

Multi-loop Integrand Reduction with Computational Algebraic Geometry

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Abstract. We discuss recent progress in multi-loop integrand reduction methods. Motivated by the possibility of an automated construction of multi-loop amplitudes via generalized unitarity cuts we describe a procedure to obtain a general parameterisation of any multi-loop integrand in a renormalizable gauge theory. The method relies on computational algebraic geometry techniques such as Gröbner bases and primary decomposition of ideals. We present some results for two and three loop amplitudes obtained with the help of the MACAULAY2 computer algebra system and the Mathematica package BASISDET.

1. Introduction

There has been huge progress recently in the programme to fully automate one-loop amplitude computations. Much of this success has been possible thanks to the development of new reductions methods such as unitarity [1, 2], generalized unitarity [3] and integrand reduction [4] which allow efficient generation of one-loop amplitudes from simpler tree-level building blocks. We refer to G. Ossola's contribution to these proceedings for more details of the status of this field [5].

There has been some interest in extending these techniques to multi-loop amplitudes in the hope of improving our predictions of multi-particle production beyond $2 \rightarrow 2$ scattering processes. Thanks to developments in infra-red subtraction techniques full NNLO QCD computations with coloured initial and final states have been now been achieved for $pp \rightarrow H + j$ [6], $pp \rightarrow t\bar{t}$ [7] and $pp \rightarrow 2j$ [8] which give hope that a similar level of precision prediction may also be possible for $2 \rightarrow 3$ processes in the near future. In order to combat the rapid growth of the Feynman diagrams in such computations, on-shell techniques only work with physical degrees of freedom and simplify intermediate steps.

The traditional approach to multi-loop Feynman diagram computations using integration-by-parts identities has been well studied since its proposal more than 30 years ago [9]. Several automated public codes have been written and new developments are constantly increasing the range of applicability, see for example the contribution of R. Lee to these proceedings [10, 11]. Nevertheless complete solutions for multi-scale problems are limited and are computationally intensive.



Since a complete integral basis for a general multi-loop amplitude is still unknown it has been difficult to proceed with the conventional unitarity construction of loop amplitudes. Nevertheless progress in understanding the planar basis by explicit construction of unitarity compatible IBP relations [12, 13]. This procedure has been a key ingredient in the maximal unitarity method proposed by Kosower and Larsen [14], where master integral coefficients are computed directly from on-shell cuts into tree-level amplitudes after finding a specific contour of a complex integral. The method has now been applied to a number of high-multiplicity cases of the two-loop planar double box [15, 16, 17] and also to non-planar topologies [18].

The extension of the OPP integrand reduction method was first proposed by Mastrolia and Ossola [19]. A key insight in how to construct a general procedure has come from algebraic geometry. As we will see, the problem reduces to operations on (quadratic) polynomial equations with many variables. Such operations are well known in mathematics and many techniques, and programs, now exist for dealing with computational algebraic geometry problems. The most well known of these is Buchberger's algorithm for the construction of Gröbner bases used in multivariate polynomial division. This direction has been explored both in the context of reducing Feynman diagrams [20, 21, 22] and using generalized unitarity [23, 24, 25]. More formal applications of computational algebraic geometry have also been used to analyze the general structure of two and higher loop amplitudes [26, 27].

In these proceedings we present a brief overview of the methods presented in refs. [23, 24, 25]. We describe the integrand reduction procedure valid in D -dimensions for an arbitrary loop amplitude. We outline applications at both two and three loops and consider the outlook for the future.

2. Multi-loop integrand reduction with computational algebraic geometry

The multi-loop amplitudes in this paper will refer to colour ordered primitive amplitudes which have a fixed ordering for the external legs and a well defined set of internal propagators. An L -loop amplitude in D dimensions can be schematically represented as,

$$A_n^{(L),[D]} = \int \prod_{i=1}^L \frac{d^D k_i}{(2\pi)^D} \frac{N(\{k\}, \{p\})}{\prod_{l=1}^{L(L+9)/2} D_l(\{k\}, \{p\})}, \quad (1)$$

where D_l are the denominators appearing in the parent topology, $\{k\}$ are the loop integration variables and $\{p\}$ are the external momenta. The numerator function N is a function of Lorentz products of loop momenta and external momenta, external polarization vectors or external spinor wave-functions. All kinematic quantities can be represented in terms of invariants or scalar products of loop momenta with vectors in a basis $\{e_1, \dots, e_4\}$ which can be constructed from the independent external momenta plus additional spurious directions as proposed by van Neerven and Vermaseren [28]. We will write all scalar products as $k_i \cdot e_j = x_{ij}$. We must also include extra-dimensional variables $\mu_{ij} = -k_i^{[-2\epsilon]} \cdot k_j^{[-2\epsilon]}$ where $k_i = \bar{k}_i + k_i^{[-2\epsilon]}$ and \bar{k}_i is the four-dimensional components of k_i . The number of propagators in $D = 4 - 2\epsilon$ dimensions is $L(L+9)/2$ coming from the $4L$ components of the 4-dimensional loop momenta plus an additional $L(L+1)/2$ components needed to describe the extra dimensional terms.

