



Geodesic normal forms and Hecke algebras for the complex reflection groups $G(de, e, n)$ [☆]



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ABSTRACT

We establish geodesic normal forms for the general series of complex reflection groups $G(de, e, n)$ by using the presentations of Corran–Picantin and Corran–Lee–Lee of $G(e, e, n)$ and $G(de, e, n)$ for $d > 1$, respectively. This requires the elaboration of a combinatorial technique in order to explicitly determine minimal word representatives of the elements of $G(de, e, n)$. Using these geodesic normal forms, we construct natural bases for the Hecke algebras associated with the complex reflection groups $G(e, e, n)$ and $G(d, 1, n)$. As an application, we obtain a new proof of the BMR freeness conjecture for these groups.

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1. Introduction

Complex reflection groups are finite groups generated by complex reflections. Recall that a complex reflection is a linear transformation of finite order that fixes a hyperplane pointwise. These groups include finite real reflection groups, also known as finite Coxeter groups. It is well known that every complex reflection group is a direct product of irreducible ones. The irreducible complex reflection groups have been classified by Shephard and Todd [20] in 1954. The classification includes the general 3-parameter series $G(de, e, n)$ that can be easily described in terms of monomial matrices and 34 exceptional groups denoted by G_4, G_5, \dots, G_{37} .

Broué, Malle and Rouquier [6] managed to attach a complex braid group to each complex reflection group. This generalizes the notion of Artin groups attached to finite Coxeter groups. Extending earlier results in [5], they also managed to generalize the definition of the (Iwahori–)Hecke algebra for real reflection groups to arbitrary complex reflection groups by using their definition of the complex braid group. Actually, the Hecke algebra is defined as a quotient of the complex braid group algebra by some polynomial relations. It is believed that nice properties of these objects in the case of real reflection groups could be extended to the general case of complex reflection groups. In [6], it was stated a number of important conjectures about the complex braid groups and the Hecke algebras.

One of these interesting conjectures is the so-called “BMR freeness conjecture”. It states that the Hecke algebra is a free module of rank equal to the order of the associated complex reflection group. This property is valid for the (Iwahori–)Hecke algebra attached to any finite real reflection group (see [3]), where a basis is constructed from geodesic normal forms in the finite Coxeter group due to Matsumoto’s property (see [16]). The BMR freeness conjecture can be easily reduced to the case of irreducible complex reflection groups. During the past two decades, a proof of this conjecture for each case of the classification of Shephard and Todd has been established involving the results of a number of authors. As we are interested in the case of the general series of complex reflection groups, we mention that this conjecture has been established for $G(d, 1, n)$ (see Ariki–Koike [2] and Bremke–Malle [4]) and for $G(de, e, n)$ by Ariki (see [1] and Appendix A.2 of [19]). A list of references for the proof of the BMR freeness conjecture can be found in the next section.

An important constraint in the proof of this conjecture, and in the theory of Hecke algebras for complex reflection groups in general, is the failure of an analogue of Matsumoto’s property. That is, we are not able to easily establish a canonical basis of the Hecke algebra. The proof of the BMR freeness conjecture was obtained by sometimes tedious and lengthy computations in order to explicitly construct bases for the Hecke algebras. It is then of importance to find nice bases for these algebras. In this paper, we construct bases for the Hecke algebras attached to the complex reflection groups $G(e, e, n)$ and $G(d, 1, n)$. We also establish that these bases never coincide with the Ariki basis [1] for the case of $G(e, e, n)$ and with the Ariki–Koike basis [2] for $G(d, 1, n)$.

In order to establish these bases, our attention is firstly shifted to the complex reflection groups $G(de, e, n)$. We construct geodesic normal forms for these groups by using the presentations of Corran–Picantin [10] and Corran–Lee–Lee [9] of $G(e, e, n)$ and $G(de, e, n)$ for $d > 1$, respectively. The geodesic normal forms are easy to describe. They generalize our construction in [18] for $G(e, e, n)$ to all the cases of the general series of complex reflection groups.

We establish that these geodesic normal forms provide bases for the Hecke algebras attached to $G(e, e, n)$ and $G(d, 1, n)$. Since these bases are constructed from geodesic normal forms in the complex reflection groups, they are natural bases for the associated Hecke algebras. Note that the geodesic normal forms for $G(e, e, n)$ have been already used in our previous work [18] in order to construct intervals in $G(e, e, n)$ that give rise to nice structures (called interval Garside structures) for the associated complex braid groups.

The article is organized as follows. In Section 2, we provide a basic background material and recall the BMR freeness conjecture. The geodesic normal forms for the complex reflection groups $G(de, e, n)$ are constructed in Section 3, which sets the stage for our later work. The techniques used are elementary and

the associated combinatorial characterizations are very explicit. In Section 4, the attention shifts to the Hecke algebras. Actually, we establish presentations for the Hecke algebras attached to $G(de, e, n)$, by using the presentations of Corran–Picantin and Corran–Lee–Lee of the associated complex braid groups that we recall in the same section. The remaining part of the article establishes that the geodesic normal forms constructed in Section 3 provide natural bases for the Hecke algebras associated with the groups $G(e, e, n)$ and $G(d, 1, n)$.

2. Definitions and preliminaries

Let W be a finite subgroup of $GL_n(\mathbb{C})$ ($n \geq 1$). A complex reflection s of W is an element of finite order $d \geq 2$ such that $\text{Ker}(s - 1)$ is a hyperplane. Let \mathcal{R} be the set of complex reflections of W . We say that W is a complex reflection group if it is generated by \mathcal{R} . Let $\mathcal{A} := \{\text{Ker}(s - 1) \text{ s.t. } s \in \mathcal{R}\}$ be the hyperplane arrangement and $X := \mathbb{C}^n \setminus \bigcup \mathcal{A}$ be the hyperplane complement. The complex reflection group W acts naturally on X . Let X/W be its space of orbits. The complex braid group B attached to W is defined as follows. For details about this definition, we refer to [6].

Definition 2.1. The complex braid group attached to W is the fundamental group

$$B := \pi_1(X/W).$$

Recall that a complex reflection $s \in W$ is called distinguished if its only nontrivial eigenvalue is $\exp(2i\pi/o(s))$, where $o(s)$ denotes the order of s . For the standard notion of braided reflections that we use in the next definition, the reader may check [6]. We are ready to define the Hecke algebra associated with W (see [6] and [14]).

Definition 2.2. Let $R = \mathbb{Z}[a_{s,i}, a_{s,0}^{-1}, 0 \leq i \leq o(s) - 1]$, where s runs over the distinguished reflections, with the convention $a_{s,i} = a_{s',i}$ if s and s' are conjugates in W . The Hecke algebra $H(W)$ attached to the complex reflection group W is the quotient of the complex braid group algebra RB by the relations

$$\sigma^{o(s)} - a_{s,o(s)-1}\sigma^{o(s)-1} - \dots - a_{s,0} = 0,$$

for each braided reflection σ associated with s .

Note that it is enough to choose one such relation per conjugacy class of distinguished reflections, as all the corresponding braided reflections are conjugates in B (see [6]).

The BMR freeness conjecture proposed by Broué, Malle and Rouquier [6] in 1998 states that the Hecke algebra $H(W)$ attached to W is a free R -module of rank equal to the order of W . After two decades, the BMR freeness conjecture is proven through the results of a number of authors. Thus, we have the following theorem.

Theorem 2.3. The Hecke algebra $H(W)$ is a free R -module of rank $|W|$.

The BMR freeness conjecture can be easily reduced to the case where W is irreducible. It is true for the (Iwahori–)Hecke algebra attached to any finite Coxeter group (see Lemma 4.4.3 of [11]). Ariki and Koike [2] proved it for the case of $G(d, 1, n)$. Note that a basis for the Hecke algebra associated with $G(d, 1, n)$ is also given in [4]. Ariki defined in [1] a Hecke algebra for $G(de, e, n)$ by a presentation with generators and relations. He also proved that it is a free module of rank $|G(de, e, n)|$. The Hecke algebra defined by Ariki is isomorphic to the Hecke algebra defined by Broué, Malle, and Rouquier in [6] for $G(de, e, n)$. The details

why this is true can be found in Appendix A.2 of [19]. Hence one gets a proof of Theorem 2.3 for the general series of complex reflection groups.

Concerning the exceptional complex reflection groups, Marin proved the conjecture for G_4 , G_{25} , G_{26} , and G_{32} in [12] and [13]. Marin and Pfeiffer proved it for G_{12} , G_{22} , G_{24} , G_{27} , G_{29} , G_{31} , G_{33} , and G_{34} in [15]. In her PhD thesis and in the article that followed (see [7] and [8]), Chavli proved the validity of this conjecture for G_5 , G_6 , \dots , G_{16} . Recently, Marin proved the conjecture for G_{20} and G_{21} (see [14]) and finally Tsushioka for G_{17} , G_{18} and G_{19} (see [21]). Hence we obtain a proof of Theorem 2.3 for all the cases of irreducible complex reflection groups.

3. Geodesic normal forms for $G(de, e, n)$

This section sets the stage for our later work by establishing a set of geodesic normal forms for the complex reflection groups $G(de, e, n)$, using the generating sets introduced by Corran–Picantin [10] and Corran–Lee–Lee [9] for $G(e, e, n)$ and $G(de, e, n)$ for $d > 1$, respectively. The case of $G(e, e, n)$ has been already done in our previous work (see Section 3 of [18]). We generalize the combinatorial techniques used there to the case of $G(de, e, n)$ for $d > 1$. We obtain natural and explicit geodesic normal forms for these groups.

3.1. Presentations for $G(de, e, n)$

The complex reflection group $G(de, e, n)$ is the group of monomial matrices whose nonzero entries are de -th roots of unity and their product is a d -th root of unity.

Set $d = 1$ and let $e \geq 1$ and $n \geq 2$. Corran–Picantin discovered in [10] a presentation of the complex reflection group $G(e, e, n)$ that is defined as follows.

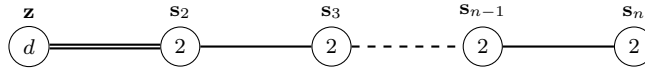
Definition 3.1. The complex reflection group $G(e, e, n)$ is defined by a presentation with generators: $\{\mathbf{t}_i \mid i \in \mathbb{Z}/e\mathbb{Z}\} \cup \{\mathbf{s}_3, \mathbf{s}_4, \dots, \mathbf{s}_n\}$ and relations:

1. $\mathbf{t}_i \mathbf{t}_{i-1} = \mathbf{t}_j \mathbf{t}_{j-1}$ for $i, j \in \mathbb{Z}/e\mathbb{Z}$,
2. $\mathbf{t}_i \mathbf{s}_3 \mathbf{t}_i = \mathbf{s}_3 \mathbf{t}_i \mathbf{s}_3$ for $i \in \mathbb{Z}/e\mathbb{Z}$,
3. $\mathbf{s}_j \mathbf{t}_i = \mathbf{t}_i \mathbf{s}_j$ for $i \in \mathbb{Z}/e\mathbb{Z}$ and $4 \leq j \leq n$,
4. $\mathbf{s}_i \mathbf{s}_{i+1} \mathbf{s}_i = \mathbf{s}_{i+1} \mathbf{s}_i \mathbf{s}_{i+1}$ for $3 \leq i \leq n-1$,
5. $\mathbf{s}_i \mathbf{s}_j = \mathbf{s}_j \mathbf{s}_i$ for $|i-j| > 1$,
6. $\mathbf{t}_i^2 = 1$ for $i \in \mathbb{Z}/e\mathbb{Z}$ and $\mathbf{s}_j^2 = 1$ for $3 \leq j \leq n$.

The matrices in $G(e, e, n)$ that correspond to the generators of this presentation are given by $\mathbf{t}_i \mapsto t_i := \begin{pmatrix} 0 & \zeta_e^{-i} & 0 \\ \zeta_e^i & 0 & 0 \\ 0 & 0 & I_{n-2} \end{pmatrix}$ for $0 \leq i \leq e-1$, and $\mathbf{s}_j \mapsto s_j := \begin{pmatrix} I_{j-2} & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & I_{n-j} \end{pmatrix}$ for $3 \leq j \leq n$. To avoid confusion, we use regular letters for matrices and bold letters for words over the generating set of the presentation of Corran–Picantin.

Remark 3.2.

1. For $e = 1$ and $n \geq 2$, we obtain the classical presentation of the symmetric group S_n .
2. For $e = 2$ and $n \geq 2$, we obtain the classical presentation of the Coxeter group of type D_n .

Fig. 1. Diagram for the presentation of $G(d, 1, n)$.

The attention shifts now to the case $d > 1$. Let $d > 1$, $e \geq 1$ and $n \geq 2$. There exists a presentation of the complex reflection group $G(de, e, n)$ discovered by Corran–Lee–Lee in [9].

Definition 3.3. The complex reflection group $G(de, e, n)$ is defined by a presentation with set of generators: $\mathbf{X} = \{\mathbf{z}\} \cup \{\mathbf{t}_i \mid i \in \mathbb{Z}/de\mathbb{Z}\} \cup \{\mathbf{s}_3, \mathbf{s}_4, \dots, \mathbf{s}_n\}$ and relations as follows.

1. $\mathbf{z}\mathbf{t}_i = \mathbf{t}_{i-e}\mathbf{z}$ for $i \in \mathbb{Z}/de\mathbb{Z}$,
2. $\mathbf{z}\mathbf{s}_j = \mathbf{s}_j\mathbf{z}$ for $3 \leq j \leq n$,
3. Relations 1 to 5 of Definition 3.1 by replacing e by de ,
4. $\mathbf{z}^d = 1$, $\mathbf{t}_i^2 = 1$ for $i \in \mathbb{Z}/de\mathbb{Z}$, and $\mathbf{s}_j^2 = 1$ for $3 \leq j \leq n$.

The generators of this presentation correspond to the following $n \times n$ matrices. The generator \mathbf{t}_i is represented by the matrix $t_i = \begin{pmatrix} 0 & \zeta_{de}^{-i} & 0 \\ \zeta_{de}^i & 0 & 0 \\ 0 & 0 & I_{n-2} \end{pmatrix}$ for $i \in \mathbb{Z}/de\mathbb{Z}$, \mathbf{z} by the diagonal matrix $z = \text{Diag}(\zeta_d, 1, \dots, 1)$ where $\zeta_d = \exp(2i\pi/d)$, and \mathbf{s}_j by the transposition matrix $s_j = (j-1, j)$ for $3 \leq j \leq n$. To avoid confusion, we use regular letters for matrices and bold letters for words over \mathbf{X} . Denote by X the set $\{z, t_0, t_1, \dots, t_{de-1}, s_3, \dots, s_n\}$.

Proposition 3.4. Let $e = 1$. The presentation given in Definition 3.3 is equivalent to the classical presentation of the complex reflection group $G(d, 1, n)$ that can be described by the following well-known diagram (see [6]).

Proof. Let $e = 1$. Relation 1 of Definition 3.3 becomes $\mathbf{z}\mathbf{t}_1 = \mathbf{t}_0\mathbf{z}$, that is $\mathbf{t}_1 = \mathbf{z}^{-1}\mathbf{t}_0\mathbf{z}$. Also by Relation 3 of Definition 3.3, we have $\mathbf{t}_k = \mathbf{z}^{-k}\mathbf{t}_0\mathbf{z}^k$ for $1 \leq k \leq d-1$. If we remove $\mathbf{t}_1, \dots, \mathbf{t}_{d-1}$ from the set of generators and replace every occurrence of \mathbf{t}_k in the defining relations with $\mathbf{z}^{-k}\mathbf{t}_0\mathbf{z}^k$ for $1 \leq k \leq d-1$, we recover the classical presentation of the complex reflection group $G(d, 1, n)$. Note that we replace \mathbf{t}_0 by \mathbf{s}_2 in the set of generators of this presentation. \square

Remark 3.5. For $d = 2$, the presentation described by the diagram of Fig. 1 corresponds to the presentation of the Coxeter group of type B_n .

From now on, we set the following convention.

Convention 3.6. A decreasing-index expression of the form $\mathbf{s}_i\mathbf{s}_{i-1} \dots \mathbf{s}_{i'}$ is the empty word when $i < i'$ and an increasing-index expression of the form $\mathbf{s}_i\mathbf{s}_{i+1} \dots \mathbf{s}_{i'}$ is the empty word when $i > i'$. Similarly, in $G(de, e, n)$, a decreasing-index product of the form $s_i s_{i-1} \dots s_{i'}$ is equal to I_n when $i < i'$ and an increasing-index product of the form $s_i s_{i+1} \dots s_{i'}$ is equal to I_n when $i > i'$, where I_n is the identity $n \times n$ matrix. We also set that \mathbf{z}^0 is the empty word.

3.2. Minimal word representatives

Consider the complex reflection group $G(de, e, n)$ with $d > 1$, $e > 1$ and $n \geq 2$. The case of $G(d, 1, n)$ will be established in the next subsection. Recall that \mathbf{X} denotes the set of the generators of the presentation of Corran–Lee–Lee of $G(de, e, n)$ and X the set of the corresponding matrices.

Denote by $\ell(\mathbf{w})$ the word length over \mathbf{X} of the word $\mathbf{w} \in \mathbf{X}^*$. Let us recall the following definition.

