

Circuit level, 32 nm, 1-bit MOSSI-ULP adder: power, PDP and area efficient base cell for unsigned multiplier

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Abstract: Scaling the feature size under 0.1 micron leads to the domination of leakage power along with the consumption by interconnects also. Among alternatives for power optimization, the circuit level approach is the best option for low power consumption which is used in this work to incorporate a novel MOS Switch Integrated Ultra-Low Power (MOSSI-ULP) 1-bit full adder using 32 nm BPTM file. It is then used as a base cell in an array multiplier. Its performance is compared with similar kind of adders like SERF, ULPFA, TGA, TFA and BBL-PT. The analysis shows that the proposed design consumes a maximum of 11 times lower average power than ULPFA. Though MOSSI-ULP switches with a moderate frequency, it requires 17 times less PDP than ULPFA and maximum of 3 times less area.

Keywords: MOSSI-ULP, low power, PDP, area efficient, VLSI, adder

Classification: Integrated circuits

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1 Introduction

Feature size scaling yields several benefits for low-power operation. The improved device performance with V_{dd} under 1 V is one of them. It also ensures the circuit to operate with lower threshold voltage (V_t). Improved interconnect technology is used to minimize the parasitic effects. Reduced junction capacitance is another result of scaling the device. Availability of multiple and variable threshold devices lead to MTCMOS technology. This results in good management of active and standby power trade-off and higher density of integration. The sleep transistor and power gating are the familiar methods in connection with scaling the devices.

Apart from scaling of CMOS devices, high performance is more complex to obtain due to the effect of velocity saturation [1]. At short channel length L , the drain current per channel width (W) is no more equal to $(V_{gs} - V_t)^2/L$, where V_{gs} is the gate-to-source voltage and V_t is the threshold voltage. The drain current is proportional to $V_{sat} * (V_{gs} - V_t)$, where V_{sat} is the saturation velocity [2]. At sub-micron technology the difference between V_{dd} and V_t is very small and the static leakage is dominant to increase the static power consumption. This situation increases the design complexity and hence care is needed. This paper implements the 1-bit full adder by integrating PMOS and NMOS transistors as switches using 32 nm model file named as MOS Switch Integrated Ultra Low Power full adder (MOSSI-ULP). This is used as base cell in a multiplier array and compared with five more adders of this kind.

2 Previous un-conventional 1-bit full adders

This section deals with the previous non-static CMOS adder implementations. In a conventional CMOS circuit, all the MOS transistors are swinging with the help of supply rail with respect to ground. On the contrary, the methods taken for comparison with MOSSI-ULP are propagating mostly with the applied input signals and at few nodes the V_{dd} is applied (static inverters) to form hybrid kind of designs. The Transmission Gate Adder (TGA) and Transmission Function Adder (TFA) are similar circuits which use the back to back coupled PMOS and NMOS devices to offer dynamic switching [3]. The Branch Based Logic and Pass Transistor combined (BBL-PT) adder and Ultra Low Power Full Adder (ULPFA) use the *XOR* function for sum circuits and static CMOS circuit for carry [4]. The Static Energy Recovery Full adder (SERF) eliminates the path to ground and avoids short circuit power P_{sc} [5]. These are implemented as base cells individually in a multiplier to analyze and compare with the proposed MOSSI-ULP adder.

3 Proposed MOSSI-ULP 1-bit adder

The transmission gates use the parallel combination of PMOS and NMOS to get a specific function. But the proposed MOSSI-ULP adder has two PMOS transistors at the first stage, a static inverter at the intermediate stage and two pairs of back to back coupled PMOS and NMOS transistors at the output end: one pair for the sum output S and the other for the carry output C_{out} as in Fig. 1 (a). The remaining part of this section describes the functions and derivations of the proposed adder.