The integrand reduction method begins by reducing the numerator into terms proportional to the denominators D_l plus a remainder which is integrand of the topology in question,

$$N(\{k\}, \{p\}) = \Delta_M(\{k\}, \{p\}) + \sum_{l=1}^M a_l D_l(\{k\}, \{p\}) \quad (2)$$

where $M = \min(n + 3(L-1), L(L+9)/2)$ represents the maximum cut topology. The sum of terms proportional to denominators D_l will all vanish when all propagators are put on-shell.

To perform this operation we must first parametrise the integrand Δ_M in terms of rational coefficients multiplying monomials of scalar products,

$$\Delta_M(\{k\}, \{p\}) = \sum_k c_k \times \prod_{i=1}^L \prod_{j=1}^4 x_{ij}^{\alpha_{k;ij}} \prod_{l=1}^L \prod_{m=l}^L (\mu_{lm})^{\beta_{k,lm}}, \quad (3)$$

where c_i are rational functions of the external kinematics and the $\alpha_{k;ij}$ are integer powers of the Lorentz products x_{ij} and $\beta_{k,lm}$ are the powers of the extra dimensional parameters μ_{lm} . The notation used in eq. (3) includes an over complete set of scalar products (i.e. the matrices $\alpha_{k,ij}$ will contain many zeros). For this reason it is necessary to distinguish reducible scalar products (RSPs), which are linear in the propagators, from the irreducible scalar products (ISPs) which satisfy non-linear identities. The basis of ISPs can be constructed from the system of propagators as we will describe in the section 2.1.

Having obtained this representation we can then proceed to compute the coefficients c_k from products of tree-level amplitudes. This requires solving the system of on-shell constraints $D_l(\{k\}, \{p\}) = 0$ which will in general have many possible families of solutions $\{k^{(s)}\}$. Each of these solutions can be parametrised in terms of a number of free variables, $\{\tau_1, \dots, \tau_{L(L+9)/2-M}\}$ ¹. On each cut solution the integrand factorises into a product of tree-level amplitudes that can be written as a polynomial in τ_i ,

$$\Delta_M(\{k^{(s)}\}, \{p\}) = \prod_{V=1}^{M-1} A^{(0)}(V, \{\{k^{(s)}\}\}), \quad (4)$$

$$= \sum_i d_i \prod_{j=1}^{L(L+9)/2-M} \tau_j^{\gamma_{k;j}}, \quad (5)$$

where we have used a schematic notation for the on-shell tree-level amplitude appearing at each vertex of the cut-topology, $A^{(0)}(V, \{\{k^{(s)}\}\})$. $\gamma_{k;j}$ are integer powers that follow from $\alpha_{k;ij}$ and $\beta_{k,lm}$. Combining the information obtained in eqs. (3) and (5) allows to construct a linear system relating the rational coefficients of the integrand with the coefficients d_i that can be obtained from tree-level amplitudes,

$$B \cdot \vec{c} = \vec{d}, \quad (6)$$

where $\vec{d} = (d_1^{(1)}, \dots, d_{n(1)}^{(1)}, d_1^{(2)}, \dots, d_{n(S)}^{(S)})$ and $n(s)$ is the dimension of the on-shell solution s of which there are total of S solutions. Solving this system then determines the coefficients c_i .

After the form of Δ_M has been determined we can proceed with lower propagator cuts by performing the OPP style subtraction of the singularities,

$$\Delta_m(\{k^{(s)}\}, \{p\}) = \prod_{V=1}^{m-1} A^{(0)}(V, \{\{k^{(s)}\}\}) - \sum_{T=m+1}^M \sum_{t \in T} \frac{\Delta_{T;t}(\{k^{(s)}\}, \{p\})}{\prod_{l \in P(t)/P(m)} D_l(\{k^{(s)}\}, \{p\})} \quad (7)$$

where T is the number of propagators in the higher topologies. For each multiplicity we must sum over all the possible topologies t each of which will have a set of additional propagators in comparison with the topology m in question. The sum is represented as $P(t)/P(m)$.