Definition 3.7. Let w be an element of $G(de, e, n)$. We define $\ell(w)$ to be the minimal word length $\ell(\mathbf{w})$ of a word \mathbf{w} over \mathbf{X} that represents w . A reduced expression of w is any word representative of w of word length $\ell(w)$.

Our aim is to represent each element of $G(de, e, n)$ by a reduced word over \mathbf{X} . This requires the elaboration of a combinatorial technique in order to determine a reduced expression decomposition over \mathbf{X} of an element of $G(de, e, n)$.

We introduce Algorithm 1 below that produces a word $RE(w)$ over \mathbf{X} for a given element w in $G(de, e, n)$. This algorithm generalizes the one introduced in Section 3 of our previous work [18] that corresponds to the case of $G(e, e, n)$. Note that we use Convention 3.6 in the elaboration of the algorithm. Later on, we prove that its output $RE(w)$ is a reduced expression over \mathbf{X} of $w \in G(de, e, n)$.

Let $w_n := w \in G(de, e, n)$. For i from n to 2, the i -th step of Algorithm 1 transforms the block diagonal matrix $\left(\begin{array}{c|c} w_i & 0 \\ \hline 0 & I_{n-i} \end{array} \right)$ into a block diagonal matrix $\left(\begin{array}{c|c} w_{i-1} & 0 \\ \hline 0 & I_{n-i+1} \end{array} \right) \in G(de, e, n)$. Actually, for $2 \leq i \leq n$, there exists a unique c with $1 \leq c \leq n$ such that $w_i[i, c] \neq 0$. At each step i of Algorithm 1, if $w_i[i, c] = 1$, we shift it into the diagonal position $[i, i]$ by right multiplication by transpositions. If $w_i[i, c] \neq 1$, we shift it into the first column by right multiplication by transpositions, transform it into 1 by right multiplication by an element of $\{t_0, t_1, \dots, t_{de-1}\}$, and then shift the 1 obtained into the diagonal position $[i, i]$. Finally, we get $w_1 = \zeta_d^k$ for some $0 \leq k \leq d-1$, where ζ_d^k is equal to the product of the nonzero entries of w . By multiplying $\left(\begin{array}{c|c} w_1 & 0 \\ \hline 0 & I_{n-1} \end{array} \right)$ on the right by z^{-k} , we get the identity matrix I_n .

Example 3.8. We apply Algorithm 1 to $w := \begin{pmatrix} \zeta_9 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & \zeta_9 \\ 0 & \zeta_9 & 0 & 0 \end{pmatrix} \in G(9, 3, 4)$.

Step 1 ($i = 4, k = 0, c = 1$): $w' := ws_2 = \begin{pmatrix} 0 & \zeta_9 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & \zeta_9 \\ \zeta_9 & 0 & 0 & 0 \end{pmatrix}$, then $w' := w't_1 = \begin{pmatrix} \zeta_9^2 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & \zeta_9 \\ 0 & 1 & 0 & 0 \end{pmatrix}$, then

$$w' := w's_3s_4 = \begin{pmatrix} \zeta_9^2 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & \boxed{\zeta_9} & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}.$$

Step 2 ($i = 3, k = 1, c = 3$): $w' := w's_3s_2 = \begin{pmatrix} 0 & \zeta_9^2 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ \zeta_9 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$, then $w' := w't_1 = \begin{pmatrix} \zeta_9^3 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$, then

$$w' := w's_3 = \begin{pmatrix} \zeta_9^3 & 0 & 0 & 0 \\ 0 & \boxed{1} & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}.$$

Step 3 ($i = 2, k = 0, c = 2$): $w' = \begin{pmatrix} \boxed{\zeta_9^3} & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} = \begin{pmatrix} \boxed{\zeta_3} & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$.

Step 4 ($k = 1$): $w' := w'z^{-1} = I_4$.

Hence $RE(w) = \mathbf{zs}_3\mathbf{t}_1\mathbf{s}_2\mathbf{s}_3\mathbf{s}_4\mathbf{s}_3\mathbf{t}_1\mathbf{s}_2 = \mathbf{zs}_3\mathbf{t}_1\mathbf{t}_0\mathbf{s}_3\mathbf{s}_4\mathbf{s}_3\mathbf{t}_1\mathbf{t}_0$ (since $\mathbf{s}_2 = \mathbf{t}_0$).

Input : w , a matrix in $G(de, e, n)$ with $d > 1$, $e > 1$, and $n \geq 2$.
Output: $RE(w)$, a word over \mathbf{X} .

Local variables: w' , $RE(w)$, i , U , V , c , k .

Initialisation: $U := [1, \zeta_{de}, \zeta_{de}^2, \dots, \zeta_{de}^{e-1}]$ ($U[1] = 1$), $V := [1, \zeta_d, \zeta_d^2, \dots, \zeta_d^{d-1}]$ ($V[1] = 1$),
 $s_2 := t_0$, $s_2 := \mathbf{t}_0$, $RE(w) := \varepsilon$: the empty word, $w' := w$.

```

for  $i$  from  $n$  down to 2 do
   $c := 1$ ;  $k := 0$ ;
  while  $w'[i, c] = 0$  do
     $c := c + 1$ ;
  end
  #Then  $w'[i, c]$  is the root of unity on the row  $i$ ;
  while  $U[k + 1] \neq w'[i, c]$  do
     $k := k + 1$ ;
  end
  #Then  $w'[i, c] = \zeta_{de}^k$ .

  if  $k \neq 0$  then
     $w' := w' s_c s_{c-1} \dots s_3 s_2 t_k$ ; #Then  $w'[i, 2] = 1$ ;
     $RE(w) := \mathbf{t}_k s_2 s_3 \dots s_c RE(w)$ ;
     $c := 2$ ;
  end
   $w' := w' s_{c+1} \dots s_{i-1} s_i$ ; #Then  $w'[i, i] = 1$ ;
   $RE(w) := \mathbf{s}_i s_{i-1} \dots s_{c+1} RE(w)$ ;
end
 $k := 0$ ;
while  $V[k + 1] \neq w'[1, 1]$  do
   $k := k + 1$ ;
end
#Then  $w'[1, 1] = \zeta_d^k$ ;
 $w' := w' z^{-k}$ ; #Then  $w' = I_n$ ;
if  $k \neq 0$  then
   $RE(w) = \mathbf{z}^k RE(w)$ ;
end
Return  $RE(w)$ ;

```

Algorithm 1: A word over \mathbf{X} corresponding to a matrix $w \in G(de, e, n)$.

The next lemma follows immediately from Algorithm 1. It explains how we can easily obtain the blocks defined in the algorithm.

Lemma 3.9. For $2 \leq i \leq n$, suppose $w_i[i, c] \neq 0$. The block w_{i-1} is obtained by

- removing the row i and the column c from w_i , then by
- multiplying the first column of the new matrix by $w_i[i, c]$.

Moreover, if we denote by a_i the unique nonzero entry on the row i of w , we have $w_1 = \prod_{i=1}^n a_i = \zeta_d^k$ for $0 \leq k \leq d - 1$.

Example 3.10. Let w be as in Example 3.8, where $n = 4$. The block w_3 is obtained by removing the row number 4 and the second column from $w_4 = w$, to obtain $\begin{pmatrix} \zeta_9 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & \zeta_9 \end{pmatrix}$, then by multiplying the first column of this matrix by ζ_9 . The same can be said for the other block w_2 . Finally, the block w_1 is equal to ζ_3 which corresponds to the product of the nonzero entries of w .

Definition 3.11. Let $1 \leq i \leq n$. Let $w_i[i, c] \neq 0$ for $1 \leq c \leq i$.

- If $w_1 = \zeta_d^k$ with $0 \leq k \leq d-1$, we define $RE_1(w)$ to be the word \mathbf{z}^k .
- If $w_i[i, c] = 1$, we define $RE_i(w)$ to be the word $\mathbf{s}_i \mathbf{s}_{i-1} \cdots \mathbf{s}_{c+1}$ (decreasing-index expression).
- If $w_i[i, c] = \zeta_{de}^k$ with $k \neq 0$, we define $RE_i(w)$ to be the word

$$\begin{array}{ll} \mathbf{s}_i \cdots \mathbf{s}_3 \mathbf{t}_k & \text{if } c = 1, \\ \mathbf{s}_i \cdots \mathbf{s}_3 \mathbf{t}_k \mathbf{t}_0 & \text{if } c = 2, \\ \mathbf{s}_i \cdots \mathbf{s}_3 \mathbf{t}_k \mathbf{t}_0 \mathbf{s}_3 \cdots \mathbf{s}_c & \text{if } c \geq 3. \end{array}$$

The output $RE(w)$ of Algorithm 1 is a concatenation of the words $RE_1(w)$, $RE_2(w)$, \dots , and $RE_n(w)$ obtained at each step i from n to 1 of Algorithm 1. Then, we have $RE(w) = RE_1(w)RE_2(w) \cdots RE_n(w)$.

Example 3.12. Let w be defined as in Example 3.8. We have

$$RE(w) = \underbrace{\mathbf{z}}_{RE_1(w)} \underbrace{\mathbf{s}_3 \mathbf{t}_1 \mathbf{t}_0 \mathbf{s}_3}_{RE_3(w)} \underbrace{\mathbf{s}_4 \mathbf{s}_3 \mathbf{t}_1 \mathbf{t}_0}_{RE_4(w)}.$$

In this example, $RE_2(w)$ is the empty word.

Proposition 3.13. Let $w \in G(de, e, n)$. The word $RE(w)$ given by Algorithm 1 is a word representative over \mathbf{X} of $w \in G(de, e, n)$.

Proof. Let $w \in G(de, e, n)$ such that the product of all the nonzero entries of w is equal to ζ_d^k for some $0 \leq k \leq d-1$. Algorithm 1 transforms the matrix w into I_n by multiplying it on the right by elements of X . We get $wx_1 \cdots x_{r-1}x_r = I_n$, where x_1, \dots, x_{r-1} are elements of $X \setminus \{z\}$ and $x_r = z^{-k}$. Hence $w = x_r^{-1}x_{r-1}^{-1} \cdots x_1^{-1} = z^k x_{r-1} \cdots x_1$ since $x_i^2 = 1$ for $1 \leq i \leq r-1$. The output $RE(w)$ of Algorithm 1 is $RE(w) = \mathbf{z}^k \mathbf{x}_{r-1} \cdots \mathbf{x}_1$. Hence it is a word representative over \mathbf{X} of $w \in G(de, e, n)$. \square

The following proposition will prepare us to prove that the output of Algorithm 1 is a reduced expression over \mathbf{X} of a given element $w \in G(de, e, n)$.

Proposition 3.14. Let w be an element of $G(de, e, n)$. For all $x \in X$, we have

$$\ell(RE(xw)) \leq \ell(RE(w)) + 1.$$

Proof. For $1 \leq i \leq n$, there exists a unique c_i such that $w[i, c_i] \neq 0$. We denote $w[i, c_i]$ by a_i . We have $\prod_{i=1}^n a_i = \zeta_d^k$ for some $0 \leq k \leq d-1$.

Case 1. $a = 1$ Suppose $x = s_i$ for $3 \leq i \leq n$.

A similar case is done in our previous work [18]. We get (1), (2), (3), and (4) as in the proof of Proposition 3.11 (Case 1) in [18]. For completeness, we include this part of the proof in this article.

Set $w' := s_i w$. Since the left multiplication by the matrix x exchanges the rows $i-1$ and i of w and the other rows remain the same, by Definition 3.11 and Lemma 3.9, we have:

$$\begin{aligned} RE_{i+1}(xw)RE_{i+2}(xw) \cdots RE_n(xw) &= RE_{i+1}(w)RE_{i+2}(w) \cdots RE_n(w), \\ RE_2(xw)RE_3(xw) \cdots RE_{i-2}(xw) &= RE_2(w)RE_3(w) \cdots RE_{i-2}(w). \end{aligned}$$

Then, in order to prove our property, we have to compare $\ell_1 := \ell(RE_{i-1}(w)RE_i(w))$ and $\ell_2 := \ell(RE_{i-1}(xw)RE_i(xw))$.

Subcase 1.1. Suppose $c_{i-1} < c_i$. By Lemma 3.9, the rows $i-1$ and i of the blocks w_i and w'_i are of the form:

$$\begin{array}{lcl} w_i : & \begin{array}{c} i-1 \\ i \end{array} & \begin{array}{|c|} \hline \begin{array}{ccccc} \cdots & c & \cdots & c' & \cdots & i \end{array} \\ \hline \begin{array}{ccccc} b_{i-1} & & & & \end{array} \\ \hline \begin{array}{ccccc} & & & a_i & \end{array} \\ \hline \end{array} \\ \\ w'_i : & \begin{array}{c} i-1 \\ i \end{array} & \begin{array}{|c|} \hline \begin{array}{ccccc} \cdots & c & \cdots & c' & \cdots & i \end{array} \\ \hline \begin{array}{ccccc} & & & a_i & \end{array} \\ \hline \begin{array}{ccccc} b_{i-1} & & & & \end{array} \\ \hline \end{array} \end{array}$$

with $c < c'$ and where we write b_{i-1} instead of a_{i-1} since a_{i-1} may change when applying Algorithm 1 if $c_{i-1} = 1$, that is a_{i-1} on the first column of w (see the second item of Lemma 3.9).

We will discuss different cases depending on the values of a_i and b_{i-1} .

• Suppose $a_i = 1$.

– Suppose $b_{i-1} = 1$.

We have $RE_i(w) = s_i \cdots s_{c'+2} s_{c'+1}$ and $RE_{i-1}(w) = s_{i-1} \cdots s_{c+2} s_{c+1}$. Furthermore, we have $RE_i(xw) = s_i \cdots s_{c+2} s_{c+1}$ and $RE_{i-1}(xw) = s_{i-1} \cdots s_{c'+1} s_{c'}$.

It follows that $\ell_1 = ((i-1) - (c+1) + 1) + (i - (c'+1) + 1) = 2i - c - c' - 1$ and $\ell_2 = ((i-1) - c' + 1) + (i - (c+1) + 1) = 2i - c - c'$ hence $\ell_2 = \ell_1 + 1$.

– Suppose $b_{i-1} = \zeta_{de}^k$ for some $1 \leq k \leq de - 1$.

We have $RE_i(w) = s_i \cdots s_{c'+2} s_{c'+1}$ and $RE_{i-1}(w) = s_{i-1} \cdots s_3 t_k t_0 s_3 \cdots s_c$. Furthermore, we have $RE_i(xw) = s_i \cdots s_3 t_k t_0 s_3 \cdots s_c$ and $RE_{i-1}(xw) = s_{i-1} \cdots s_{c'}$.

It follows that $\ell_1 = (((i-1) - 3 + 1) + 2 + (c - 3 + 1)) + (i - (c' + 1) + 1) = 2i + c - c' - 3$ and $\ell_2 = ((i-1) - c' + 1) + ((i-3+1) + 2 + (c-3+1)) = 2i + c - c' - 2$ hence $\ell_2 = \ell_1 + 1$.

It follows that

$$\text{if } a_i = 1, \text{ then } \ell(RE(s_i w)) = \ell(RE(w)) + 1. \quad (1)$$

• Suppose now that $a_i = \zeta_{de}^k$ with $1 \leq k \leq de - 1$.

– Suppose $b_{i-1} = 1$.

We have $RE_i(w) = s_i \cdots s_3 t_k t_0 s_3 \cdots s_{c'}$ and $RE_{i-1}(w) = s_{i-1} \cdots s_{c+1}$.

Also, we have $RE_i(xw) = s_i \cdots s_{c+1}$ and $RE_{i-1}(xw) = s_{i-1} \cdots s_3 t_k t_0 s_3 \cdots s_{c'-1}$.

It follows that $\ell_1 = ((i-1) - (c+1) - 1) + ((i-3+1) + 2 + (c' - 3 + 1)) = 2i - c + c' - 5$ and $\ell_2 = (((i-1) - 3 + 1) + 2 + ((c' - 1) - 3 + 1)) + (i - (c+1) - 1) = 2i - c + c' - 6$ hence $\ell_2 = \ell_1 - 1$.

– Suppose $b_{i-1} = \zeta_{de}^{k'}$ for some $1 \leq k' \leq de - 1$.

We have $RE_i(w) = s_i \cdots s_3 t_k t_0 s_3 \cdots s_{c'}$ and $RE_{i-1}(w) = s_{i-1} \cdots s_3 t_{k'} t_0 s_3 \cdots s_c$.

Also, we have $RE_i(xw) = s_i \cdots s_3 t_{k'} t_0 s_3 \cdots s_c$ and $RE_{i-1}(xw) = s_{i-1} \cdots s_3 t_k t_0 s_3 \cdots s_{c'-1}$.

It follows that $\ell_1 = ((i-1) - 3 + 1) + 2 + (c - 3 + 1) + (i - 3 + 1) + 2 + (c' - 3 + 1) = 2i + c + c' - 5$ and $\ell_2 = ((i-1) - 3 + 1) + 2 + ((c' - 1) - 3 + 1) + (i - 3 + 1) + 2 + (c - 3 + 1) = 2i + c + c' - 6$ hence $\ell_2 = \ell_1 - 1$.