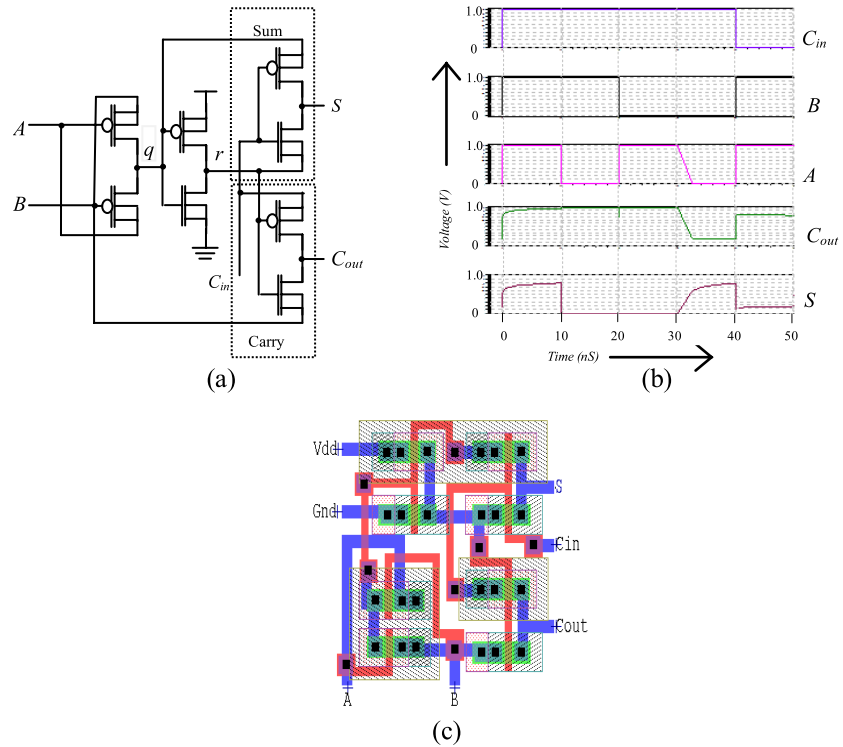


Fig. 1. (a) Proposed MOSSI-ULP Adder (b) Output waveform and (c) Layout

The drains of the two PMOS transistors at the first stage are coupled together and the node is named as q . The input A is applied to gate of the PMOS at the top as well as source of the bottom transistor. The other input B is applied to the gate of the PMOS transistor at the bottom and also to source of the upper PMOS. The upper transistor is switching $A'B$ while the other predicts AB' . At node q these two signals are combined to give a wired *or* function. It is expressed as $q = A'B + AB'$ which is a *xor* function constructed with the help of PMOS transistors as switches only (proved as in Fig. 3 (f)). When $V_g = 0$, the PMOS conducts and yields $V_d = V_{gs}$. Where V_d is the drain voltage which is at logic '0' for $V_{gs} = 0$ and logic '1' for $V_d = 1$ V. On the other hand the PMOS transistor is at cut-off for $|V_{gs}| < |V_{TP}|$ and hence it offers V_{TP} as logic '0' termed as weak '0' [6, 7]. Here V_{TP} is the threshold voltage of PMOS (for 32 nm, $V_{TP} = -0.35$ V). These two conditions are expressed as

$$V_{gs} < V_{TP} \text{ or } |V_{gs}| > |V_{TP}| ; \text{ for conduction} \quad (1)$$

$$V_{gs} > V_{TP} \text{ or } |V_{gs}| < |V_{TP}| ; \text{ for cut-off} \quad (2)$$

At the node r the static inverter gives the complement of q as $r=q'=(A'B+AB')'$ which is a *xnor* function. With respect to C_{in} , the *Sum* output ' S ' is expressed as

$$S = C_{in}'q + C_{in}r$$

$$S = C_{in}'(A'B + AB') + C_{in}(A'B + AB')' = A \oplus B \oplus C_{in} \quad (3)$$

