tag{4.22}$$

Next, using the facts that $|s(\lambda)| = 1$ and $|1 - \kappa s(\lambda)| = |1 - \overline{\kappa s(\lambda)}|$, we conclude that

$$\max_{\lambda \in [c-d, c+d]} \left| \frac{s(\lambda)^{m-a}[1 - \kappa s(\lambda)]^a}{1 - \overline{\kappa s(\lambda)}} \right| = \max_{\lambda \in [c-d, c+d]} |1 - \kappa s(\lambda)|^{a-1} = (1 + \kappa)^{a-1} \tag{4.23}$$

and this maximum is attained at $\lambda = c + d$ for which $s(\lambda) = -1$. Finally, combining (4.23) and (4.22), we obtain that

$$\frac{1}{m^a\kappa^{m-a+1}} \max_{\lambda \in [c-d, c+d]} |p_m(\lambda)| = \frac{2(\kappa^{-1} - \kappa)(1 + \kappa)^{a-1}}{a!} + O(m^{-1}) \quad \text{as } m \rightarrow \infty,$$

from which (4.17) follows. \square

Now, using (4.9) and (4.17), we have the following convergence result that is the main result of this section.

Corollary 4.6. *With the same notation as in Theorem 4.5, we have*

$$\|x_m - A^D b - \tilde{x}_0\| = O(m^{a+2\hat{k}-2}\kappa^m) \quad \text{as } m \rightarrow \infty,$$

where \hat{k} is as defined in (4.10).

Theorem 4.5 implies that

$$\max_{\lambda \in [c-d, c+d]} |p_m(\lambda)| \approx \frac{2(\kappa^{-1} - \kappa)(1 + \kappa)^{a-1}}{a!} m^a \kappa^{m-a+1}.$$

On the other hand, the Bernstein result as applied by Eiermann and Starke to the polynomials $\{p_m\}$ developed in their paper, see [6, p. 314], gives that their residual polynomials satisfy that

$$\max_{\lambda \in [c-d, c+d]} |p_m(\lambda)| \approx \frac{2(\kappa^{-1} - \kappa)^a}{a!} m^a \kappa^m. \tag{4.24}$$

Now,

$$\frac{2(\kappa^{-1} - \kappa)^a}{a!} m^a \kappa^m < \frac{2(\kappa^{-1} - \kappa)(1 + \kappa)^{a-1}}{a!} m^a \kappa^{m-a+1},$$

because $1 > -\kappa$. Therefore, our polynomials are not “near-optimal”. However, the residual polynomials $\{p_m\}$ constructed by Eiermann and Starke in [6] cannot be computed by means of short recurrences as we have developed for the present residuals in Section 3. Such short recurrences make for the efficient implementation of semi-iterative methods. In this regard please see also the comments on Hanke and Hochbruck [8, pp. 90, 93].

5. Numerical examples

In this section we use the algorithm developed in Section 3 to compute the eigenprojection $Z_A := I - AA^D$ onto the eigenspace of A corresponding to the eigenvalue 0 of three singular matrices whose index exceeds 1.

If we take $b = 0$ in (1.1) then Corollary 4.6 implies

$$\lim_{m \rightarrow \infty} x_m = \tilde{x}_0 = (I - AA^D)x_0.$$

Now, if we choose x_0 as the i th column of I , the above expression represents the i th column of the eigenprojection Z_A .

First, consider the following singular M -matrix:

$$A_1 = \begin{bmatrix} 1 & -1 & 0 & 0 & 0 & 0 \\ -1 & 1 & 0 & 0 & 0 & 0 \\ -1 & -1 & 1 & -1 & 0 & 0 \\ -1 & -1 & -1 & 1 & 0 & 0 \\ -1 & -1 & -1 & 0 & 2 & -1 \\ -1 & -1 & 0 & -1 & -1 & 2 \end{bmatrix}.$$

Observe that $\sigma(A_1) = \{0, 0, 1, 2, 2, 3\}$ and $a = 2$. So we can choose $c = 2$ and $d = 1$ in (1.2). Using the algorithm of Section 3, with the polynomials $t_m(\lambda)$ defined by (3.6), and stopping when

$$\frac{\|x_{m+1} - x_m\|_\infty}{\|x_m\|_\infty} \leq 10^{-15},$$

we obtain, after 35 iterations, that

$$Z_{A_1} = \begin{bmatrix} 0.500000000000001 & 0.499999999999999 & 0 & 0 & 0 & 0 \\ 0.499999999999999 & 0.500000000000001 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0.500000000000001 & 0.499999999999999 & 0 & 0 \\ 0 & 0 & 0.499999999999999 & 0.500000000000001 & 0 & 0 \\ 0 & 0 & 0.500000000000005 & 0.499999999999997 & -0.000000000000004 & 0.000000000000002 \\ 0 & 0 & 0.499999999999997 & 0.500000000000005 & 0.000000000000002 & -0.000000000000004 \end{bmatrix}.$$

The exact eigenprojection is given by

$$\begin{bmatrix} 0.5 & 0.5 & 0 & 0 & 0 & 0 \\ 0.5 & 0.5 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0.5 & 0.5 & 0 & 0 \\ 0 & 0 & 0.5 & 0.5 & 0 & 0 \\ 0 & 0 & 0.5 & 0.5 & 0 & 0 \\ 0 & 0 & 0.5 & 0.5 & 0 & 0 \end{bmatrix}.$$