It follows that

$$\text{if } a_i \neq 1, \text{ then } \ell(RE(s_i w)) = \ell(RE(w)) - 1. \quad (2)$$

Subcase 1.2. Suppose $c_{i-1} > c_i$. Recall that $w' = s_i w$. If $w'[i-1, c'_{i-1}]$ and $w'[i, c'_i]$ denote the nonzero entries of w' on the rows $i-1$ and i , respectively, we have $w'[i-1, c'_{i-1}] = a_i$ and $w'[i, c'_i] = a_{i-1}$. For w' , we have $c'_{i-1} < c'_i$, in which case the preceding analysis would give:

$$\begin{aligned} &\text{if } a_{i-1} = 1, \text{ then } \ell(RE(s_i(s_i w))) = \ell(RE(s_i w)) + 1, \\ &\text{if } a_{i-1} \neq 1, \text{ then } \ell(RE(s_i(s_i w))) = \ell(RE(s_i w)) - 1. \end{aligned}$$

Hence, since $s_i^2 = 1$, we get the following:

$$\text{if } a_{i-1} = 1, \text{ then } \ell(RE(s_i w)) = \ell(RE(w)) - 1. \quad (3)$$

$$\text{if } a_{i-1} \neq 1, \text{ then } \ell(RE(s_i w)) = \ell(RE(w)) + 1. \quad (4)$$

Case 2. Suppose $x = t_i$ for $0 \leq i \leq de - 1$.

Set $w' := t_i w$. By the left multiplication by t_i , we have that the last $n - 2$ rows of w and w' are the same. Hence, by Definition 3.11 and Lemma 3.9, we have:

$RE_3(xw)RE_4(xw) \cdots RE_n(xw) = RE_3(w)RE_4(w) \cdots RE_n(w)$. In order to prove our property in this case, we should compare $\ell_1 := \ell(RE_1(w)RE_2(w))$ and $\ell_2 := \ell(RE_1(xw)RE_2(xw))$.

Subcase 2.1. Consider the case where $c_1 < c_2$. Since $c_1 < c_2$, by Lemma 3.9, the blocks w_2 and w'_2 are of the form: $w_2 = \begin{pmatrix} b_1 & 0 \\ 0 & a_2 \end{pmatrix}$ and $w'_2 = \begin{pmatrix} 0 & \zeta_{de}^{-i} a_2 \\ \zeta_{de}^i b_1 & 0 \end{pmatrix}$ with b_1 instead of a_1 since a_1 may change when applying Algorithm 1 if $c_1 = 1$.

- Suppose $a_2 = 1$.

We have $b_1 = \zeta_d^k$ hence $\ell_1 = k$. We also have $RE_2(xw) = t_{i+ke}$ and $RE_1(xw) = z^k$. Hence we get $\ell_2 = k + 1$. It follows that when $c_1 < c_2$,

$$\text{if } a_2 = 1, \text{ then } \ell(RE(t_i w)) = \ell(RE(w)) + 1. \quad (5)$$

- Suppose $a_2 = \zeta_{de}^{k'}$ for some $1 \leq k' \leq de - 1$.

We have $b_1 = \zeta_{de}^{ke-k'}$. We get $RE_2(w) = t_{k'} t_0$ and $RE_1(w) = z^k$. Thus, $\ell_1 = k + 2$. We also get $RE_2(xw) = t_{ke+i-k'}$ and $RE_1(xw) = z^k$. Thus, $\ell_2 = k + 1$. It follows that when $c_1 < c_2$,

$$\text{if } a_2 \neq 1, \text{ then } \ell(RE(t_i w)) = \ell(RE(w)) - 1. \quad (6)$$

Subcase 2.2. Consider the case where $c_1 > c_2$. Since $c_1 > c_2$, by Lemma 3.9, the blocks w_2 and w'_2 are of the form: $w_2 = \begin{pmatrix} 0 & a_1 \\ b_2 & 0 \end{pmatrix}$ and $w'_2 = \begin{pmatrix} \zeta_{de}^{-i} b_2 & 0 \\ 0 & \zeta_{de}^i a_1 \end{pmatrix}$ with b_2 instead of a_2 since a_2 may change when applying Algorithm 1 if $c_2 = 1$.

- Suppose $a_1 \neq \zeta_{de}^{-i}$.

We have $\ell_1 = k + 1$, and since $\zeta_{de}^i a_1 \neq 1$, we have $\ell_2 = k + 2$. Hence when $c_1 > c_2$,

$$\text{if } a_1 \neq \zeta_{de}^{-i}, \text{ then } \ell(RE(t_i w)) = \ell(RE(w)) + 1. \quad (7)$$

- Suppose $a_1 = \zeta_{de}^{-i}$. We have $b_2 = \zeta_{de}^{i+ek}$. We get $\ell_1 = k + 1$ and $\ell_2 = k$. Hence when $c_1 > c_2$,

$$\text{if } a_1 = \zeta_{de}^{-i}, \text{ then } \ell(RE(t_i w)) = \ell(RE(w)) - 1. \quad (8)$$

Case 3: Suppose $x = z$.

Set $w' := zw$. By the left multiplication by z , we have that the last $n - 1$ rows of w and w' are the same. Hence, by Definition 3.11 and Lemma 3.9, we have:

$RE_2(xw)RE_3(xw) \cdots RE_n(xw) = RE_2(w)RE_3(w) \cdots RE_n(w)$. In order to prove our property in this case, we should compare $\ell_1 := \ell(RE_1(w))$ and $\ell_2 := \ell(RE_1(xw))$.

We get w_1 is equal to b_1 and $w'_1 = \zeta_d b_1$ with b_1 instead of a_1 since a_1 may change when applying Algorithm 1 if $c_1 = 1$. We have $b_1 = \prod_{i=1}^n a_i = \zeta_d^k$ for some $0 \leq k \leq d-1$. Hence if $k \neq d-1$, we get $\ell_1 = k$ and $\ell_2 = k+1$ and if $k = d-1$, we get $\ell_1 = d-1$ and $\ell_2 = 0$. It follows that

$$\ell(RE(zw)) \leq \ell(RE(w)) + 1. \quad \square \quad (9)$$

The next proposition establishes that Algorithm 1 produces geodesic normal forms for $G(de, e, n)$.

Proposition 3.15. *Let w be an element of $G(de, e, n)$. The word $RE(w)$ is a reduced expression over \mathbf{X} of w .*

Proof. We must prove that $\ell(w) = \ell(RE(w))$. Let $\mathbf{x}_1 \mathbf{x}_2 \cdots \mathbf{x}_r$ be a reduced expression over \mathbf{X} of w . Hence $\ell(w) = \ell(\mathbf{x}_1 \mathbf{x}_2 \cdots \mathbf{x}_r) = r$. Since $RE(w)$ is a word representative over \mathbf{X} of w , we have $\ell(RE(w)) \geq \ell(\mathbf{x}_1 \mathbf{x}_2 \cdots \mathbf{x}_r) = r$.

We prove that $\ell(RE(w)) \leq r$. Write w as $x_1 x_2 \cdots x_r$ where x_1, x_2, \dots, x_r are the matrices of $G(de, e, n)$ corresponding to $\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_r$. By Proposition 3.14, we have: $\ell(RE(w)) = \ell(RE(x_1 x_2 \cdots x_r)) \leq \ell(RE(x_2 x_3 \cdots x_r)) + 1 \leq \ell(RE(x_3 \cdots x_r)) + 2 \leq \cdots \leq r$. Hence $\ell(RE(w))$ is equal to $\ell(w)$. This establishes that $RE(w)$ is a reduced expression over \mathbf{X} of w . \square

Remark 3.16. Geodesic normal forms for the complex reflection groups $G(e, e, n)$ have been already established in our previous work [18]. They are explicitly defined by an algorithm (similar to Algorithm 1). Let $w \in G(e, e, n)$. The output of the algorithm is the word also denoted by $RE(w)$ and defined as a concatenation of the words $RE_2(w), RE_3(w), \dots, RE_n(w)$ introduced as in Definition 3.11. It describes a minimal word representative of the element $w \in G(e, e, n)$.

As a direct consequence of Algorithm 1 and Proposition 3.15, the next statement characterizes the elements of $G(de, e, n)$ that are of maximal length over the generating set of Corran–Lee–Lee.

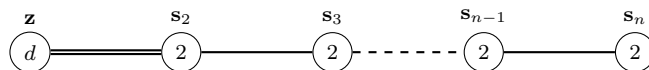
Proposition 3.17. *Let $d > 1, e > 1$ and $n \geq 2$. The maximal length of an element of $G(de, e, n)$ over the generating set of Corran–Lee–Lee is $n(n-1) + d - 1$. It is realized for diagonal matrices w such that for all $2 \leq i \leq n$, we have $w[i, i] = \zeta_{de}^{k_i}$ with $1 \leq k_i \leq de - 1$ and $w[1, 1] = \zeta_{de}^x$ with $x + (k_2 \cdots k_n) = e(d-1)$. A minimal word representative of such an element is of the form*

$$\mathbf{z}^{d-1}(\mathbf{t}_{k_2} \mathbf{t}_0)(\mathbf{s}_3 \mathbf{t}_{k_3} \mathbf{t}_0 \mathbf{s}_3) \cdots (\mathbf{s}_n \cdots \mathbf{s}_3 \mathbf{t}_{k_n} \mathbf{t}_0 \mathbf{s}_3 \cdots \mathbf{s}_n),$$

with $1 \leq k_2, \dots, k_n \leq de - 1$. The number of elements that are of maximal length is then $(de - 1)^{(n-1)}$.

3.3. The case of $G(d, 1, n)$

We establish a similar construction for the case of $G(d, 1, n)$ for $d > 1$ and $n \geq 2$. We recall the diagram of the presentation of $G(d, 1, n)$:



Denote by \mathbf{X} the set $\{\mathbf{z}, \mathbf{s}_2, \dots, \mathbf{s}_n\}$ of the generators. The generator \mathbf{z} corresponds to the matrix $z := \text{Diag}(\zeta_d, 1, \dots, 1)$ in $G(d, 1, n)$ with $\zeta_d = \exp(2i\pi/d)$ and \mathbf{s}_j corresponds to the transposition matrix $s_j := (j-1, j)$ for $2 \leq j \leq n$. Denote by X the set $\{z, s_2, s_3, \dots, s_n\}$ of these matrices.

Input : w , a matrix in $G(d, 1, n)$, with $d > 1$ and $n \geq 2$.

Output: $RE(w)$, a word over \mathbf{X} .

Local variables: w' , $RE(w)$, i , U , c , k .

Initialisation: $U := [1, \zeta_d, \zeta_d^2, \dots, \zeta_d^{d-1}]$ ($U[1] = 1$), $RE(w) := \varepsilon$: the empty word, $w' := w$.

```

for  $i$  from  $n$  down to 1 do
     $c := 1$ ;  $k := 0$ ;
    while  $w'[i, c] = 0$  do
         $c := c + 1$ ;
    end
    #Then  $w'[i, c]$  is the root of unity on the row  $i$ ;
    while  $U[k + 1] \neq w'[i, c]$  do
         $k := k + 1$ ;
    end
    #Then  $w'[i, c] = \zeta_d^k$ .

    if  $k \neq 0$  then
         $w' := w' s_c s_{c-1} \dots s_3 s_2 z^{-k}$ ; #Then  $w'[i, 2] = 1$ ;
         $RE(w) := \mathbf{z}^k \mathbf{s}_2 \mathbf{s}_3 \dots \mathbf{s}_c RE(w)$ ;
         $c := 1$ ;
    end
     $w' := w' s_{c+1} \dots s_{i-1} s_i$ ; #Then  $w'[i, i] = 1$ ;
     $RE(w) := \mathbf{s}_i \mathbf{s}_{i-1} \dots \mathbf{s}_{c+1} RE(w)$ ;
end
Return  $RE(w)$ ;

```

Algorithm 2: A word over \mathbf{X} corresponding to an element $w \in G(d, 1, n)$.

We define Algorithm 2 that produces a word $RE(w)$ for each matrix w of $G(d, 1, n)$. This Algorithm is different than Algorithm 1. Let us explain the steps of the algorithm. Let $w_n := w \in G(d, 1, n)$. For i from n to 1, the i -th step of the algorithm transforms the block diagonal matrix $\left(\begin{array}{c|c} w_i & 0 \\ \hline 0 & I_{n-i} \end{array} \right)$ into a block diagonal matrix $\left(\begin{array}{c|c} w_{i-1} & 0 \\ \hline 0 & I_{n-i+1} \end{array} \right) \in G(d, 1, n)$. Let $w_i[i, c] \neq 0$ be the nonzero coefficient on the row i of w_i . If $w_i[i, c] = 1$, we shift it into the diagonal position $[i, i]$ by right multiplication by transpositions. If $w_i[i, c] = \zeta_d^k$ with $k \geq 1$, we shift it into position $[i, 1]$ by right multiplication by transpositions, followed by a right multiplication by z^{-k} , then we shift the 1 obtained in position $[i, 1]$ into the diagonal position $[i, i]$ by right multiplication by transpositions. Let us illustrate these operations by the following example.

Example 3.18. Let $w := \begin{pmatrix} 0 & \zeta_3 & 0 \\ 0 & 0 & \zeta_3^2 \\ \zeta_3^2 & 0 & 0 \end{pmatrix} \in G(3, 1, 3)$.

Step 1 ($i = 3, k = 2, c = 1$): $w' := wz^{-2} = \begin{pmatrix} 0 & \zeta_3 & 0 \\ 0 & 0 & \zeta_3^2 \\ 1 & 0 & 0 \end{pmatrix}$, then

$$w' := w' s_2 s_3 = \begin{pmatrix} \zeta_3 & 0 & 0 \\ 0 & \boxed{\zeta_3^2} & 0 \\ 0 & 0 & 1 \end{pmatrix}.$$

Step 2 ($i = 2, k = 2, c = 1$): $w' := w' s_2 = \begin{pmatrix} 0 & \zeta_3 & 0 \\ \zeta_3^2 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}$, then $w' := w' z^{-2} = \begin{pmatrix} 0 & \zeta_3 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}$, then

$$w' := w' s_2 = \begin{pmatrix} \boxed{\zeta_3} & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}.$$

Step 3 ($i = 1, k = 1, c = 1$): $w' := w' z^{-1} = I_3$.

Hence $RE(w) = \mathbf{z} \mathbf{s}_2 \mathbf{z}^2 \mathbf{s}_2 \mathbf{s}_3 \mathbf{s}_2 \mathbf{z}^2$.

The next lemma follows directly from Algorithm 2.

Lemma 3.19. For $2 \leq i \leq n$, let $w_i[i, c] \neq 0$ be the nonzero coefficient on the row i of w_i . The block w_{i-1} is obtained by removing the row i and the column c from w_i . Moreover, w_1 is equal to the nonzero entry on the first row of w .

Definition 3.20. Let $1 \leq i \leq n$. Let $w_i[i, c] \neq 0$ for $1 \leq c \leq i$.

- If $w_1 = \zeta_d^k$ for some $0 \leq k \leq d-1$ (this is equal to the nonzero entry on the first row of w), we define $RE_1(w)$ to be the word \mathbf{z}^k .
- If $w_i[i, c] = 1$, we define $RE_i(w)$ to be the word $\mathbf{s}_i \mathbf{s}_{i-1} \cdots \mathbf{s}_{c+1}$ (decreasing-index expression).
- If $w_i[i, c] = \zeta_d^k$ with $k \neq 0$, we define $RE_i(w)$ to be the word

$$\begin{aligned} &\mathbf{s}_i \cdots \mathbf{s}_3 \mathbf{z}^k && \text{if } c = 1, \\ &\mathbf{s}_i \cdots \mathbf{s}_3 \mathbf{s}_2 \mathbf{z}^k \mathbf{s}_2 \mathbf{s}_3 \cdots \mathbf{s}_c && \text{if } c \geq 2. \end{aligned}$$

As for Algorithm 1, the output of Algorithm 2 is equal to $RE_1(w)RE_2(w) \cdots RE_n(w)$. In Example 3.18, we have $RE(w) = \underbrace{\mathbf{z}}_{RE_1(w)} \underbrace{\mathbf{s}_2 \mathbf{z}^2 \mathbf{s}_2}_{RE_2(w)} \underbrace{\mathbf{s}_3 \mathbf{s}_2 \mathbf{z}^2}_{RE_3(w)}$.

The proof of the next proposition is similar to the proof of Proposition 3.13 and is left to the reader.

Proposition 3.21. Let $w \in G(d, 1, n)$. The word $RE(w)$ given by Algorithm 2 is a word representative over \mathbf{X} of $w \in G(d, 1, n)$.

The following proposition enables us to prove that the output of Algorithm 2 is a reduced expression over \mathbf{X} of a given element $w \in G(d, 1, n)$.

Proposition 3.22. Let w be an element of $G(d, 1, n)$. For all $x \in X$, we have

$$\ell(RE(xw)) \leq \ell(RE(w)) + 1.$$

Proof. Consider $x = s_i$ with $2 \leq i \leq n$. This case is done in the same way as Case 1 in the proof of Proposition 3.14. Consider now $x = z$. This case is done this time as Case 3 of the proof of Proposition 3.14. \square

Applying the arguments used before in the proof of Proposition 3.15, we deduce that $RE(w)$ is a reduced expression over \mathbf{X} of $w \in G(d, 1, n)$. Hence Algorithm 2 produces geodesic normal forms for $G(d, 1, n)$.