The carry output C_{out} uses the signal from node r which is computed for the sum output. This is a logic sharing approach and also one of the circuit level low power techniques. The node r is connected with the gate terminals of PMOS and NMOS transistors of the bottom pair: the path in which the drains of the pair are coupled together to produce C_{out} . The sources of the pair are driven by inputs C and B respectively to yield a wired *or* function C_{out} at the coupled drains given as

$$C_{out} = r'C_{in} + rB \quad (4)$$

Since $r = q'$, it is also true that $q = r'$. Hence Eq. (4) can be written as

$$\begin{aligned} C_{out} &= qC_{in} + rB \\ &= (AB' + A'B)C_{in} + (A'B' + AB)B \\ &= AB'C_{in} + A'BC_{in} + A'B'B + AB \\ &= A(B + B'C_{in}) + A'BC_{in} + 0 \\ &= A(B + C_{in}) + A'BC_{in} \\ C_{out} &= AB + BC_{in} + AC_{in} \end{aligned} \quad (5)$$

Eq. (3) and Eq. (5) are the sum and carry outputs of MOSSI-ULP, proved true and the signals are as in the Fig. 1 (b) while Fig. 1 (c) represents its full custom layout.

4 Simulation results and comparative analysis

32 nm Berkeley Technology Model (BPTM) file is used to integrate the adders taken for consideration and they are used as base cells to test and analyze the performances. Fig. 2 shows the simulation test set up to validate the performances of the various logic methods. This top module uses two inputs AND gates to perform multiplication (product terms) which are represented in two forms. However, these are the same structures: one is the logic symbol and the other one is the dot product notation, represented to manage the space. The half adder is also a commonly used cell. 1-bit full adder is the Design Under Test (DUT) for which the discussed adders are used individually to simulate. All the methods use the same test environment with a supply voltage of $V_{dd} = 1\text{ V}$ to find out the results uniformly as well as most accurately. Depending on the depth of vector merging cells the full adders and half adders (Fig. 3 (e)) are used.

Transistor sizing is the key factor which decides the value of V_t and reduces power consumption as well as the circuit's switching activity [8]. Here all the methods use the same channel width profiles. The width of NMOS transistors are taken twice that of the feature size (channel length: 32 nm). For PMOS the width is double that of the NMOS. The results of