As a second example, we again consider a singular M -matrix, this time of index $a=4$:

$$A_2 = \begin{bmatrix} 1.0 & -1.0 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1.0 & 1.0 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1.0 & -1.0 & 1.0 & -1.0 & 0 & 0 & 0 & 0 \\ -1.0 & -1.0 & -1.0 & 1.0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1.0 & -1.0 & -1.0 & -1.0 \\ 0 & 0 & 0 & 0 & -1.0 & 1.0 & -1.0 & -1.0 \\ 0 & 0 & 0 & -1.0 & 0 & 0 & 1.0 & -1.0 \\ 0 & 0 & 0 & 0 & 0 & 0 & -1.0 & 1.0 \end{bmatrix}.$$

Here $\sigma(A_2) = \{0, 0, 0, 0, 2, 2, 2, 2\}$. With $c=2$ and $d=1$ in (1.2) we get using the algorithm in Section 3 that after 25 iterations for columns 1, 2, 5, 6, and 7 and 45 iterations for columns 3 and 4 that

$$Z_{A_2} = \begin{bmatrix} 5.0000 \times 10^{-1} & 5.0000 \times 10^{-1} & 0 & 0 & 0 & 0 & 0 & 0 \\ 5.0000 \times 10^{-1} & 5.0000 \times 10^{-1} & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 5.0000 \times 10^{-1} & 5.0000 \times 10^{-1} & 0 & 0 & 0 & 0 \\ 0 & 0 & 5.0000 \times 10^{-1} & 5.0000 \times 10^{-1} & 0 & 0 & 0 & 0 \\ 0 & 0 & 1.2500 \times 10^{-1} & -1.2500 \times 10^{-1} & 5.0000 \times 10^{-1} & 5.0000 \times 10^{-1} & 0 & 0 \\ 0 & 0 & 1.2500 \times 10^{-1} & -1.2500 \times 10^{-1} & 5.0000 \times 10^{-1} & 5.0000 \times 10^{-1} & 0 & 0 \\ -1.2500 \times 10^{-1} & -1.2500 \times 10^{-1} & -5.7858 \times 10^{-12} & 2.5000 \times 10^{-1} & 0 & 0 & 5.0000 \times 10^{-1} & 5.0000 \times 10^{-1} \\ 1.2500 \times 10^{-1} & 1.2500 \times 10^{-1} & -2.5000 \times 10^{-1} & 5.3423 \times 10^{-11} & 0 & 0 & 5.0000 \times 10^{-1} & 5.0000 \times 10^{-1} \end{bmatrix}.$$

The exact eigenprojection is given here by

$$\begin{bmatrix} 0.5000 & 0.5000 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0.5000 & 0.5000 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0.5000 & 0.5000 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0.5000 & 0.5000 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0.1250 & -0.1250 & 0.5000 & 0.5000 & 0 & 0 \\ 0 & 0 & 0.1250 & -0.1250 & 0.5000 & 0.5000 & 0 & 0 \\ -0.1250 & -0.1250 & 0 & 0.2500 & 0 & 0 & 0.5000 & 0.5000 \\ 0.1250 & 0.1250 & -0.2500 & 0 & 0 & 0 & 0.5000 & 0.5000 \end{bmatrix}.$$

Finally, we consider a singular matrix A with $a = 3$.

$$A_3 = \begin{bmatrix} 5 & -1 & -1 & -1 & -1 & 0 & -1 \\ 1 & 3 & -1 & -1 & -1 & 0 & -1 \\ 0 & 0 & 3 & -1 & -1 & 0 & -1 \\ 0 & 0 & 1 & 1 & -1 & 0 & -1 \\ 0 & 0 & 0 & 0 & 1 & 0 & -1 \\ 0 & 0 & 0 & 0 & 1 & 0 & -1 \\ 0 & 0 & 0 & 0 & 0 & 1 & -1 \end{bmatrix}.$$

Here $\sigma(A_3) = \{0, 0, 0, 2, 2, 4, 4\}$, so we can take $c = 3$ and $d = 1$. Then, using the algorithm in Section 3 we get, after 51 iterations for columns 1, 2, 3 and 4, 29 iterations for column 5, and 6 iterations for columns 6 and 7 that

$$Z_{A_3} = \begin{bmatrix} -0.3908 \times 10^{-12} & 0.1799 \times 10^{-12} & 0.0267 \times 10^{-12} & 0.1909 \times 10^{-12} & 1.0000 & 0 & 0 \\ -0.1846 \times 10^{-12} & -0.0252 \times 10^{-12} & 0.0270 \times 10^{-12} & 0.1909 \times 10^{-12} & 1.0000 & 0 & 0 \\ 0 & 0 & -0.1840 \times 10^{-12} & 0.1909 \times 10^{-12} & 1.0000 & 0 & 0 \\ 0 & 0 & -0.1890 \times 10^{-12} & 0.1970 \times 10^{-12} & 1.0000 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1.0000 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1.0000 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1.0000 \end{bmatrix}.$$

The exact value is

$$\begin{bmatrix} 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}.$$

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