As a direct application, the next statement characterizes the elements of $G(d, 1, n)$ that are of maximal length. Note that this statement was also observed in [4].

Proposition 3.23. Let $d > 1$ and $n \geq 2$. There exists a unique element of maximal length of $G(d, 1, n)$. Its minimal word representative is of the form

$$\mathbf{z}^{d-1}(\mathbf{s}_2 \mathbf{z}^{d-1} \mathbf{s}_2)(\mathbf{s}_3 \mathbf{s}_2 \mathbf{z}^{d-1} \mathbf{s}_2 \mathbf{s}_3) \cdots (\mathbf{s}_n \cdots \mathbf{s}_2 \mathbf{z}^{d-1} \mathbf{s}_2 \cdots \mathbf{s}_n).$$

Its length is then equal to $n(n+d-2)$.

Remark 3.24. When $d = 2$, the group $G(2, 1, n)$ is the Coxeter group of type B_n . By Proposition 3.23, the longest element is of the form

$$\mathbf{z}(\mathbf{s}_2\mathbf{z}\mathbf{s}_2)(\mathbf{s}_3\mathbf{s}_2\mathbf{z}\mathbf{s}_2\mathbf{s}_3)\cdots(\mathbf{s}_n\cdots\mathbf{s}_2\mathbf{z}\mathbf{s}_2\cdots\mathbf{s}_n).$$

Its length is equal to n^2 which is already known for Coxeter groups of type B_n , see Example 1.4.6 of [11].

4. The Hecke algebras $H(de, e, n)$

The Hecke algebras $H(de, e, n)$ attached to the general series of complex reflection groups $G(de, e, n)$ are defined as quotients of the corresponding complex braid group algebras by some polynomial relations. Let $B(de, e, n)$ denote the complex braid group attached to $G(de, e, n)$, as defined in [6]. We establish presentations for the Hecke algebras $H(de, e, n)$ by using the presentations of Corran–Picantin [10] and Corran–Lee–Lee [9] of the complex braid groups $B(e, e, n)$ and $B(de, e, n)$ for $d > 1$, respectively.

4.1. The Hecke algebras $H(e, e, n)$

Corran–Picantin introduced in [10] a presentation for the complex braid groups $B(e, e, n)$. Although we are not going to use this result, we mention that they also established nice structures (called Garside structures) for these groups. The presentation of Corran–Picantin is as follows.

Definition 4.1. Let $e \geq 1$ and $n \geq 2$. The group $B(e, e, n)$ is defined by a presentation with set of generators $\{t_i \mid i \in \mathbb{Z}/e\mathbb{Z}\} \cup \{s_3, s_4, \dots, s_n\}$ in bijection with $\{\mathbf{t}_i \mid i \in \mathbb{Z}/e\mathbb{Z}\} \cup \{\mathbf{s}_3, \mathbf{s}_4, \dots, \mathbf{s}_n\}$ and relations 1 to 5 of Definition 3.1.

Adding the quadratic relations to all the generators, we get the presentation of Corran–Picantin of $G(e, e, n)$ given earlier in Definition 3.1.

The next definition establishes a presentation of the Hecke algebra $H(e, e, n)$ attached to the group $G(e, e, n)$, by using the presentation of Corran–Picantin of the complex braid group $B(e, e, n)$.

Definition 4.2. Let $e \geq 1$ and $n \geq 2$. We exclude the case $(n = 2, e \text{ even})$, see Remark 4.3 below. Let $R_0 = \mathbb{Z}[a]$. The unitary associative Hecke algebra $H(e, e, n)$ is defined as the quotient of the group algebra $R_0(B(e, e, n))$ by the following relations:

1. $t_i^2 - at_i - 1 = 0$ for $i \in \mathbb{Z}/e\mathbb{Z}$,
2. $s_j^2 - as_j - 1 = 0$ for $3 \leq j \leq n$,

where $\{t_i \mid i \in \mathbb{Z}/e\mathbb{Z}\} \cup \{s_j \mid 3 \leq j \leq n\}$ is the set of generators of the presentation of Corran–Picantin of $B(e, e, n)$. Then, a presentation of $H(e, e, n)$ is obtained by adding these relations to those of the presentation of Corran–Picantin given in Definition 4.1.

Remark 4.3. For the case $(n = 2, e \text{ even})$, there exist two conjugacy classes of the reflections t_i , for $i \in \mathbb{Z}/e\mathbb{Z}$ in the complex reflection group. In this case, we define the Hecke algebra $H(e, e, 2)$ in the same way as in Definition 4.2 over $R_0 = \mathbb{Z}[a_1, a_2]$ with two types of polynomial relations for each conjugacy class of the t_i 's: $t_i^2 - a_1 t_i - 1 = 0$ for the first conjugacy class and $t_j^2 - a_2 t_j - 1 = 0$ for the second.

Note that we use the polynomial ring R_0 instead of the usual Laurent polynomial ring R introduced in Definition 2.2 and we use normalized polynomial relations in the definition of $H(e, e, n)$. Actually, by a result of Marin (see Proposition 2.3 in [14]) applied to the case of $G(e, e, n)$, the BMR freeness conjecture for this case is equivalent to the fact that $H(e, e, n)$ is a free R_0 -module of rank equal to the order of $G(e, e, n)$. We will also use a polynomial ring and normalized relations in the definition of the Hecke algebras attached to all the groups $G(de, e, n)$ in the next subsection.

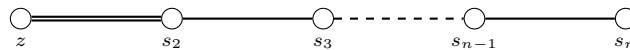
4.2. The general case

Corran–Lee–Lee [9] established a presentation for the complex braid group $B(de, e, n)$ that give rise to nice structures (called quasi-Garside structures) for these groups. The presentation is defined as follows.

Definition 4.4. Let $d > 1$, $e \geq 1$ and $n \geq 2$. The group $B(de, e, n)$ is defined by a presentation with set of generators $\{z\} \cup \{t_i \mid i \in \mathbb{Z}\} \cup \{s_3, s_4, \dots, s_n\}$ in bijection with $\{\mathbf{z}\} \cup \{\mathbf{t}_i \mid i \in \mathbb{Z}/d\mathbb{Z}\} \cup \{\mathbf{s}_3, \mathbf{s}_4, \dots, \mathbf{s}_n\}$ and relations 1 to 3 of Definition 3.3.

Note that $B(de, e, n)$ is isomorphic to $B(2e, e, n)$ for $d > 1$. The parameter d makes an appearance when it comes to the complex reflection group $G(de, e, n)$. Adding the relation $z^d = 1$ and the quadratic relations to all the other generators of the presentation of Corran–Lee–Lee of $B(de, e, n)$, we obtain the presentation of $G(de, e, n)$ given earlier in Definition 3.3. Actually, with the additional relation $z^d = 1$, we have $t_{i+de} = z^d t_{i+de} = t_i z^d = t_i$ for all $i \in \mathbb{Z}$.

Note that for $e = 1$, it is readily checked that the presentation of Corran–Lee–Lee is equivalent to the classical presentation of $B(d, 1, n)$ that is isomorphic to $B(2, 1, n)$ for $d > 1$. Note that we replace t_0 by s_2 in the set of generators. The (well-known) diagram that describes the classical presentation of $B(d, 1, n)$ is the following.



We are ready to establish a presentation of the Hecke algebra $H(de, e, n)$ attached to the group $G(de, e, n)$, by using the presentation of Corran–Lee–Lee of the complex braid group $B(de, e, n)$. Similarly to the case of $H(e, e, n)$, we also define the Hecke algebra $H(de, e, n)$ over a polynomial ring R_0 and use normalized polynomial relations.

Definition 4.5. Let $d > 1$, $e \geq 1$ and $n \geq 2$. We exclude the case $(n = 2, e \text{ even})$, see Remark 4.6 below. Let $R_0 = \mathbb{Z}[a, b_1, b_2, \dots, b_{d-1}]$. The unitary associative Hecke algebra $H(de, e, n)$ is defined as the quotient of the group algebra $R_0(B(de, e, n))$ by the following relations:

1. $z^d - b_1 z^{d-1} - b_2 z^{d-2} - \dots - b_{d-1} z - 1 = 0$,
2. $t_i^2 - at_i - 1 = 0$ for $i \in \mathbb{Z}$,
3. $s_j^2 - as_j - 1 = 0$ for $3 \leq j \leq n$,

where $\{z\} \cup \{t_i \mid i \in \mathbb{Z}\} \cup \{s_j \mid 3 \leq j \leq n\}$ is the set of generators of the presentation of Corran–Lee–Lee of $B(de, e, n)$. Then, a presentation of $H(de, e, n)$ is obtained by adding these relations to those given in Definition 4.4.

Remark 4.6. When $(n = 2, e \text{ even})$, the Hecke algebra $H(de, e, 2)$ can be defined over $R_0[a_1, a_2, b_1, b_2, \dots, b_{d-1}]$ in the same way as in the previous definition, but with two types of polynomial relations for the t_i 's (due to the existence of two conjugacy classes of the t_i 's in $G(de, e, 2)$), as established before in Remark 4.3.

Remark 4.7. The generators of Corran–Picantin for $B(e, e, n)$ and of Corran–Lee–Lee for $B(de, e, n)$ include the generators of Broué–Malle–Rouquier $\mathbf{t}_0, \mathbf{t}_1, \mathbf{s}_3, \mathbf{s}_4, \dots, \mathbf{s}_n$ for $B(e, e, n)$ and $\mathbf{z}, \mathbf{t}_0, \mathbf{t}_1, \mathbf{s}_3, \mathbf{s}_4, \dots, \mathbf{s}_n$ for $B(de, e, n)$ that are distinguished braided reflections (see [6]). By Relation 1 of Definition 3.1 and Relation 3 of Definition 3.3, the generators \mathbf{t}_i ($i \neq 0, 1$) are all conjugate to either \mathbf{t}_0 or \mathbf{t}_1 . Hence their corresponding braided reflections are conjugate in the corresponding complex braid group (see also [6]). The definition of $H(e, e, n)$ and $H(de, e, n)$ then coincides with Marin's definition of H_0 (see the paragraph after Theorem 2.2 in [14]), as it is defined to be the quotient of the complex braid group algebra over R_0 by the same normalized polynomial relations.

5. Bases for the Hecke algebras $H(e, e, n)$

The Hecke algebra $H(e, e, n)$ is described in Definition 4.2 by a presentation with generating set $\{t_0, t_1, \dots, t_{e-1}, s_3, \dots, s_n\}$. It is defined over $R_0 = \mathbb{Z}[a]$. Note that we will replace t_0 by s_2 in some cases in order to simplify notations. Using the geodesic normal forms of $G(e, e, n)$ introduced in Section 3, we construct a natural basis for $H(e, e, n)$ that is different from the one introduced by Ariki in [1].

Let us define the following subsets of $H(e, e, n)$:

$$\Lambda_2 = \{1, \\ t_k \quad \text{for } 0 \leq k \leq e-1, \\ t_k t_0 \quad \text{for } 1 \leq k \leq e-1\},$$

and for $3 \leq i \leq n$,

$$\Lambda_i = \{1, \\ s_i \cdots s_{i'} \quad \text{for } 3 \leq i' \leq i, \\ s_i \cdots s_3 t_k \quad \text{for } 0 \leq k \leq e-1, \\ s_i \cdots s_3 t_k s_2 \cdots s_{i'} \quad \text{for } 1 \leq k \leq e-1 \text{ and } 2 \leq i' \leq i\}.$$

Define $\Lambda = \Lambda_2 \cdots \Lambda_n$ to be the set of the products $a_2 \cdots a_n$, where $a_2 \in \Lambda_2, \dots, a_n \in \Lambda_n$. Remark that this set corresponds to all the reduced words $RE(w)$ of the form $RE_2(w)RE_3(w) \cdots RE_n(w)$ introduced in Section 3 (see Definition 3.11 and Remark 3.16). Recall that $R_0 = \mathbb{Z}[a]$ (see Definition 4.2). The aim of this section is to establish the following.

Theorem 5.1. *The set Λ provides an R_0 -basis of the Hecke algebra $H(e, e, n)$ for $e \geq 1$ and $n \geq 2$.*

In order to prove this theorem, it is shown in Proposition 2.3(i) of [14] that it is enough to find a spanning set of $H(e, e, n)$ over R_0 of $|G(e, e, n)|$ elements. This is a general fact about Hecke algebras associated to complex reflection groups. We have $|\Lambda_2| = 2e$, $|\Lambda_3| = 3e$, \dots , and $|\Lambda_n| = ne$ by the definition of Λ_2, \dots , and Λ_n . Thus, $|\Lambda|$ is equal to $e^{n-1}n!$ that is the order of $G(e, e, n)$. If we manage to prove that Λ is a spanning set of $H(e, e, n)$ over R_0 , then we get Theorem 5.1. Denote by $\text{Span}(S)$ the sub- R_0 -module of $H(e, e, n)$ generated by S .

We prove Theorem 5.1 by induction on $n \geq 2$. Propositions 5.4 and 5.6 below correspond to the cases $n = 2$ and $n = 3$, respectively. Suppose that $\Lambda_2 \cdots \Lambda_{n-1}$ is an R_0 -basis of $H(e, e, n-1)$. As mentioned before, in order to prove that $\Lambda = \Lambda_2 \cdots \Lambda_n$ is an R_0 -basis of $H(e, e, n)$, it is enough to show that it is an R_0 -generating set of $H(e, e, n)$, that is Λ stable under left multiplication by $t_0, \dots, t_{e-1}, s_3, \dots, s_n$. Since $\Lambda_2 \cdots \Lambda_{n-1}$ is an R_0 -basis of $H(e, e, n-1)$, the set $\Lambda_2 \cdots \Lambda_n$ is stable under left multiplication by $t_0, \dots, t_{e-1}, s_3, \dots, s_{n-1}$. We prove that it is stable under left multiplication by s_n , that is $s_n(a_2 \cdots a_n) = a_2 \cdots a_{n-2} s_n(a_{n-1} a_n)$ belongs to $\text{Span}(\Lambda)$ for $a_2 \in \Lambda_2, \dots, a_n \in \Lambda_n$ by checking all the different possibilities for $a_{n-1} \in \Lambda_{n-1}$ and $a_n \in \Lambda_n$.

Assume $n > 3$. If $a_{n-1} = 1$ or $a_n = 1$, it is readily checked that $s_n(a_{n-1} a_n)$ belongs to $\text{Span}(\Lambda_{n-1} \Lambda_n)$. If $a_{n-1} = s_{n-1} \cdots s_i$ for $2 \leq i \leq n-1$, we distinguish 3 different cases for a_n that belongs to Λ_n . This is done in Lemmas 5.8, 5.9 and 5.10 below. If $a_{n-1} = s_{n-1} \cdots s_3 t_k$ for $0 \leq k \leq e-1$, we also distinguish 3 different cases for $a_n \in \Lambda_n$. This is done in Lemmas 5.11, 5.12 and 5.13. Finally, if $a_{n-1} = s_{n-1} \cdots s_3 t_k s_2 \cdots s_i$ for $1 \leq k \leq e-1$ and $2 \leq i \leq n-1$, we also have 3 different cases for $a_n \in \Lambda_n$ (see Lemmas 5.14, 5.15 and 5.16 below).

Let us start by establishing the following two preliminary lemmas.

Lemma 5.2. For $i, j \in \mathbb{Z}/e\mathbb{Z}$, we have $t_j t_i \in \text{Span}(\Lambda_2)$.

Proof. If $i = j$, then we have $t_i^2 = at_i + 1 \in \text{Span}(\Lambda_2)$. Suppose $i \neq j$. We have $t_i t_{i-1} = t_j t_{j-1}$. If we multiply by t_j on the left and by t_{i-1} on the right, we get $t_j t_i t_{i-1}^2 = t_j^2 t_{j-1} t_{i-1}$. Using the quadratic relations, we have $t_j t_i (at_{i-1} + 1) = (at_j + 1) t_{j-1} t_{i-1}$, that is $at_j t_i t_{i-1} + t_j t_i = at_j t_{j-1} t_{i-1} + t_{j-1} t_{i-1}$. Replacing $t_i t_{i-1}$ by $t_j t_{j-1}$ and $t_j t_{j-1}$ by $t_i t_{i-1}$, we get $at_j^2 t_{j-1} + t_j t_i = at_i t_{i-1}^2 + t_{j-1} t_{i-1}$. Using the quadratic relations, we have $a(at_j + 1) t_{j-1} + t_j t_i = at_i (at_{i-1} + 1) + t_{j-1} t_{i-1}$, that is $a^2 t_1 t_0 + at_{j-1} + t_j t_i = a^2 t_1 t_0 + at_i + t_{j-1} t_{i-1}$. Simplifying this relation, we get

$$t_j t_i = t_{j-1} t_{i-1} + a(t_i - t_{j-1}).$$

Now, we apply the same operations to compute $t_{j-1} t_{i-1}$ and so on until we arrive to a term of the form $t_k t_0$ for some $k \in \mathbb{Z}/e\mathbb{Z}$. Thus, if $i \neq j$, then $t_j t_i$ belongs to $\text{Span}(\Lambda_2 \setminus \{1\})$. \square

Lemma 5.3. For $1 \leq k \leq e-1$, we have $t_k t_0 \in R_0(t_1 t_0) + R_0(t_1 t_0)^2 + \cdots + R_0(t_1 t_0)^k + R_0 t_1 + R_0 t_2 + \cdots + R_0 t_{k-1}$.