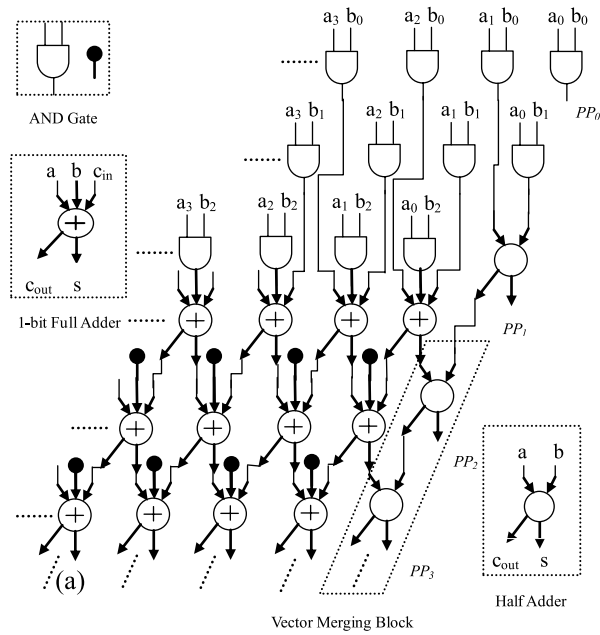


Fig. 2. Multiplier array using RCA blocks to test the adders

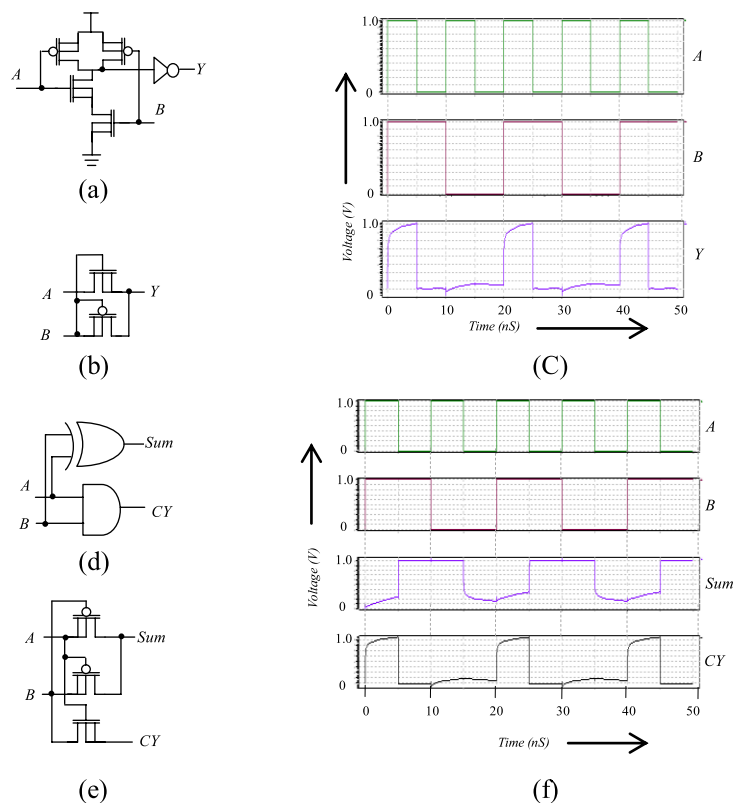


Fig. 3. (a) Static CMOS *AND* gate (b) *AND* gate with only two transistors (c) Voltage waveforms of 2T *AND* gate (d) Conventional half adder (e) Half adder used in the RCA with only three transistors (3T) (f) Voltage waveforms of the 3T half adder

adders chosen are listed in Table I.

The RCA structure uses 2T *AND* gates as in Fig. 3 (b) to produce partial products which needs reduced power than the consumption by

Table I. Simulation results and comparison of 1-bit full adders

Type of Adder ($V_{dd}=1V$)	Average Power (nW)	Delay (nS)	PDP (pWS)	Device Count	Number of Nodes	Area (λ^2)	Freq. (MHz)	Switching Factor A
MOSSI-ULP	27.32	22.66	0.62	171	794	1600	44.13	0.78
SERF	127.99	25.35	3.24	188	850	1920	39.45	3.82
TFA	202.98	4.87	0.99	236	1050	2880	205.34	0.94
TGA	185.94	4.91	0.91	268	1194	2880	203.67	0.76
BBL-PT	104.96	25.32	2.66	300	1370	4480	39.49	1.94
ULPFA	311.42	35.09	10.93	348	1594	5440	28.5	6.86

conventional gate as in Fig. 3 (a). For an $m \times n$ array with regular *AND* gates the transistors can be calculated as

$$TC_{AND} = 6mn \quad (6)$$

Where, TC_{AND} is the transistor count of *AND* gate, m is the number of digits of the multiplicand and n is for multiplier. For a 2T gate used in RCA, it is calculated as

$$TC_{AND} = 2mn \quad (7)$$

Hence the 2T structure needs 3 times fewer transistors than the regular gate and its waveforms are as in Fig. 3 (c). Fig. 3 (d) shows the conventional half adder which can be replaced with a circuit as in Fig. 3 (e) and its signals are given in Fig. 3 (f).

Short circuit (P_{sc}) and leakage (P_{leak}) power are the parts of total power dissipation. The former one happens when the output is in a meta-stable state while the latter is caused by the leakage current flow during the input voltage less than V_t [9].

At deep sub-micron technology, P_{sc} and P_{leak} are less important than the dynamic power due to its dominance and dependency on frequency given by

$$P_{dyn} = \alpha C_L V_{dd} * V_i f = \alpha C_L V_{dd}^2 f \quad (8)$$

In Eq. (8), at full swing signal the input voltage $V_i = V_{dd}$ and the terms are reduced as V_{dd}^2 . The circuit with smaller load capacitance, higher delay and few nodes will consumes low power. Hence with least nodes the proposed MOSSI-ULP consumes 13 times low power than the ULPFA.