If we restrict to the case of 4 dimensions $M = \min(n + 3(L - 1), 4L)$ and all coefficients $\beta_{k,lm} = 0$. There are two important details that need to be clarified before we can show that such an approach will work:

¹ The existence of this parametrization is only guaranteed in $D = 4 - 2\epsilon$ dimensions. In cases where the cut equations have genus > 1 this form will not exist (for example the 10-leg two-loop double box with 2 massless momenta at each of the 6 vertices).

- (i) How do we define the integrand parameterisation? (i.e. compute $\alpha_{k,ij}$ and $\beta_{k,lm}$)
- (ii) How can we determine the number of on-shell solutions S to define \vec{d} ?

2.1. Integrand parameterisation

The solution to first question above has been presented using multi-variate polynomial division with Gröbner basis Zhang [24] and also Mastrolia, Mirabella, Ossola and Peraro [20].

In this section we will construct, using the language of algebraic geometry, an ideal from the set of propagators for a topology with m denominators,

$$P = \{D_1, \dots, D_m\} \quad (8)$$

whose integrand we will denote as Δ_m . Using the van Neerven Vermaseren construction we are able to construct a basis where external momenta $\{p_1, \dots, p_x\}$ span the physical space and the directions $\{\omega_1, \dots, \omega_4\}$, orthogonal to the physical space, span the spurious, or trivial, space:

$$e = \{e_1, \dots, e_4\}. \quad (9)$$

All of the Lorentz products appearing in the expansion of P can be written in terms of the basis products $k_i \cdot e_j = x_{ij}$ using the 4×4 Gram matrix $[G_4]_{ij} = e_i \cdot e_j$,

$$k_i \cdot v_j = (k_i \cdot e_1 \ k_i \cdot e_2 \ k_i \cdot e_3 \ k_i \cdot e_4) \cdot G_4^{-1} \cdot \begin{pmatrix} v_j \cdot e_1 \\ v_j \cdot e_2 \\ v_j \cdot e_3 \\ v_j \cdot e_4 \end{pmatrix} \quad (10)$$

and

$$k_i \cdot k_j = -\mu_{ij} + (k_i \cdot e_1 \ k_i \cdot e_2 \ k_i \cdot e_3 \ k_i \cdot e_4) \cdot G_4^{-1} \cdot \begin{pmatrix} k_j \cdot e_1 \\ k_j \cdot e_2 \\ k_j \cdot e_3 \\ k_j \cdot e_4 \end{pmatrix} \quad (11)$$

The above assumes that we have all external vectors in four dimensions and hence only obtain contributions from the $4D$ parts of the loop momenta $k_i \cdot e_j = \bar{k}_i \cdot e_j$. To find the constraints from the propagators on the integrand we will use eqs. (10) and (11) to translate eq. (8) into polynomials in x_{ij} . The RSPs could be identified from linear relations among the equations however an efficient method is to compute the Gröbner basis using a graded monomial ordering [24].

It is helpful to consider an example at this point so we take the case of the two-loop planar double box in four dimensions. The list of propagators is,

$$P = \{k_1^2, (k_1 - p_1)^2, (k_1 - p_1 - p_2)^2, k_2^2, (k_2 - p_4)^2, (k_2 - p_3 - p_4)^2, (k_1 + k_2)^2\} \quad (12)$$

and we will span k_1 and k_2 with the basis,

$$\vec{e} = (p_1, p_2, p_4, \omega) \quad (13)$$

where $\omega^\nu = 4i\varepsilon^{\nu 124}$. The Gram matrix is explicitly

$$G_4 = \frac{1}{2} \begin{pmatrix} 0 & s & t & 0 \\ s & 0 & u & 0 \\ 0 & u & t & 0 \\ 0 & 0 & 0 & stu \end{pmatrix} \quad (14)$$