Proof. We prove the property by induction on k . The property is clearly satisfied for $k = 1$. Let $k \geq 2$. Suppose $t_{k-1} t_0 \in R_0(t_1 t_0) + R_0(t_1 t_0)^2 + \cdots + R_0(t_1 t_0)^{k-1} + R_0 t_1 + R_0 t_2 + \cdots + R_0 t_{k-2}$. We have that $t_{k+1} t_k = t_k t_{k-1}$. Multiplying by t_{k+1} on the left and by t_0 on the right, we get $t_{k+1}^2 t_k t_0 = t_{k+1} t_k t_{k-1} t_0$. Using the quadratic relations and replacing $t_{k+1} t_k$ with $t_1 t_0$, we get $(at_{k+1} + 1) t_k t_0 = t_1 t_0 t_{k-1} t_0$. After simplifying this relation, we have $t_k t_0 = t_1 t_0 (t_{k-1} t_0) - a(t_1 t_0) t_0$. Using the induction hypothesis, we replace $t_{k-1} t_0$ by its value and we get $t_k t_0 \in t_1 t_0 (R_0(t_1 t_0) + R_0(t_1 t_0)^2 + \cdots + R_0(t_1 t_0)^{k-1} + R_0 t_1 + R_0 t_2 + \cdots + R_0 t_{k-2}) + R_0(t_1 t_0) t_0$. This is equal to $R_0(t_1 t_0)^2 + R_0(t_1 t_0)^3 + \cdots + R_0(t_1 t_0)^k + R_0(t_1 t_0) t_1 + R_0(t_1 t_0) t_2 + \cdots + R_0(t_1 t_0) t_{k-2} + R_0(t_1 t_0) t_0$. Now $(t_1 t_0) t_m$ is equal to $t_{m+1} t_m t_0 \in R_0 t_1 t_0 + R_0 t_{m+1}$ for $1 \leq m \leq k-2$ and $(t_1 t_0) t_0 \in R_0(t_1 t_0) + R_0 t_1$. It follows that $t_k t_0 \in R_0(t_1 t_0) + R_0(t_1 t_0)^2 + \cdots + R_0(t_1 t_0)^k + R_0 t_1 + R_0 t_2 + \cdots + R_0 t_{k-1}$. \square

As a direct consequence of Lemma 5.2, we have the following.

Proposition 5.4. Let $x = t_l$ with $l \in \mathbb{Z}/e\mathbb{Z}$. For all $a_2 \in \Lambda_2$, we have xa_2 belongs to $\text{Span}(\Lambda_2)$.

Remark 5.5. Recall that we excluded the case $(n = 2, e \text{ even})$ in Definition 4.2 (see Remark 4.3). Consider the Hecke algebra $H(e, e, 2)$ with e even. Similarly to Proposition 5.4, one shows that Λ_2 is stable under left multiplication by all the t_i 's for $i \in \mathbb{Z}/e\mathbb{Z}$. Then, applying Proposition (2.3) (i) of [14], we also get that Λ_2 is an R_0 -basis of $H(e, e, 2)$ for the case e even.

Proposition 5.6. For all $a_2 \in \Lambda_2$ and $a_3 \in \Lambda_3$, the element $s_3(a_2 a_3)$ belongs to $\text{Span}(\Lambda_2 \Lambda_3)$.

Proof. The case where $a_1 \in \Lambda_1$ and $a_2 = 1$ is obvious. The case where $a_1 = 1$ and $a_2 \in \Lambda_2$ is also obvious.

Case 1. Suppose $a_2 = t_k$ for $0 \leq k \leq e-1$ and $a_3 = s_3$.

We have $s_3 t_k s_3 = t_k s_3 t_k \in \text{Span}(\Lambda_2 \Lambda_3)$.

Case 2. Suppose $a_2 = t_k$ for $0 \leq k \leq e-1$ and $a_3 = s_3 t_l$ for $0 \leq l \leq e-1$. We have $s_3 t_k s_3 t_l = t_k s_3 t_k t_l$. After replacing $t_k t_l$ by its decomposition over Λ_2 (see Lemma 5.2), we directly have $s_3 t_k s_3 t_l \in \text{Span}(\Lambda_2 \Lambda_3)$.

Case 3. Suppose $a_2 = t_k$ for $0 \leq k \leq e-1$ and $a_3 = s_3 t_l t_0$ or $a_3 = s_3 t_l t_0 s_3$ for $1 \leq l \leq e-1$. We have $s_3 t_k s_3 t_l t_0 s_3 = t_k s_3 t_k t_l t_0 s_3$. By replacing $t_k t_l$ by its value (see Lemma 5.2), we obviously have $s_3 t_k s_3 t_l t_0$ and $s_3 t_k s_3 t_l t_0 s_3$ belong to $\text{Span}(\Lambda_2 \Lambda_3)$.

Case 4. Suppose $a_2 = t_k t_0$ for $1 \leq k \leq e-1$ and $a_3 = s_3$. We have $s_3(a_2 a_3) = s_3 t_k t_0 s_3 \in \text{Span}(\Lambda_2 \Lambda_3)$.

Case 5. Suppose $a_2 = t_k t_0$ with $1 \leq k \leq e-1$ and $a_3 = s_3 t_l$ with $0 \leq l \leq e-1$. We have $s_3(a_2 a_3) = s_3 t_k t_0 s_3 t_l$. Recall that by Lemma 5.3, we have $t_k t_0 \in R_0(t_1 t_0) + R_0(t_1 t_0)^2 + \cdots + R_0(t_1 t_0)^k + R_0 t_1 + R_0 t_2 + \cdots + R_0 t_{k-1}$. Replacing $t_k t_0$ by its value, we have to deal with the following two terms:

$s_3 t_x s_3 t_l$ with $1 \leq x \leq k-1$ and $s_3(t_1 t_0)^x s_3 t_l$ with $1 \leq x \leq k$. The first term is done in Case 2. For the second term, we decrease the power of $(t_1 t_0)$ and use $t_1 t_0 = t_{l+1} t_l$ to get $s_3(t_1 t_0)^{x-1} t_{l+1} t_l s_3 t_l$. We apply a braid relation and then get $s_3(t_1 t_0)^{x-1} t_{l+1} s_3 t_l s_3$. Again, we decrease the power of $(t_1 t_0)$ and use $t_1 t_0 = t_{l+2} t_{l+1}$. We get $s_3(t_1 t_0)^{x-2} t_{l+2} t_{l+1} s_3 t_l s_3 \in R_0 s_3(t_1 t_0)^{x-2} t_{l+2} t_{l+1} s_3 t_l s_3 + R_0 s_3(t_1 t_0)^{x-2} t_{l+2} s_3 t_l s_3$. We continue by decreasing the power of $(t_1 t_0)$ and we get in the next step that $s_3(a_2 a_3)$ belongs to $R_0 s_3(t_1 t_0)^{x-3} t_{l+1} s_3 t_l s_3^2 + R_0 s_3(t_1 t_0)^{x-3} t_{l+2} s_3 t_l s_3^2 + R_0 s_3(t_1 t_0)^{x-3} t_{l+1} s_3 t_l s_3^2 + R_0 s_3(t_1 t_0)^{x-3} t_{l+3} s_3 t_l s_3$. Inductively, we arrive to terms of the form $s_3 t_1 t_0 t_{x'} s_3 t_l (s_3)^{x''}$ ($0 \leq x' \leq e-1$ and $x'' \in \mathbb{N}$). Replace $t_1 t_0$ by $t_{x'+1} t_{x'}$, we get $s_3 t_{x'+1} (t_{x'})^2 s_3 t_l (s_3)^{x''}$ which belongs to $R_0 s_3 t_{x'+1} t_{x'} s_3 t_l (s_3)^{x''} + R_0 s_3 t_{x'+1} s_3 t_l (s_3)^{x''}$. Replacing $t_{x'+1} t_{x'}$ by $t_{l+1} t_l$ and applying a braid relation in the first term, we get $R_0 s_3 t_{l+1} s_3 t_l (s_3)^{x''+1} + R_0 s_3 t_{x'+1} s_3 t_l (s_3)^{x''}$. Since $s_3^2 = a s_3 + 1$, it remains to deal with these 2 terms:

$$s_3 t_x s_3 t_l \text{ and } s_3 t_x s_3 t_l s_3, \text{ for some } 0 \leq x \leq e-1.$$

It is readily checked that these terms belong to $\text{Span}(\Lambda_2 \Lambda_3)$.

Case 6. Suppose $a_2 = t_k t_0$ and $a_3 = s_3 t_l t_0$ or $a_3 = s_3 t_l t_0 s_3$ ($1 \leq k, l \leq e-1$).

By Case 5, we get two terms of the form $s_3 t_x s_3 t_l$ and $s_3 t_x s_3 t_l s_3$. Multiplying them on the right by t_0 then by s_3 , we get that $s_3(a_2 a_3)$ obviously belongs to $\text{Span}(\Lambda_2 \Lambda_3)$. \square

In the sequel, we will indicate by (1) the operation that shifts the underlined letters to the left, by (2) the operation that applies braid relations, and by (3) the one that applies quadratic relations. The following lemma is useful in the proofs of Lemmas 5.13, 5.15 and 5.16 below. Denote by S_{n-1}^* the set of the words over $\{t_0, \dots, t_{e-1}, s_3, \dots, s_{n-1}\}$.

Lemma 5.7. *Let $3 \leq i \leq n$. We have $s_n \cdots s_4 s_3^2 s_4 \cdots s_i$ belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$.*

Proof. If $i = 3$, we have $s_n \cdots s_4 s_3^2 \in R_0 s_n \cdots s_4 s_3 + R_0 s_n \cdots s_4$ (the last term is equal to 1 if $n = 3$). If $i = 4$, we have $s_n \cdots s_4 s_3^2 s_4 \in R_0 s_n \cdots s_4 s_3 s_4 + R_0 s_n \cdots s_4^2 \in R_0 s_3 s_n \cdots s_3 + R_0 s_n \cdots s_4 + R_0 s_n \cdots s_5$. The last term is equal to 1 if $n = 4$. Let $i \geq 5$. We have $s_n \cdots s_4 s_3^2 s_4 \cdots s_i$ belongs to $R_0 s_n \cdots s_4 s_3 s_4 \cdots s_i + R_0 s_n \cdots s_5 s_4^2 s_5 \cdots s_i$. We apply the quadratic relation $s_4^2 = a s_4 + 1$ to the second term and get $R_0 s_n \cdots s_5 s_4 s_5 \cdots s_i + R_0 s_n \cdots s_6 s_5^2 s_6 \cdots s_i$. And so on, we apply quadratic relations. We get terms of the form $s_n \cdots s_{k+1} s_k s_{k+1} \cdots s_i$ with $k+1 \leq i$ and a term of the form $s_n \cdots s_{i+1} s_i s_{i-1}^2 s_i$. We have $s_n \cdots s_{i+1} s_i s_{i-1}^2 s_i$ belongs to $R_0 s_n \cdots s_{i+1} s_i s_{i-1} s_i + R_0 s_n \cdots s_{i+1} s_i + R_0 s_n \cdots s_{i+1} \subseteq R_0 s_{i-1} s_n \cdots s_{i-1} + R_0 s_n \cdots s_{i+1} s_i + R_0 s_n \cdots s_{i+1}$ (the last term is equal to 1 if $i = n$). Hence it belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$.

The other terms are of the form $s_n \cdots s_{k+1} s_k s_{k+1} \cdots s_i$ ($k+1 \leq i$). We have

$$\begin{aligned} s_n \cdots \underline{s_{k+1} s_k s_{k+1}} \cdots s_i & \stackrel{(2)}{=} \\ s_n \cdots s_{k+2} \underline{s_k s_{k+1} s_k s_{k+2}} \cdots s_i & \stackrel{(1)}{=} \\ s_k s_n \cdots \underline{s_{k+2} s_{k+1} s_{k+2}} s_k s_{k+3} \cdots s_i & \stackrel{(2)}{=} \\ s_k s_n \cdots \underline{s_{k+1} s_{k+2} s_{k+1} s_k s_{k+3}} \cdots s_i & \stackrel{(1)}{=} \\ s_k s_{k+1} s_n \cdots s_{k+2} s_{k+1} s_k s_{k+3} \cdots \underline{s_{i-1} s_i}. \end{aligned}$$

We apply the same operations to $\underline{s_{k+3}}, \dots, \underline{s_{i-1}}$ and get

$$\begin{aligned} s_k s_{k+1} \cdots s_{i-2} s_n \cdots \underline{s_k s_i} & \stackrel{(1)}{=} \\ s_k s_{k+1} \cdots s_{i-2} s_n \cdots \underline{s_i s_{i-1} s_i s_{i-2}} \cdots s_k & \stackrel{(2)}{=} \\ s_k s_{k+1} \cdots s_{i-2} s_n \cdots s_{i+1} \underline{s_{i-1} s_i s_{i-1} s_{i-2}} \cdots s_k & \stackrel{(1)}{=} \\ s_k s_{k+1} \cdots s_{i-1} s_n s_{n-1} \cdots s_k, \end{aligned}$$

which belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$. \square

Lemma 5.8. *If $a_{n-1} = s_{n-1}s_{n-2} \cdots s_i$ with $3 \leq i \leq n-1$ and $a_n = s_ns_{n-1} \cdots s_{i'}$ with $3 \leq i' \leq n$, then $s_n(a_{n-1}a_n)$ belongs to $\text{Span}(\Lambda_{n-1}\Lambda_n)$.*

Proof. *Case 1.* Suppose $i < i'$. We have $s_n(a_{n-1}a_n)$ is equal to

$$\begin{aligned} & s_ns_{n-1}s_{n-2} \cdots s_i \underline{s_n} s_{n-1} \cdots s_{i'} & \underline{(1)} \\ & \underline{s_ns_{n-1}s_n} s_{n-2} \cdots s_i s_{n-1} \cdots s_{i'} & \underline{(2)} \\ & s_{n-1}s_ns_{n-1}s_{n-2} \cdots s_i \underline{s_{n-1}} \cdots s_{i'} & \underline{(1)} \\ & s_{n-1}s_n \underline{s_{n-1}s_{n-2}s_{n-1}} \cdots s_i s_{n-2} \cdots s_{i'} & \underline{(2)} \\ & s_{n-1}s_n \underline{s_{n-2}s_{n-1}s_{n-2}} \cdots s_i s_{n-2} \cdots s_{i'} & \underline{(1)} \\ & s_{n-1}s_n \underline{s_{n-2}s_{n-1}s_{n-2}} \cdots s_i \underline{s_{n-2}} \cdots s_{i'}. \end{aligned}$$

We apply the same operations to the underlined letters $\underline{s_{n-2}} \cdots \underline{s_{i'}}$ in order to get $s_{n-1}s_{n-2} \cdots s_{i'-1}s_ns_{n-1} \cdots s_i$ which belongs to $\text{Span}(\Lambda_{n-1}\Lambda_n)$.