Fig. 4 (a) shows the comparisons of average power with V_{dd} ranging from 0.8 V to 1.6 V. The TGA and TFA are switching faster than the other designs due to the availability of supply voltage at their intermediate nodes. For all the adder types, the delay (t_d) as a function of V_{dd} has no major variation because the difference in delay is distinct only when V_{dd} is less than V_t or different feature sizes are used for comparison. It is given as,

$$t_d = \left[\frac{1}{k_n W_{effn}} + \frac{1}{k_p W_{effp}} \right] \frac{C_L V_{dd}}{(V_{dd} - V_t)} \quad (9)$$

Where k is the trans-conductance and W is the width of MOS transistors. Suffix n and p refer NMOS and PMOS transistors respectively. The switching factor is calculated by taking number of nodes transferring from logic ‘0’ to ‘1’ in a unit time. Area is calculated using λ based design rules.

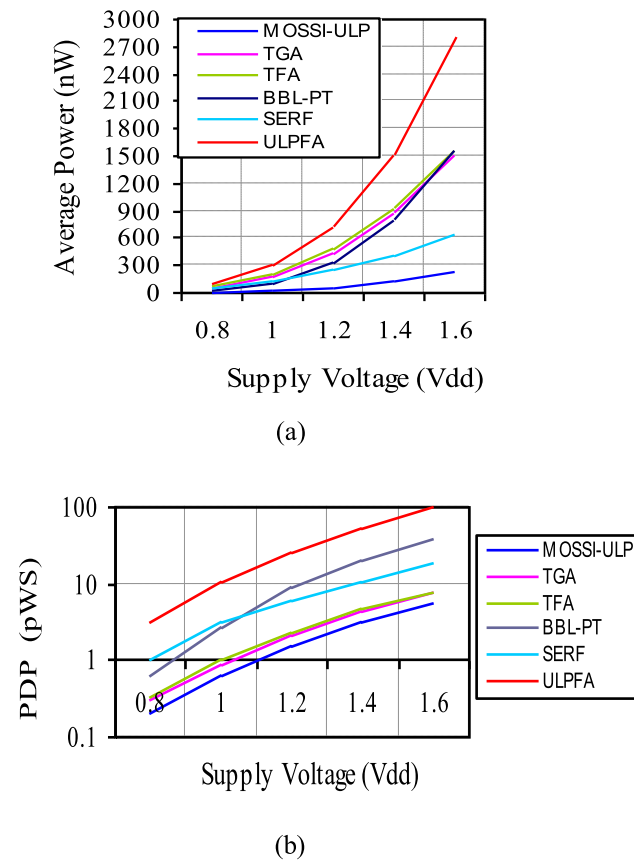


Fig. 4. Comparison of adders: (a) V_{dd} versus average power and (b) V_{dd} versus PDP

There are five layers horizontally: namely PMOS, NMOS, two metall layers and intermediate nodes with 4λ per layer and 4λ between adjacent layers [10]. Hence 8λ is the pitch. The proposed architecture utilizes the least area whereas the ULPFA requires more than thrice the requirement by the former one. The MOSSI-ULP adder is efficient also in terms of Power Delay Product (PDP) by referring Fig. 4 (b) which is given in logarithmic scale to view the comparison legibly.

5 Conclusion

Similar circuits which are passing the signals to output with the help of the applied inputs are simulated and compared with the MOSSI-ULP adder. With just 8 transistors and fewer nodes, the MOSSI-ULP consumes a minimum of 4 times lower power against the SERF adder. It switches moderately with nearly 5 times later than the TFA and TGA circuits. The TFA and TGA circuits are propagating faster due to the availability of supply voltage at their intermediate nodes. The Power Delay Product of the MOSSI-ULP is 30% lower than TFA and TGA. The proposed method operates with 17 times lowest PDP against the ULPFA. The layout area of the proposed architecture is 3.4 times compact than the other adders. The signal integrity issues may be the future scope of this work.