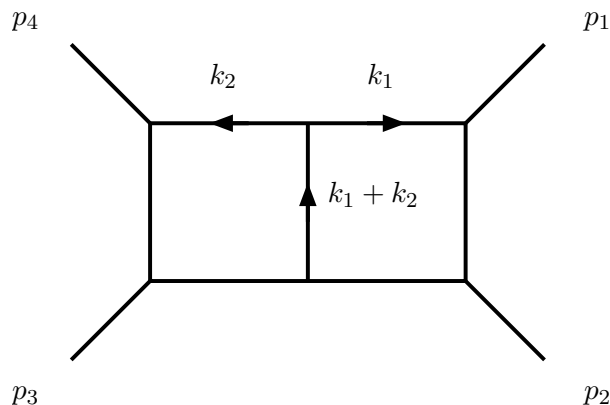


Figure 1. The two-loop planar double box configuration

where $u = -s - t$. It is then straightforward to apply eq. (10) and express P in terms of the 8 scalar products x_{ij} . We find four linear relations in the on-shell equations:

$$x_{11} = 0 \quad 2x_{12} + s = 0 \quad x_{23} = 0 \quad 2x_{22} + 2x_{21} + s = 0 \quad (15)$$

which leaves us with a set of 4 ISPs $\{x_{24}, x_{14}, x_{21}, x_{13}\}$ and an ideal

$$I = \langle -4tx_{13} + 4x_{13}^2 - 4x_{14}^2 + t^2, -4tx_{21} + 4x_{21}^2 - 4x_{24}^2 + t^2, \\ st(2x_{13} + 2x_{21}) + s(4x_{13}x_{21} - 4x_{14}x_{24}) + 8tx_{13}x_{21} - st^2 \rangle \quad (16)$$

which we should use to remove terms from the polynomial ansatz for Δ_7 restricted by imposing renormalizability constraints,

$$\Delta_7^{\text{ansatz}} = \sum_{a_1, \dots, a_4} C_{a_1 a_2 a_3 a_4} x_{14}^{a_1} x_{24}^{a_2} x_{13}^{a_3} x_{21}^{a_4} \quad (17)$$

where $a_1 + a_3 \leq 4, a_2 + a_4 \leq 4$ and $a_1 + a_2 + a_3 + a_4 \leq 6$ which is a sum of 160 terms. The polynomial division operation based on Gröbner bases is a standard algorithm in computer algebra packages and so it is now straightforward to obtain:

$$\Delta_7^{\text{ansatz}}/G(I) = \Delta_7 = \sum_{i=1} c_i x_{14}^{\alpha_{i1}} x_{24}^{\alpha_{i2}} x_{13}^{\alpha_{i3}} x_{21}^{\alpha_{i4}} \quad (18)$$

which contains 32 terms defined by:

$$\alpha = ((0, 0, 0, 0), (0, 0, 0, 1), (0, 0, 0, 2), (0, 0, 0, 3), (0, 0, 0, 4), \\ (0, 0, 1, 0), (0, 0, 1, 1), (0, 0, 1, 2), (0, 0, 1, 3), (0, 0, 1, 4), \\ (0, 0, 2, 0), (0, 0, 2, 1), (0, 0, 3, 0), (0, 0, 3, 1), (0, 0, 4, 0), \\ (0, 0, 4, 1), (0, 1, 0, 0), (0, 1, 0, 1), (0, 1, 0, 2), (0, 1, 0, 3), \\ (0, 1, 1, 0), (0, 1, 1, 1), (0, 1, 1, 2), (0, 1, 1, 3), (0, 1, 2, 0), \\ (0, 1, 3, 0), (0, 1, 4, 0), (1, 0, 0, 0), (1, 0, 0, 1), (1, 0, 1, 0), \\ (1, 0, 2, 0), (1, 0, 3, 0)) \quad (19)$$

The algorithm presented here is implemented in the public Mathematica package BASISDET by Yang Zhang. We note the the problem of determining the integrand parameterisation is rather

efficient, especially when external kinematic quantities are computed numerically, and has been applied to many complicated multi-loop topologies in 4 and $4 - 2\epsilon$ dimensions. The integrand basis can be rather large for higher loop topologies. The three-loop triple box, shown in figure 2, considered in ref. [25] contains a total 398 monomials in terms of 7 ISPs. In this case IBP reduction was used subsequently to reduce the integrand onto a basis of three master integrals.