Case 2. Suppose $i \geq i'$. We have $s_n(a_{n-1}a_n)$ is equal to $s_ns_{n-1} \cdots s_i \underline{s_n} s_{n-1} \cdots s_{i'}$. We apply the operations (1) and (2) and get $s_{n-1}s_ns_{n-1} \cdots s_i \underline{s_{n-1}} \cdots s_{i'}$. Then we apply the same operations to $\underline{s_{n-1}}$ and get $s_{n-1}s_{n-2}s_ns_{n-1} \cdots s_i \underline{s_{n-2}} \cdots s_{i'}$. Since $i \geq i'$, we write $s_{i+2}s_{i+1}s_i$ in $\underline{s_{n-2}} \cdots \underline{s_{i'}}$ and get $s_{n-1}s_{n-2}s_ns_{n-1} \cdots s_i \underline{s_{n-2}} \cdots \underline{s_{i+2}} s_{i+1}s_i \cdots s_{i'}$. Similarly, we apply the same operations to $\underline{s_{n-2}}, \dots, \underline{s_{i+2}}$ and get

$$\begin{aligned} & s_{n-1} \cdots s_{i+1}s_n \cdots \underline{s_{i+1}s_i s_{i+1}} s_i \cdots s_{i'} & \underline{(2)} \\ & s_{n-1} \cdots s_{i+1}s_n \cdots s_{i+2} \underline{s_i} s_{i+1} s_i^2 s_{i-1} \cdots s_{i'} & \underline{(1)} \\ & s_{n-1} \cdots s_i s_n \cdots s_{i+2} s_{i+1} \underline{s_i^2} s_{i-1} \cdots s_{i'} & \underline{(3)} \\ & as_{n-1} \cdots s_i s_n \cdots s_{i+1} s_i s_{i-1} \cdots s_{i'} + s_{n-1} \cdots s_i s_n \cdots s_{i+2} s_{i+1} \underline{s_{i-1}} \cdots \underline{s_{i'}}. \end{aligned}$$

The first term belongs to $\text{Span}(\Lambda_{n-1}\Lambda_n)$. For the second term, the underlined letters commute with $s_n \cdots s_{i+2}s_{i+1}$ hence they are shifted to the left. We thus get $s_n(a_{n-1}a_n)$ is equal to $as_{n-1} \cdots s_i s_n \cdots s_{i'} + s_{n-1} \cdots s_{i'} s_n \cdots s_{i+1}$ which belongs to $\text{Span}(\Lambda_{n-1}\Lambda_n)$. \square

Lemma 5.9. *If $a_{n-1} = s_{n-1} \cdots s_i$ with $3 \leq i \leq n-1$ and $a_n = s_n \cdots s_3 t_k$ with $0 \leq k \leq e-1$, then $s_n(a_{n-1}a_n)$ belongs to $\text{Span}(\Lambda_{n-1}\Lambda_n)$.*

Proof. This corresponds to the case $i' = 3$ in the proof of Lemma 5.8 with a right multiplication by t_k for $0 \leq k \leq e-1$. Since $i \geq 3$, by the case $i \geq i'$ of Lemma 5.8, we have $s_n(a_{n-1}a_n) = as_{n-1} \cdots s_i s_n \cdots s_3 t_k + s_{n-1} \cdots s_3 s_n \cdots s_{i+1} t_k$. In the second term, t_k commutes with $s_n \cdots s_{i+1}$ hence it is shifted to the left. We get $s_n(a_{n-1}a_n) = as_{n-1} \cdots s_i s_n \cdots s_3 t_k + s_{n-1} \cdots s_3 t_k s_n \cdots s_{i+1}$ which belongs to $\text{Span}(\Lambda_{n-1}\Lambda_n)$. \square

Lemma 5.10. *If $a_{n-1} = s_{n-1} \cdots s_i$ with $3 \leq i \leq n-1$ and $a_n = s_n \cdots s_3 t_k s_2 s_3 \cdots s_{i'}$ with $1 \leq k \leq e-1$ and $2 \leq i' \leq n$, then $s_n(a_{n-1}a_n)$ belongs to $\text{Span}(\Lambda_{n-1}\Lambda_n)$.*

Proof. According to Lemma 5.9, we have $s_n(a_{n-1}a_n) = as_{n-1} \cdots s_i s_n \cdots s_3 t_k s_2 \cdots s_{i'} + s_{n-1} \cdots s_3 t_k s_n \cdots s_{i+1} s_2 \cdots s_{i'}$. The first term is an element of $\text{Span}(\Lambda_{n-1}\Lambda_n)$. We check that the second term also belongs to $\text{Span}(\Lambda_{n-1}\Lambda_n)$. Actually,

Case 1. Suppose $i' < i$. The second term is equal to $s_{n-1} \cdots s_3 t_k s_n \cdots s_{i+1} \underline{s_2} \cdots \underline{s_{i'}}$. The underlined letters commute with $s_n \cdots s_{i+1}$ and are shifted to the left. We get $s_{n-1} \cdots s_3 t_k s_2 \cdots s_{i'} s_n \cdots s_{i+1} \in \text{Span}(\Lambda_{n-1}\Lambda_n)$.

Case 2. Suppose $i' \geq i$. We write $s_{i-1}s_i s_{i+1}$ in $s_2 \cdots s_{i'}$ and get

$$\begin{aligned}
s_{n-1} \cdots s_3 t_k s_n \cdots s_{i+1} (s_2 \cdots s_{i'}) &= \\
s_{n-1} \cdots s_3 t_k s_n \cdots s_{i+1} (\underline{s_2} \cdots \underline{s_{i-1}} s_i s_{i+1} \cdots s_{i'}) &\stackrel{(1)}{=} \\
s_{n-1} \cdots s_3 t_k s_2 \cdots s_{i-1} s_n \cdots \underline{s_{i+1}} (\underline{s_i s_{i+1}} \cdots s_{i'}) &\stackrel{(2)}{=} \\
s_{n-1} \cdots s_3 t_k s_2 \cdots s_{i-1} s_n \cdots \underline{s_i s_{i+1}} s_i (s_{i+2} \cdots s_{i'}) &\stackrel{(1)}{=} \\
s_{n-1} \cdots s_3 t_k s_2 \cdots s_{i-1} s_i s_n \cdots s_{i+1} s_i (\underline{s_{i+2}} \cdots s_{i'}) &\stackrel{(1)}{=} \\
s_{n-1} \cdots s_3 t_k s_2 \cdots s_{i-1} s_i s_n \cdots \underline{s_{i+2} s_{i+1} s_{i+2}} s_i (s_{i+3} \cdots s_{i'}) &\stackrel{(2)}{=} \\
s_{n-1} \cdots s_3 t_k s_2 \cdots s_{i-1} s_i s_n \cdots \underline{s_{i+1}} s_{i+2} s_{i+1} s_i (s_{i+3} \cdots s_{i'}) &\stackrel{(1)}{=} \\
s_{n-1} \cdots s_3 t_k s_2 \cdots s_i s_{i+1} s_n \cdots s_{i+2} s_{i+1} s_i (\underline{s_{i+3}} \cdots \underline{s_{i'}}). &
\end{aligned}$$

We apply the same operations to the underlined letters $\underline{s_{i+3}}, \dots, \underline{s_{i'}}$. We finally get $s_{n-1} \cdots s_3 t_k s_2 \cdots s_{i'-1} s_n \cdots s_i \in \text{Span}(\Lambda_{n-1} \Lambda_n)$. \square

Lemma 5.11. *If $a_{n-1} = s_{n-1} \cdots s_3 t_k$ with $0 \leq k \leq e-1$ and $a_n = s_n \cdots s_i$ with $3 \leq i \leq n$, then $s_n(a_{n-1} a_n)$ belongs to $\text{Span}(\Lambda_{n-1} \Lambda_n)$.*

Proof. We have $s_n(a_{n-1} a_n)$ is equal to

$$\begin{aligned}
s_n s_{n-1} \cdots s_3 t_k \underline{s_n} \cdots s_i &\stackrel{(1)}{=} \\
\underline{s_n s_{n-1} s_n} \cdots s_3 t_k s_{n-1} \cdots s_i &\stackrel{(2)}{=} \\
s_{n-1} s_n s_{n-1} \cdots s_3 t_k \underline{s_{n-1}} \cdots s_i &\stackrel{(1)}{=} \\
s_{n-1} s_n \underline{s_{n-1} s_{n-2} s_{n-1}} \cdots s_3 t_k s_{n-2} \cdots s_i &\stackrel{(2)}{=} \\
s_{n-1} s_n \underline{s_{n-2} s_{n-1} s_{n-2}} \cdots s_3 t_k s_{n-2} \cdots s_i &\stackrel{(1)}{=} \\
s_{n-1} s_{n-2} s_n s_{n-1} \cdots s_3 t_k \underline{s_{n-2}} \cdots \underline{s_i}. &
\end{aligned}$$

Now we apply the same operations for $\underline{s_{n-2}}, \dots, \underline{s_i}$.

Case 1. Suppose $i = 3$. We get $s_{n-1} \cdots s_3 s_n s_{n-1} \cdots s_3 t_k s_3$. Next, we apply a braid relation to get $s_{n-1} \cdots s_3 s_n s_{n-1} \cdots \underline{t_k} s_3 t_k$, then we shift $\underline{t_k}$ to the left and we finally get $s_{n-1} \cdots s_3 t_k s_n s_{n-1} \cdots s_3 t_k$ which belongs to $\text{Span}(\Lambda_{n-1} \Lambda_n)$.

Case 2. Suppose $i > 3$. We directly get $s_{n-1} \cdots s_i s_{i-1} s_n \cdots s_3 t_k$ that also belongs to $\text{Span}(\Lambda_{n-1} \Lambda_n)$. \square

Lemma 5.12. *If $a_{n-1} = s_{n-1} \cdots s_3 t_k$ with $0 \leq k \leq e-1$ and $a_n = s_n \cdots s_3 t_l$ with $0 \leq l \leq e-1$, then $s_n(a_{n-1} a_n)$ belongs to $\text{Span}(\Lambda_{n-1} \Lambda_n)$.*

Proof. By Lemma 5.11, one can write $s_n(a_{n-1} a_n) = s_{n-1} \cdots s_3 t_k s_n s_{n-1} \cdots s_3 t_l$. By Lemma 5.2 where we compute $t_k t_l$, we directly deduce that $s_n(a_{n-1} a_n)$ belongs to $\text{Span}(\Lambda_{n-1} \Lambda_n)$. \square

Lemma 5.13. *If $a_{n-1} = s_{n-1} \cdots s_3 t_k$ with $0 \leq k \leq e-1$ and $a_n = s_n \cdots s_3 t_l s_2 \cdots s_i$ with $2 \leq i \leq n$ and $1 \leq l \leq e-1$, then $s_n(a_{n-1} a_n)$ belongs to $\text{Span}(\Lambda_{n-1}^* \Lambda_n)$.*

Proof. By the previous lemma, we have $s_n(a_{n-1} a_n) = s_{n-1} \cdots s_3 t_k s_n \cdots s_3 t_l t_l (s_2 \cdots s_i)$. By Lemma 5.2, the case $i = 2$ is obvious. Suppose $i \geq 3$. After replacing $t_k t_l$ by its value given in Lemma 5.2, we have two different terms in $s_n(a_{n-1} a_n)$ of the form $s_{n-1} \cdots s_3 t_k s_n \cdots s_3 t_x (s_2 \cdots s_i)$ with $0 \leq x \leq e-1$ and of the form $s_{n-1} \cdots s_3 t_k s_n \cdots s_3 t_x (s_3 \cdots s_i)$ with $0 \leq x \leq e-1$.

For terms of the form $s_{n-1} \cdots s_3 t_k s_n \cdots s_3 t_x (s_3 \cdots s_i)$ with $0 \leq x \leq e-1$, we have

$$\begin{aligned}
 s_{n-1} \cdots s_3 t_k s_n \cdots \underline{s_3 t_x (s_3 \cdots s_i)} & \stackrel{(2)}{=} \\
 s_{n-1} \cdots s_3 t_k s_n \cdots \underline{t_x s_3 t_x (s_4 \cdots s_i)} & \stackrel{(1)}{=} \\
 s_{n-1} \cdots s_3 t_k t_x s_n \cdots s_3 t_x (\underline{s_4 \cdots s_i}) & \stackrel{(1)}{=} \\
 s_{n-1} \cdots s_3 t_k t_x s_n \cdots \underline{s_4 s_3 s_4 t_x (s_5 \cdots s_i)} & \stackrel{(2)}{=} \\
 s_{n-1} \cdots s_3 t_k t_x s_n \cdots \underline{s_3 s_4 s_3 t_x (s_5 \cdots s_i)} & \stackrel{(1)}{=} \\
 s_{n-1} \cdots s_3 t_k t_x s_3 s_n \cdots s_3 t_x (\underline{s_5 \cdots s_i}).
 \end{aligned}$$

We apply the same operations for the underlined letters to get $s_{n-1} \cdots s_3 t_k t_x s_3 \cdots s_{i-1} \underline{s_n \cdots s_3 t_x}$ which belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$.

Consider the terms of the form $s_{n-1} \cdots s_3 t_k s_n \cdots s_3 t_x (s_2 \cdots s_i)$ with $0 \leq x \leq e-1$.

If $x \neq 0$, they belong to $\text{Span}(\Lambda_{n-1} \Lambda_n)$.

If $x = 0$, we have $s_{n-1} \cdots s_3 t_k s_n \cdots s_3 t_0 (s_2 s_3 \cdots s_i) \in R_0 s_{n-1} \cdots s_3 t_k s_n \cdots s_3 t_0 s_3 \cdots s_i + R_0 s_{n-1} \cdots s_3 t_k s_n \cdots s_4 s_3^2 s_4 \cdots s_i$. The first term corresponds to the previous case (with $x = 0$) and then belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$. By Lemma 5.7, the second term also belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$. \square

Lemma 5.14. *If $a_{n-1} = s_{n-1} \cdots s_3 t_k s_2 \cdots s_i$ with $2 \leq i \leq n-1$, $1 \leq k \leq e-1$ and $a_n = s_n \cdots s_{i'}$ with $3 \leq i' \leq n$, then $s_n(a_{n-1} a_n)$ belongs to $\text{Span}(\Lambda_{n-1} \Lambda_n)$.*

Proof. *Case 1.* Suppose $i < i'$. We have $s_n(a_{n-1} a_n)$ is equal to

$$\begin{aligned}
 s_n \cdots s_3 t_k s_2 \cdots s_i \underline{s_n \cdots s_{i'}} & \stackrel{(1)}{=} \\
 \underline{s_n s_{n-1} s_n} \cdots s_3 t_k s_2 \cdots s_i s_{n-1} \cdots s_{i'} & \stackrel{(2)}{=} \\
 s_{n-1} s_n s_{n-1} \cdots s_3 t_k s_2 \cdots s_i \underline{s_{n-1} \cdots s_{i'}} & \stackrel{(1)}{=} \\
 s_{n-1} s_n \underline{s_{n-1} s_{n-2} s_{n-1}} \cdots s_3 t_k s_2 \cdots s_i s_{n-2} \cdots s_{i'} & \stackrel{(2)}{=} \\
 s_{n-1} s_n \underline{s_{n-2} s_{n-1} s_{n-2}} \cdots s_3 t_k s_2 \cdots s_i s_{n-2} \cdots s_{i'} & \stackrel{(1)}{=} \\
 s_{n-1} s_n s_{n-2} s_{n-1} s_{n-2} \cdots s_3 t_k s_2 \cdots s_i \underline{s_{n-2} \cdots s_{i'+1} s_{i'}}.
 \end{aligned}$$

We apply the same operations to the underlined letters and we get $s_{n-1} \cdots s_{i'} s_n \cdots s_3 t_k s_2 \cdots s_i s_{i'}$.

Subcase 1.1. Suppose $i' = i+1$. We directly have $s_n(a_{n-1} a_n) \in \text{Span}(\Lambda_{n-1} \Lambda_n)$.

Subcase 1.2. Suppose $i' > i+1$. Then we write $s_{i'+1} s_{i'} s_{i'-1}$ in the underlined word of $s_{n-1} \cdots s_{i'} \underline{s_n \cdots s_3 t_k s_2} \cdots s_i s_{i'}$ and get

$$\begin{aligned}
 s_{n-1} \cdots s_{i'} s_n \cdots s_{i'+1} s_{i'} s_{i'-1} \cdots s_3 t_k s_2 \cdots s_i \underline{s_{i'}} & \stackrel{(1)}{=} \\
 s_{n-1} \cdots s_{i'} s_n \cdots s_{i'+1} \underline{s_{i'} s_{i'-1} s_{i'}} \cdots s_3 t_k s_2 \cdots s_i & \stackrel{(2)}{=} \\
 s_{n-1} \cdots s_{i'} s_n \cdots s_{i'+1} \underline{s_{i'-1} s_{i'} s_{i'-1}} \cdots s_3 t_k s_2 \cdots s_i & \stackrel{(1)}{=} \\
 s_{n-1} \cdots s_{i'-1} s_n \cdots s_3 t_k s_2 \cdots s_i \text{ which belongs to } \text{Span}(\Lambda_{n-1} \Lambda_n).
 \end{aligned}$$

Case 2. Suppose $i \geq i'$. We have $s_n(a_{n-1} a_n)$ is equal to $s_n \cdots s_3 t_k s_2 \cdots s_i \underline{s_n \cdots s_{i'}}$. We shift $\underline{s_n}$ to the left and apply a braid relation to get $s_{n-1} s_n s_{n-1} \cdots s_3 t_k s_2 \cdots s_i \underline{s_{n-1} \cdots s_{i'}}$. Write $s_{i+2} s_{i+1}$ in the underlined word to get $s_{n-1} s_n s_{n-1} \cdots s_3 t_k s_2 \cdots s_i \underline{s_{n-1} \cdots s_{i+2} s_{i+1} \cdots s_{i'}}$. We apply the same operations to the underlined letters to get

$$\begin{aligned}
 s_{n-1} \cdots s_{i+1} s_n \cdots s_3 t_k s_2 \cdots s_i \underline{s_i s_{i+1} s_i s_{i-1} \cdots s_{i'}} & \stackrel{(2)}{=} \\
 s_{n-1} \cdots s_{i+1} s_n \cdots s_3 t_k s_2 \cdots s_i \underline{s_{i+1} s_i s_{i+1} s_{i-1} \cdots s_{i'}} & \stackrel{(1)}{=}
 \end{aligned}$$

$s_{n-1} \cdots s_i s_n \cdots s_3 t_k s_2 \cdots s_{i-1} s_i s_{i+1} s_{i-1} \cdots s_{i'}$ (The details of the computation are left to the reader). Then we have

$s_{n-1} \cdots s_i s_n \cdots s_3 t_k s_2 \cdots s_{i-1} s_i s_{i+1} \underline{s_{i-1}} \cdots s_{i'}$ $\stackrel{(1)}{=}$
 $s_{n-1} \cdots s_i s_n \cdots s_3 t_k s_2 \cdots \underline{s_{i-1} s_i s_{i-1}} s_{i+1} s_{i-2} \cdots s_{i'}$ $\stackrel{(2)}{=}$
 $s_{n-1} \cdots s_i s_n \cdots s_3 t_k s_2 \cdots \underline{s_i s_{i-1} s_i s_{i+1} s_{i-2}} \cdots s_{i'}$ $\stackrel{(1)}{=}$
 $s_{n-1} \cdots s_{i-1} s_n \cdots s_3 t_k s_2 \cdots s_{i-2} s_{i-1} s_i s_{i+1} \underline{s_{i-2}} \cdots \underline{s_{i'+1}} s_{i'}$. We apply the same operations to the underlined letters and we finally get

$s_{n-1} \cdots s_{i'+1} s_n \cdots s_3 t_k \underline{s_2 s_3 \cdots s_i s_{i+1} s_{i'}}$. We write $s_{i'-1} s_{i'} s_{i'+1}$ in the underlined word and get $s_{n-1} \cdots s_{i'+1} s_n \cdots s_3 t_k s_2 s_3 \cdots s_{i'-1} s_{i'} s_{i'+1} \cdots s_i s_{i+1} \underline{s_{i'}}$. We shift $\underline{s_{i'}}$ to the left. We get $s_{n-1} \cdots s_{i'+1} s_n \cdots s_3 t_k s_2 s_3 \cdots s_{i'-1} \underline{s_{i'} s_{i'+1} s_{i'}}$. We apply a braid relation and get $s_{n-1} \cdots s_{i'+1} s_n \cdots s_3 t_k s_2 s_3 \cdots s_{i'-1} \underline{s_{i'+1} s_{i'} s_{i'+1}} \cdots s_i s_{i+1}$. Now we shift $\underline{s_{i'+1}}$ to the left and get

$$\begin{aligned}
 s_{n-1} \cdots s_{i'+1} s_n \cdots \underline{s_{i'+1} s_{i'} s_{i'+1}} s_{i'-1} \cdots s_3 t_k s_2 \cdots s_{i+1} & \stackrel{(2)}{=} \\
 s_{n-1} \cdots s_{i'+1} s_n \cdots \underline{s_{i'} s_{i'+1} s_{i'} s_{i'-1}} \cdots s_3 t_k s_2 \cdots s_{i+1} & \stackrel{(1)}{=} \\
 s_{n-1} \cdots s_{i'+1} s_{i'} s_n \cdots s_{i'+1} s_{i'} s_{i'-1} \cdots s_3 t_k s_2 \cdots s_{i+1} & = \\
 s_{n-1} \cdots s_{i'} s_n \cdots s_3 t_k s_2 \cdots s_{i+1}, & \text{ which belongs to } \text{Span}(\Lambda_{n-1} \Lambda_n). \quad \square
 \end{aligned}$$

Lemma 5.15. *If $a_{n-1} = s_{n-1} \cdots s_3 t_k s_2 \cdots s_i$ with $2 \leq i \leq n-1$, $1 \leq k \leq e-1$ and $a_n = s_n s_{n-1} \cdots s_3 t_l$ with $0 \leq l \leq e-1$, then $s_n(a_{n-1} a_n)$ belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$.*

Proof. By the final result of the computations in Lemma 5.14, we have $s_n(a_{n-1} a_n)$ is equal to $s_{n-1} \cdots s_3 s_n \cdots s_3 t_k s_2 s_3 \cdots s_{i+1} \underline{t_l}$. We shift $\underline{t_l}$ to the left and get $s_{n-1} \cdots s_3 s_n \cdots s_3 t_k s_2 s_3 t_l \cdots s_{i+1}$. By Case 5 of Proposition 5.6, we have to deal with the following two terms:

- $s_{n-1} \cdots s_3 s_n \cdots s_3 t_x s_3 t_l s_4 \cdots s_{i+1}$ and
- $s_{n-1} \cdots s_3 s_n \cdots s_3 t_x s_3 t_l s_3 s_4 \cdots s_{i+1}$ with $1 \leq x, l \leq e-1$.

The first term is of the form

$$\begin{aligned}
 s_{n-1} \cdots s_3 s_n \cdots \underline{s_3 t_x s_3 t_l} s_4 \cdots s_{i+1} & \stackrel{(2)}{=} \\
 s_{n-1} \cdots s_3 s_n \cdots s_4 \underline{t_x} s_3 t_l s_4 \cdots s_{i+1} & \stackrel{(1)}{=} \\
 s_{n-1} \cdots s_3 t_x s_n \cdots s_3 t_x \underline{s_4} \cdots s_{i+1} & \stackrel{(1)}{=} \\
 s_{n-1} \cdots s_3 t_x s_n \cdots \underline{s_4 s_3 s_4} t_x t_l s_5 \cdots s_{i+1} & \stackrel{(2)}{=} \\
 s_{n-1} \cdots s_3 t_x s_n \cdots \underline{s_3 s_4 s_3} t_x t_l s_5 \cdots s_{i+1} & \stackrel{(1)}{=} \\
 s_{n-1} \cdots s_3 t_x s_3 s_n \cdots s_4 s_3 t_x t_l \underline{s_5} \cdots \underline{s_{i+2}} s_{i+1}. & \text{ We apply the same operations to the underlined letters, we get} \\
 s_{n-1} \cdots s_3 t_x s_3 \cdots s_{i-1} s_n \cdots s_4 s_3 t_x t_l \underline{s_{i+1}}. & \text{ Finally, we shift } \underline{s_{i+1}} \text{ to the left and get } s_{n-1} \cdots s_3 t_x s_3 \cdots s_i s_n \cdots \\
 s_4 s_3 t_x t_l. & \text{ Since } 2 \leq i \leq n-1 \text{ and by the computation of } t_x t_l \text{ in Lemma 5.2, the lemma is satisfied for this case.}
 \end{aligned}$$

The second term is equal to

$$\begin{aligned}
 s_{n-1} \cdots s_3 s_n \cdots \underline{s_3 t_x s_3} t_l s_3 s_4 \cdots s_{i+1} & \stackrel{(2)}{=} \\
 s_{n-1} \cdots s_3 s_n \cdots \underline{t_x} s_3 t_l s_3 s_4 \cdots s_{i+1} & \stackrel{(1)}{=} \\
 s_{n-1} \cdots s_3 t_x s_n \cdots s_3 t_x t_l s_3 s_4 \cdots s_{i+1}. & \text{ We replace } t_x t_l \text{ by its value given in Lemma 5.2, we get terms of the} \\
 \text{three following forms:} &
 \end{aligned}$$

- $s_{n-1} \cdots s_3 t_x s_n \cdots s_3 t_m t_0 s_3 s_4 \cdots s_{i+1}$ with $1 \leq m \leq e-1$,
- $s_{n-1} \cdots s_3 t_x s_n \cdots s_3 t_m s_3 s_4 \cdots s_{i+1}$ with $0 \leq m \leq e-1$,
- $s_{n-1} \cdots s_3 t_x s_n \cdots s_4 s_3^2 s_4 \cdots s_{i+1}$.

The first term belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$. The third term also belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$. This is done by using the computation in the proof of Lemma 5.7. For the second term, we have

$$s_{n-1} \cdots s_3 t_x s_n \cdots s_4 \underline{s_3 t_m s_3} s_4 \cdots s_{i+1} \quad \stackrel{(2)}{=}$$

$$s_{n-1} \cdots s_3 t_x s_n \cdots \underline{s_4 t_m} s_3 t_m \underline{s_4} \cdots s_{i+1} \quad \stackrel{(1)}{=}$$

$$s_{n-1} \cdots s_3 t_x t_m s_n \cdots \underline{s_4 s_3 s_4} t_m s_5 \cdots s_{i+1} \quad \stackrel{(2)}{=}$$

$$s_{n-1} \cdots s_3 t_x t_m s_n \cdots \underline{s_3 s_4 s_3} t_m s_5 \cdots s_{i+1} \quad \stackrel{(1)}{=}$$

$s_{n-1} \cdots s_3 t_x t_m s_n \cdots s_4 s_3 t_m \underline{s_5} \cdots \underline{s_{i+2}} s_{i+1}$. We apply the same operations to the underlined letters and get $s_{n-1} \cdots s_3 t_x t_m s_3 \cdots s_{i-1} s_n \cdots s_3 t_m s_{i+1}$. Now we shift $\underline{s_{i+1}}$ to the left and finally get $s_{n-1} \cdots s_3 t_x t_m s_3 \cdots s_i s_n \cdots s_3 t_m$ with $2 \leq i \leq n-1$ which belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$.

Note that for $i = 2$, we get terms that are equal to the itemized terms (given at the beginning of this proof) after replacing $s_4 \cdots s_{i+1}$ by 1. \square

Lemma 5.16. *If $a_{n-1} = s_{n-1} \cdots s_3 t_k s_2 \cdots s_i$ with $2 \leq i \leq n-1$, $1 \leq k \leq e-1$ and $a_n = s_n \cdots s_3 t_l s_2 \cdots s_{i'}$ with $1 \leq l \leq e-1$ and $2 \leq i' \leq n$, then $s_n(a_{n-1} a_n)$ belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$.*

Proof. According to the computation in the proof of Lemma 5.15, we get the following possible terms. They appear in the proof of Lemma 5.15 in the following order.

1. $s_n \cdots s_3 t_x t_l$ with $0 \leq x, l \leq e-1$,
2. $s_n \cdots s_3 t_m t_0 s_3 \cdots s_{i+1}$ with $1 \leq m \leq e-1$,
3. $s_n \cdots s_3 t_m$ with $0 \leq m \leq e-1$,
4. $s_n \cdots s_4 s_3^2 s_4 \cdots s_{i+1}$

We show that the product on the right by $s_2 \cdots s_{i'}$ of each of the previous terms belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$.

Case 1. Consider the first term $s_n \cdots s_3 t_x t_l (s_2 \cdots s_{i'})$ with $0 \leq x, l \leq e-1$. We replace $t_x t_l$ by its decomposition given in Lemma 5.2, we get these terms

- $s_n \cdots s_3 t_m t_0 s_3 \cdots s_{i'}$ with $1 \leq m \leq e-1$,
- $s_n \cdots s_3 t_m s_3 \cdots s_{i'}$ with $0 \leq m \leq e-1$,
- $s_n \cdots s_4 s_3^2 s_4 \cdots s_{i'}$.

The first term belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$. The third one is done in Lemma 5.7. For the second one, we have

$$s_n \cdots \underline{s_3 t_m s_3} \cdots s_{i'} \quad \stackrel{(2)}{=}$$

$$s_n \cdots \underline{t_m} s_3 t_m \underline{s_4} \cdots s_{i'} \quad \stackrel{(1)}{=}$$

$$t_m s_n \cdots \underline{s_4 s_3 s_4} t_m s_5 \cdots s_{i'} \quad \stackrel{(2)}{=}$$

$$t_m s_n \cdots \underline{s_3 s_4 s_3} t_m s_5 \cdots s_{i'} \quad \stackrel{(1)}{=}$$

$t_m s_3 s_n \cdots s_4 s_3 t_m \underline{s_5} \cdots \underline{s_{i'-1}} s_{i'}$. We apply the same operations to $\underline{s_5}, \dots, \underline{s_{i'-1}}$ and get $t_m s_3 \cdots s_{i'-2} s_n \cdots s_3 t_m \underline{s_{i'}}$. We shift $\underline{s_{i'}}$ to the left and finally get $t_m s_3 \cdots s_{i'-1} \underline{s_n \cdots s_3 t_m}$ which belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$.

We now consider *Case 3* because we use the computation we made in Case 1. In this case, the term is of the form $s_n \cdots s_3 t_m (s_2 \cdots s_{i'})$ with $0 \leq m \leq e-1$. If $m \neq 0$, then it belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$. If $m = 0$, we get two terms $s_n \cdots s_3 t_0 s_3 \cdots s_{i'}$ and $s_n \cdots s_4 s_3^2 s_4 \cdots s_{i'}$. The first term is done in Case 1. The second term is done in Lemma 5.7.

Consider *Case 4*. We replace $s_n \cdots s_4 s_3^2 s_4 \cdots s_{i+1}$ by its decomposition given by the computation in the proof of Lemma 5.7. We multiply each term of the decomposition by $s_2 \cdots s_{i'}$ on the right and we prove that it belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$ in the same way as the proof of Lemma 5.10.

Finally, it remains to show that the term corresponding to *Case 2* belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$. It is of the form

$$s_n \cdots s_3 t_m t_0 s_3 \cdots s_{i+1} (s_2 \cdots s_{i'}) \text{ with } 1 \leq m \leq e-1.$$

Suppose $i' \leq i$. We have

$$\begin{aligned} s_n \cdots s_3 t_m s_2 s_3 \cdots s_{i+1} \underline{s_2} \cdots s_{i'} & \quad \underline{(1)} \\ s_n \cdots s_3 t_m \underline{s_2 s_3 s_2} s_4 \cdots s_{i+1} s_3 \cdots s_{i'} & \quad \underline{(2)} \\ s_n \cdots \underline{s_3 t_m s_3} s_2 s_3 s_4 \cdots s_{i+1} s_3 \cdots s_{i'} & \quad \underline{(2)} \\ s_n \cdots \underline{t_m} s_3 t_m s_2 s_3 s_4 \cdots s_{i+1} s_3 \cdots s_{i'} & \quad \underline{(1)} \\ t_m s_n \cdots s_3 t_m s_2 s_3 s_4 \cdots s_{i+1} \underline{s_3} \cdots s_{i'} & \quad \underline{(1)} \\ t_m s_n \cdots s_3 t_m s_2 \underline{s_3 s_4 s_3} s_5 \cdots s_{i+1} s_4 \cdots s_{i'} & \quad \underline{(2)} \\ t_m s_n \cdots s_3 t_m s_2 \underline{s_4 s_3 s_4} s_5 \cdots s_{i+1} s_4 \cdots s_{i'} & \quad \underline{(1)} \\ t_m s_n \cdots \underline{s_4 s_3 s_4} t_m s_2 s_3 s_4 s_5 \cdots s_{i+1} s_4 \cdots s_{i'} & \quad \underline{(2)} \\ t_m s_n \cdots \underline{s_3 s_4 s_3} t_m s_2 s_3 s_4 s_5 \cdots s_{i+1} s_4 \cdots s_{i'} & \quad \underline{(1)} \\ t_m s_3 s_n \cdots s_4 s_3 t_m s_2 s_3 s_4 s_5 \cdots s_{i+1} \underline{s_4} \cdots \underline{s_{i'-1}} \underline{s_{i'}}. & \end{aligned}$$

We apply the same operations to $\underline{s_4}, \dots, \underline{s_{i'-1}}$ to get $t_m s_3 \cdots s_{i'-1} s_n \cdots s_3 t_m s_2 s_3 \cdots s_i s_{i+1} \underline{s_{i'}}$. We shift $\underline{s_{i'}}$ to the left and finally get $t_m s_3 \cdots s_{i'} s_n \cdots s_3 t_m s_2 s_3 \cdots s_i s_{i+1}$. Since $i' \leq i \leq n-1$, this term satisfies the property of the lemma.

Suppose $i' > i$. As previously, we have

$$\begin{aligned} s_n \cdots s_3 t_m s_2 s_3 \cdots s_{i+1} \underline{s_2} \cdots s_{i'} & \quad \underline{(1)} \\ t_m s_n \cdots s_3 t_m s_2 s_3 \cdots s_{i+1} \underline{s_3} \cdots s_{i'} & \quad \underline{(1)} \\ t_m s_3 s_n \cdots s_3 t_m s_2 s_3 \cdots s_{i+1} s_4 \cdots s_{i'}. & \end{aligned}$$

Now we write $s_i s_{i+1}$ in $s_4 \cdots s_{i'}$ and get $t_m s_3 s_n \cdots s_3 t_m s_2 s_3 \cdots s_{i+1} \underline{s_4} \cdots \underline{s_i} s_{i+1} \cdots s_{i'}$. We apply the same operations to $\underline{s_4}, \dots, \underline{s_i}$ to get $t_m s_3 \cdots s_i s_n \cdots s_3 t_m s_2 s_3 \cdots s_i s_{i+1}^2 s_{i+2} \cdots s_{i'}$. Applying a quadratic relation, we finally get $at_m s_3 \cdots s_i s_n \cdots s_3 t_m s_2 s_3 \cdots s_{i'} + t_m s_3 \cdots s_i s_n \cdots s_3 t_m s_2 s_3 \cdots s_i s_{i+2} \cdots s_{i'}$. The first term satisfies the property of the lemma.