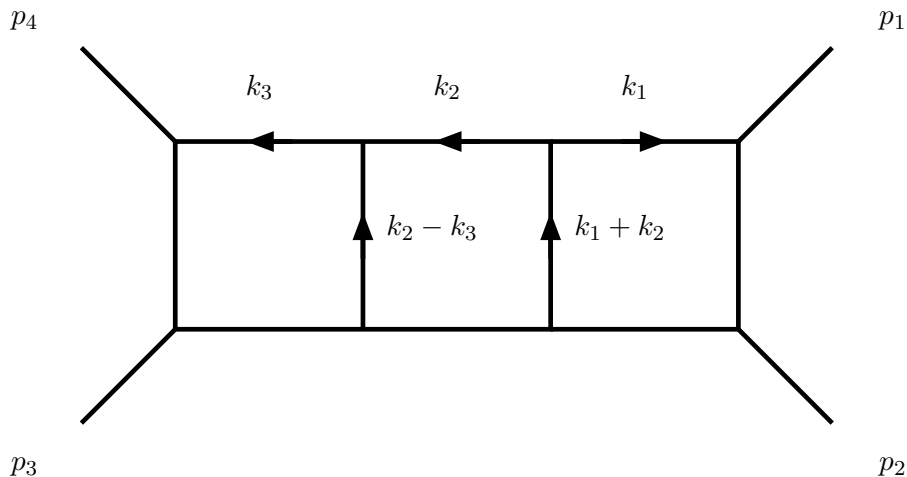


Figure 2. The three-loop triple box configuration

2.2. On-shell solutions

The solution to the second question has been formulated by Zhang in ref. [24]. In general the geometry of the solutions space can be very complicated but algebraic geometry again provides a technique to systematically analyze the zeros of the propagator ideal I , which will denote $\mathcal{Z}(I)$ or the zero locus of I . The famous Lasker-Noether theorem for primary decomposition of ideals proves that any ideal may be decomposed into a set of smaller primary ideals,

$$I = \bigcap_{s=1}^S I_s \quad (20)$$

where $\dim(I_s) = n(s)$. The zero locus of I follows from this since,

$$\mathcal{Z}(I) = \bigcup_{s=1}^S \mathcal{Z}(I_s). \quad (21)$$

Algorithms to compute the primary decomposition are available in the MACAULAY2 computer algebra package [29]. For the two loop double box considered in the previous section one can use this technique to reproduce the 6 independent solutions first obtained by Kosower and Larsen [14]. The three-loop triple box has a total of 14 solutions and explicit parameterisation have been presented in ref. [25].

2.3. D -dimensional cuts

One more aspect of the linear system in eq. (6) that remains is to prove that the matrix B does indeed have the property $\text{rank}(M) = \dim(\vec{c})$. This is not guaranteed in four dimensions. However, as shown recently in ref. [30], by going to D dimensions this problem is avoided at the cost of creating larger linear systems relating the integrand coefficients to the tree-level amplitudes. We have observed two cases where the procedure can break down:

- (i) The propagator ideal is not radical, often written $I \neq \sqrt{I}$. In this case there will be terms in Δ that cannot be determined from the product of trees as they will vanish on each cut solution $\{k^{(s)}\}$.
- (ii) The dimension of a branch of the cut solutions is not regular. For example, in four dimensions, we can find that $\dim(I_i) \neq 4L - m$ for all primary ideals in an m propagator cut. The two-loop pentagon-triangle is an example of this: though there are seven propagators we find two families of solutions with dimension 2 instead of dimension 1.

There is a possibility these problems could be avoided by considering different topologies simultaneously but it will complicate the integrand reduction procedure. In $4 - 2\epsilon$ dimensions however this problem is side stepped since one can prove that all ideals are radical ideals [30]. In this case there will always be exactly one solution to the on-shell constraints ($S = 1$).

In order to generate the extra-dimensional dependence in the tree amplitudes it is necessary at two-loops to use at least six dimensions. In general tree amplitudes will need to be computed in higher and higher dimensions as we increase the loop order.

3. Outlook

We have presented an overview of an integrand reduction procedure valid in D dimensions which allows the rational coefficients of the integrand monomials to be extracted from tree-level amplitudes. The method has been applied in the recent computation of the planar all-plus helicity five-gluon two-loop amplitude, the first example of a $2 \rightarrow 3$ helicity amplitude in QCD [30].

There are still many questions as to whether the method can be an efficient alternative to integration-by-parts identities in the future. Of course the integrand basis can always be further reduced by IBP to reduce the number of integrals that must be computed. There is also an open question concerning the treatment of doubled propagator diagrams that can appear in the input. In this case it has been shown that there is no obstacle in performing the polynomial division [22] yet the on-shell cuts into tree amplitudes will need regulating in order to follow the approach suggested here.

The use of algebraic geometry in multi-loop computations may have many applications in the future both for integrand computations and in the computations of the integrals themselves. Progress in the computation of the loop integrals themselves, for example the differential equation of Henn [31], give hope that precision predictions of multi-particle states at NNLO may be possible within the near future.

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