For the second term, we write $s_{i+2} s_{i+1}$ in $\underline{s_n \cdots s_3}$ of $t_m s_3 \cdots s_i \underline{s_n \cdots s_3} t_m s_2 s_3 \cdots s_i s_{i+2} \cdots s_{i'}$ and get

$$\begin{aligned} t_m s_3 \cdots s_i s_n \cdots s_{i+2} s_{i+1} \cdots s_3 t_m s_2 s_3 \cdots s_i \underline{s_{i+2}} \cdots s_{i'} & \quad \underline{(1)} \\ t_m s_3 \cdots s_i s_n \cdots \underline{s_{i+2} s_{i+1} s_{i+2}} \cdots s_3 t_m s_2 s_3 \cdots s_i s_{i+3} \cdots s_{i'} & \quad \underline{(2)} \\ t_m s_3 \cdots s_i s_n \cdots \underline{s_{i+1} s_{i+2} s_{i+1}} \cdots s_3 t_m s_2 s_3 \cdots s_i s_{i+3} \cdots s_{i'} & \quad \underline{(1)} \\ t_m s_3 \cdots s_{i+1} s_n \cdots s_3 t_m s_2 s_3 \cdots s_i \underline{s_{i+3}} \cdots \underline{s_{i'}}. & \end{aligned}$$

We apply the same operations to $\underline{s_{i+3}}, \dots, \underline{s_{i'}}$ and finally get $t_m s_3 \cdots s_{i'-1} s_n \cdots s_3 t_m s_2 s_3 \cdots s_i$. Since $i' - 1 \leq n - 1$, this term belongs to $\text{Span}(S_{n-1}^* \Lambda_n)$. \square

Remark 5.17. Our basis never coincides with the Ariki basis (see Proposition 1.6 (2) in [1]) for the Hecke algebra associated with $G(e, e, n)$. For example, consider the element $t_1 t_0 t_0$ which belongs to Ariki's basis. In our basis, it is equal to the linear combination $at_1 t_0 + t_1$, where $t_1 t_0$ and t_1 are two distinct elements of our basis.

The general hope would be to construct natural bases for $H(de, e, n)$ that can be defined from geodesic normal forms in the associated complex reflection groups $G(de, e, n)$, which was established in Theorem 5.1 for the case of $H(e, e, n)$. This arises the question whether the geodesic normal forms established in Section 3 by using the presentation of Corran–Lee–Lee of $G(de, e, n)$ for $d > 1$ also provide natural bases for the corresponding Hecke algebras. Unfortunately, some intricate arguments in our proof for the case of $H(e, e, n)$ do not work in the general case of $H(de, e, n)$. Actually, on the one hand, the Hecke algebra is defined, in the general case, as a quotient of the complex braid group algebra defined by the presentation of Corran–Lee–Lee, where the generators t_i 's are defined over \mathbb{Z} , but on the other hand our geodesic normal forms are attached to the presentation of Corran–Lee–Lee of the complex reflection group $G(de, e, n)$, where the generators t_i 's are defined this time over $\mathbb{Z}/de\mathbb{Z}$. This phenomenon does not occur in the case of the presentations of Corran–Picantin of $G(e, e, n)$ and $B(e, e, n)$ used to construct a basis for $H(e, e, n)$. Nonetheless, for $e = 1$, the Hecke algebra $H(d, 1, n)$ is defined by using the classical presentation of the complex braid group $B(d, 1, n)$. Using the geodesic normal forms constructed in Subsection 3.3, we are also able to provide a natural basis for $H(d, 1, n)$. This will be established in the next section.

6. The case of $H(d, 1, n)$

Let $d > 1$ and $n \geq 2$. Let $R_0 = \mathbb{Z}[a, b_1, b_2, \dots, b_{d-1}]$. Recall that the Hecke algebra $H(d, 1, n)$ is defined as the unitary associative R_0 -algebra generated by the elements z, s_2, s_3, \dots, s_n with the following relations:

1. $zs_2zs_2 = s_2zs_2z$,
2. $zs_j = s_jz$ for $2 \leq j \leq n$,
3. $s_i s_{i+1} s_i = s_{i+1} s_i s_{i+1}$ for $2 \leq i \leq n-1$,
4. $s_i s_j = s_j s_i$ for $|i - j| > 1$,
5. $z^d - b_1 z^{d-1} - b_2 z^{d-2} - \dots - b_{d-1} z - 1 = 0$ and $s_j^2 - a s_j - 1 = 0$ for $2 \leq j \leq n$.

Using the geodesic normal forms introduced in subsection 3.3 for all the elements of $G(d, 1, n)$, we construct a basis for $H(d, 1, n)$ that is different from the one defined by Ariki–Koike in [2].

Let us introduce the following subsets of $H(d, 1, n)$.

$$\Lambda_1 = \{z^k \text{ for } 0 \leq k \leq d-1\},$$

and for $2 \leq i \leq n$,

$$\Lambda_i = \begin{cases} 1, & \\ s_i \cdots s_{i'} & \text{for } 2 \leq i' \leq i, \\ s_i \cdots s_2 z^k & \text{for } 1 \leq k \leq d-1, \\ s_i \cdots s_2 z^k s_2 \cdots s_{i'} & \text{for } 1 \leq k \leq d-1 \text{ and } 2 \leq i' \leq i. \end{cases}$$

Define $\Lambda = \Lambda_1 \Lambda_2 \cdots \Lambda_n$ to be the set of the products $a_1 a_2 \cdots a_n$, where $a_1 \in \Lambda_1, \dots, a_n \in \Lambda_n$. Remark that this set corresponds to all the reduced words $RE(w)$ of the form $RE_1(w) RE_2(w) RE_3(w) \cdots RE_n(w)$ introduced in subsection 3.3 (see Definition 3.20). In this section, we establish the following theorem.

Theorem 6.1. *The set Λ provides an R_0 -basis of the Hecke algebra $H(d, 1, n)$.*

We have $|\Lambda_1| = d$ and $|\Lambda_i| = id$ for $2 \leq i \leq n$. Then $|\Lambda|$ is equal to $d^n n!$ that is the order of $G(d, 1, n)$. Hence by Proposition 2.3(i) of [14], it is sufficient to show that Λ is an R_0 -generating set of $H(d, 1, n)$. This is proved by induction on n in much the same way as Theorem 5.1. We provide the following preliminary lemmas that are useful in the proof of the theorem.

Lemma 6.2. For $1 \leq k \leq d-1$, the element $(s_2zs_2)^k$ belongs to $\sum_{\substack{\lambda_1 \in \Lambda_1, \\ \lambda_2 \in \Lambda'_2}} R_0\lambda_1\lambda_2$, where $\Lambda'_2 = \{1, s_2, s_2z, s_2z^2, \dots, s_2z^{k-1}, s_2zs_2, s_2z^2s_2, \dots, s_2z^ks_2\}$.

Proof. The property is clear for $k = 1$. Let $k = 2$. We have $(s_2zs_2)^2 = s_2zs_2^2zs_2$. We apply a quadratic relation and get $as_2zs_2zs_2 + s_2z^2s_2$. Applying a braid relation, one gets $azs_2zs_2^2 + s_2z^2s_2$. Using a quadratic relation, this is equal to $a^2zs_2zs_2 + azs_2z + s_2z^2s_2$, where each term is of the form $\lambda_1\lambda_2$ with $\lambda_1 \in \Lambda_1$ and $\lambda_2 \in \{1, s_2, s_2z, s_2zs_2, s_2z^2s_2\}$.

Let $k \geq 3$. Suppose the property is satisfied for $(s_2zs_2)^3, \dots$, and $(s_2zs_2)^{k-1}$. We have $(s_2zs_2)^k = (s_2zs_2)^{k-1}(s_2zs_2)$. By the induction hypothesis, the terms that appear in the decomposition of $(s_2zs_2)^{k-1}$ are of the following forms.

- z^c with $0 \leq c \leq d-1$,
- $z^cs_2z^{c'}$ with $0 \leq c \leq d-1$ and $0 \leq c' \leq k-2$,
- $z^cs_2z^{c'}s_2$ with $0 \leq c \leq d-1$ and $1 \leq c' \leq k-1$.

Multiplying these terms by s_2zs_2 on the right, we get the following 3 cases.

Case 1: A term of the form $z^cs_2zs_2$ with $0 \leq c \leq d-1$. It is of the form $\lambda_1\lambda_2$ with $\lambda_1 \in \Lambda_1$ and $\lambda_2 \in \Lambda'_2$.

Case 2: A term of the form $z^cs_2z^{c'}s_2zs_2$ with $0 \leq c \leq d-1$ and $0 \leq c' \leq k-2$. We shift $z^{c'}$ to the right by applying braid relations and get $z^cs_2^2zs_2z^{c'}$. Applying a quadratic relation, this is equal to $az^cs_2zs_2z^{c'} + z^{c+1}s_2z^{c'}$. Now we shift $z^{c'}$ to the left by applying braid relations and get $az^{c+c'}s_2zs_2 + z^{c+1}s_2z^{c'}$. Each term is of the form $\lambda_1\lambda_2$ with $\lambda_1 \in \Lambda_1$ and $\lambda_2 \in \Lambda'_2$.

Case 3: A term of the form $z^cs_2z^{c'}s_2^2zs_2$ with $0 \leq c \leq d-1$ and $1 \leq c' \leq k-1$. By applying a quadratic relation, we have $z^cs_2z^{c'}s_2^2zs_2 = az^cs_2z^{c'}s_2zs_2 + z^cs_2z^{c'+1}s_2$. The first term is the same as in the previous case. Then both terms are of the form $\lambda_1\lambda_2$ with $\lambda_1 \in \Lambda_1$ and $\lambda_2 \in \Lambda'_2$. \square

Lemma 6.3. For $1 \leq k \leq d-1$, the element $(s_2zs_2)^ks_2$ belongs to $R_0(s_2zs_2)^k + R_0z(s_2zs_2)^{k-1} + \dots + R_0z^{k-1}(s_2zs_2) + R_0s_2z^k$.

Proof. For $k = 1$, we have $(s_2zs_2)s_2 = s_2zs_2^2 = as_2zs_2 + s_2z$. Then the property is satisfied for $k = 1$. Let $k \geq 2$. Suppose that the property is satisfied for $(s_2zs_2)^{k-1}$. We have $(s_2zs_2)^ks_2 = (s_2zs_2)(s_2zs_2)^{k-1}s_2$. By the induction hypothesis, it belongs to $R_0(s_2zs_2)(s_2zs_2)^{k-1} + R_0(s_2zs_2)z(s_2zs_2)^{k-2} + \dots + R_0(s_2zs_2)z^{k-2}(s_2zs_2) + R_0(s_2zs_2)s_2z^{k-1}$. Then it belongs to $R_0(s_2zs_2)^k + R_0z(s_2zs_2)^{k-1} + \dots + R_0z^{k-2}(s_2zs_2)^2 + R_0z^{k-1}(s_2zs_2) + R_0s_2z^k$. It follows that for all $1 \leq k \leq d-1$, the element $(s_2zs_2)^ks_2$ belongs to $R_0(s_2zs_2)^k + R_0z(s_2zs_2)^{k-1} + \dots + R_0z^{k-1}(s_2zs_2) + R_0s_2z^k$. \square

Lemma 6.4. For $1 \leq k \leq d-1$, the element $s_2z^ks_2$ belongs to $\sum_{\substack{\lambda_1 \in \Lambda_1, \\ \lambda_2 \in \Lambda''_2}} R_0\lambda_1\lambda_2$, where $\Lambda''_2 = \{1, s_2, s_2z, s_2z^2, \dots, s_2z^{k-1}, s_2zs_2, (s_2zs_2)^2, \dots, (s_2zs_2)^k\}$.

Proof. The lemma is satisfied for $k = 1$. For $k = 2$, we have $s_2z^2s_2 = s_2zs_2^{-1}s_2zs_2$. Using that $s_2^{-1} = s_2 - a$, we get $s_2zs_2^2s_2 - as_2zs_2zs_2$. Now we apply a braid relation then a quadratic relation and get $(s_2zs_2)^2 - azs_2zs_2^2 = (s_2zs_2)^2 - a^2zs_2zs_2 - azs_2z$ which satisfies the property we are proving.

Suppose the property is satisfied for $s_2z^{k-1}s_2$. We have $s_2z^ks_2 = s_2z^{k-1}s_2^{-1}s_2zs_2 = s_2z^{k-1}s_2zs_2zs_2 - as_2z^{k-1}s_2zs_2$ by replacing s_2^{-1} by $s_2 - a$. For the second term, we shift z^{k-1} to the right and get $s_2^2zs_2z^{k-1}$. We apply a quadratic relation to get $as_2zs_2z^{k-1} + zs_2z^{k-1}$ then we shift z^{k-1} to the left and finally get $az^{k-1}s_2zs_2 + zs_2z^{k-1}$, where each term is of the form $\lambda_1\lambda_2$ with $\lambda_1 \in \Lambda_1$ and $\lambda_2 \in \Lambda''_2$.

For the first term $s_2 z^{k-1} s_2 s_2 z s_2$, by the induction hypothesis, the terms that appear in the decomposition of $s_2 z^{k-1} s_2$ are of the following forms.

- z^c and $z^c s_2$ with $0 \leq c \leq d-1$,
- $z^c (s_2 z s_2)^{c'}$ with $0 \leq c \leq d-1$ and $1 \leq c' \leq k-1$,
- $z^c s_2 z^{c'}$ with $0 \leq c \leq d-1$ and $1 \leq c' \leq k-2$.

Multiplying these terms by $s_2 z s_2$ on the right, we get the following 3 cases.

Case 1. We have $z^c s_2 z s_2$ and $z^c s_2 s_2 z s_2 = a z^c s_2 z s_2 + z^{c+1} s_2$, where each term in both expressions is of the form $\lambda_1 \lambda_2$ with $\lambda_1 \in \Lambda_1$ and $\lambda_2 \in \Lambda_2''$.

Case 2. The term $z^c (s_2 z s_2)^{c'} s_2 z s_2 = z^c (s_2 z s_2)^{c'+1}$ is of the form $\lambda_1 \lambda_2$ with $\lambda_1 \in \Lambda_1$ and $\lambda_2 \in \Lambda_2''$ since $1 \leq c' \leq k-1$.

Case 3. We have $z^c s_2 z^{c'} s_2 z s_2 = z^c s_2^2 z s_2 z^{c'} = a z^c s_2 z s_2 z^{c'} + z^{c+1} s_2 z^{c'}$. The first term is equal to $z^{c+c'} s_2 z s_2$ and the second term is equal to $z^{c+1} s_2 z^{c'}$ with $1 \leq c' \leq k-2$. Both are of the form $\lambda_1 \lambda_2$ with $\lambda_1 \in \Lambda_1$ and $\lambda_2 \in \Lambda_2''$. \square

The following proposition ensures that the case $n = 2$ of Theorem 6.1 works properly.

Proposition 6.5. *For all $a_1 \in \Lambda_1$ and $a_2 \in \Lambda_2$, the elements $z a_1 a_2$ and $s_2 a_1 a_2$ belong to $\text{Span}(\Lambda_1 \Lambda_2)$.*

Proof. It is readily checked that $z a_1 a_2$ belongs to $\text{Span}(\Lambda_1 \Lambda_2)$. Note that when the power of z exceeds $d-1$, we use Relation 1 of Definition 4.5.

It is easily checked that if $a_1 \in \Lambda_1$ and $a_2 = 1$, the element $s_2 a_1 a_2$ belongs to $\text{Span}(\Lambda_1 \Lambda_2)$. Also, when $a_1 = 1$ and $a_2 \in \Lambda_2$, we have that $s_2 a_1 a_2$ belongs to $\text{Span}(\Lambda_1 \Lambda_2)$.

Suppose $a_1 = z^k$ with $1 \leq k \leq d-1$ and $a_2 = s_2$. We have $s_2 a_1 a_2$ is equal to $s_2 z^k s_2$. Hence it belongs to $\text{Span}(\Lambda_1 \Lambda_2)$.

Suppose $a_1 = z^k$ with $1 \leq k \leq d-1$ and $a_2 = s_2 z^{k'}$ with $1 \leq k' \leq d-1$. We have $s_2 a_1 a_2 = s_2 z^k s_2 z^{k'}$. We replace $s_2 z^k s_2$ by its decomposition given in Lemma 6.4, then we use the result of Lemma 6.2 to directly deduce that $s_2 z^k s_2 z^{k'}$ belongs to $\text{Span}(\Lambda_1 \Lambda_2)$.

Finally, suppose $a_1 = z^k$ with $1 \leq k \leq d-1$ and $a_2 = s_2 z^{k'} s_2$ with $1 \leq k' \leq d-1$. We have $s_2 a_1 a_2$ is equal to $s_2 z^k s_2 z^{k'} s_2$. We replace $s_2 z^k s_2$ by its decomposition given in Lemma 6.4. Then by the results of Lemmas 6.3 and 6.2, we deduce that $s_2 z^k s_2 z^{k'} s_2$ belongs to $\text{Span}(\Lambda_1 \Lambda_2)$. \square

As in the previous section, in order to prove Theorem 6.1, we have to proceed as in Lemmas 5.7 to 5.16. This is established in Lemmas 4.3.6 to 4.3.15 in [17]. Along with Proposition 6.5, this provides an inductive proof of Theorem 6.1 that is similar to the proof of Theorem 5.1. We conclude this section by the following remark.

Remark 6.6. For every d and n at least equal to 2, our basis never coincides with the Ariki–Koike basis (see Theorem 3.10 in [2]) as illustrated by the following example. Consider the element $s_2 z s_2^2 = s_2 z s_2 s_2$ which belongs to the Ariki–Koike basis. In our basis, it is equal to the linear combination $a s_2 z s_2 + s_2 z$, where $s_2 z s_2$ and $s_2 z$ are two distinct elements of our basis.

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